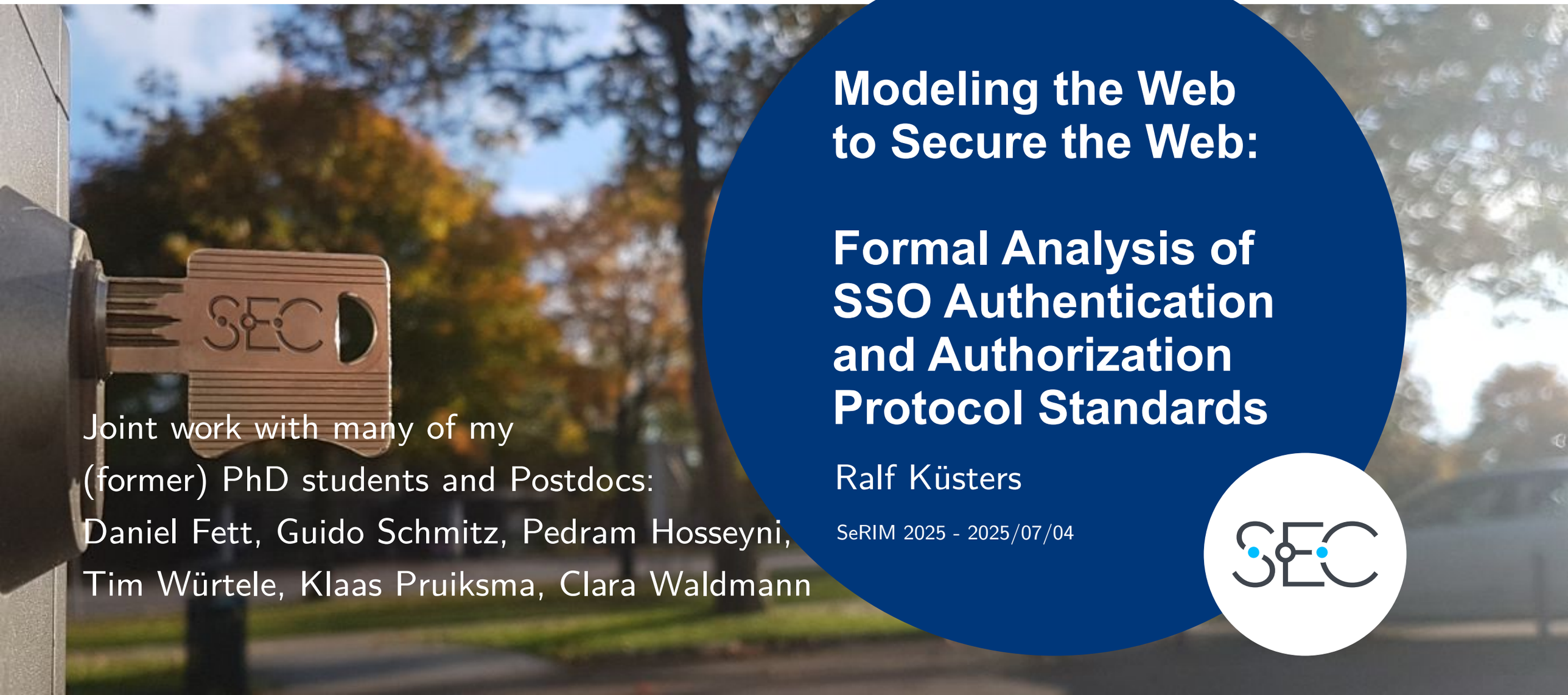
University of Stuttgart
Institute of
Information Security



Modeling the Web to Secure the Web:

Formal Analysis of SSO Authentication and Authorization Protocol Standards

Ralf Küsters

SeRIM 2025 - 2025/07/04

Joint work with many of my
(former) PhD students and Postdocs:
Daniel Fett, Guido Schmitz, Pedram Hosseyni,
Tim Würtele, Klaas Pruiksma, Clara Waldmann

Single Sign-On (SSO): Social Web SSO – The Beginning

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OAuth 2.0

OpenID Connect

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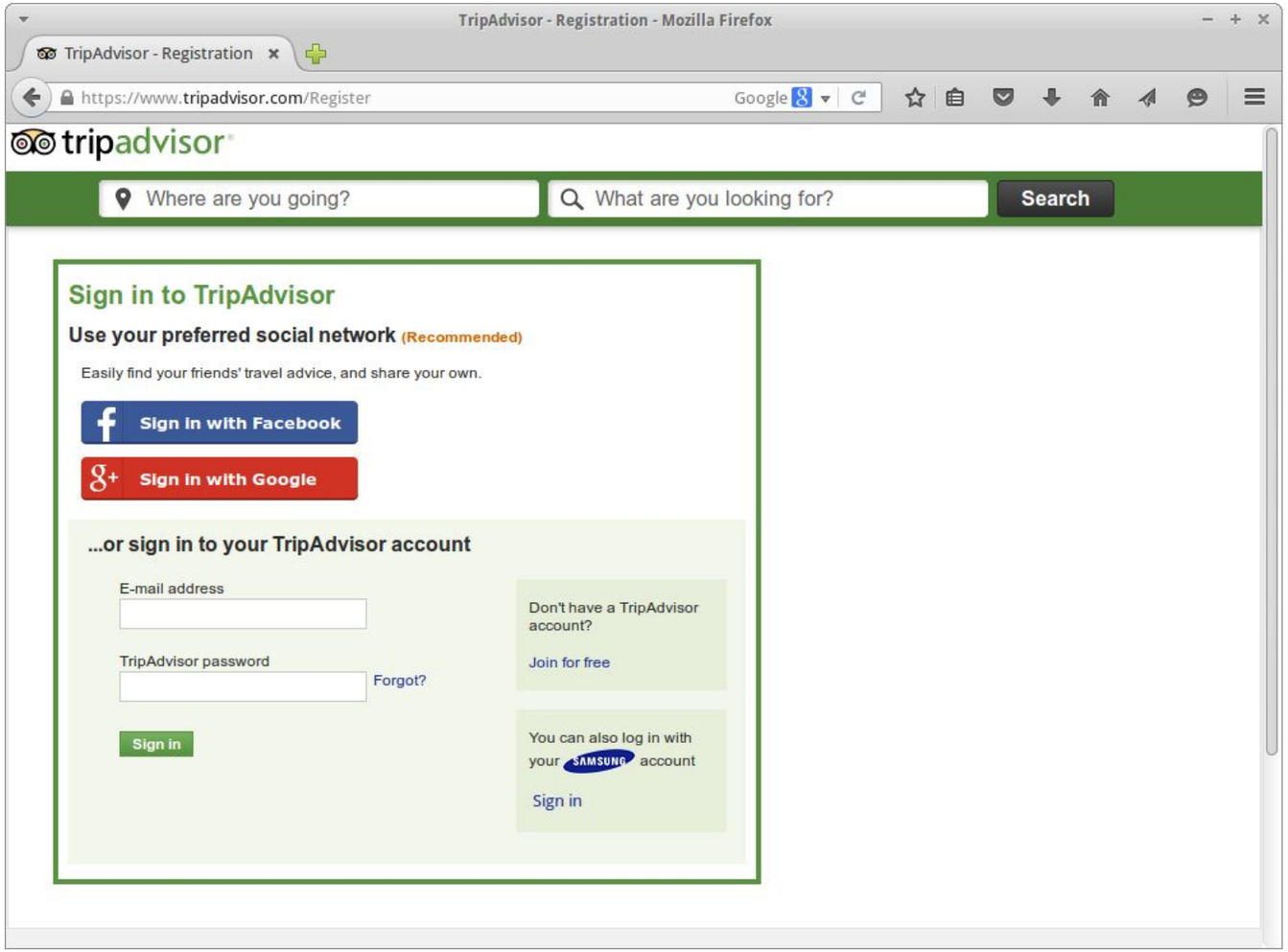

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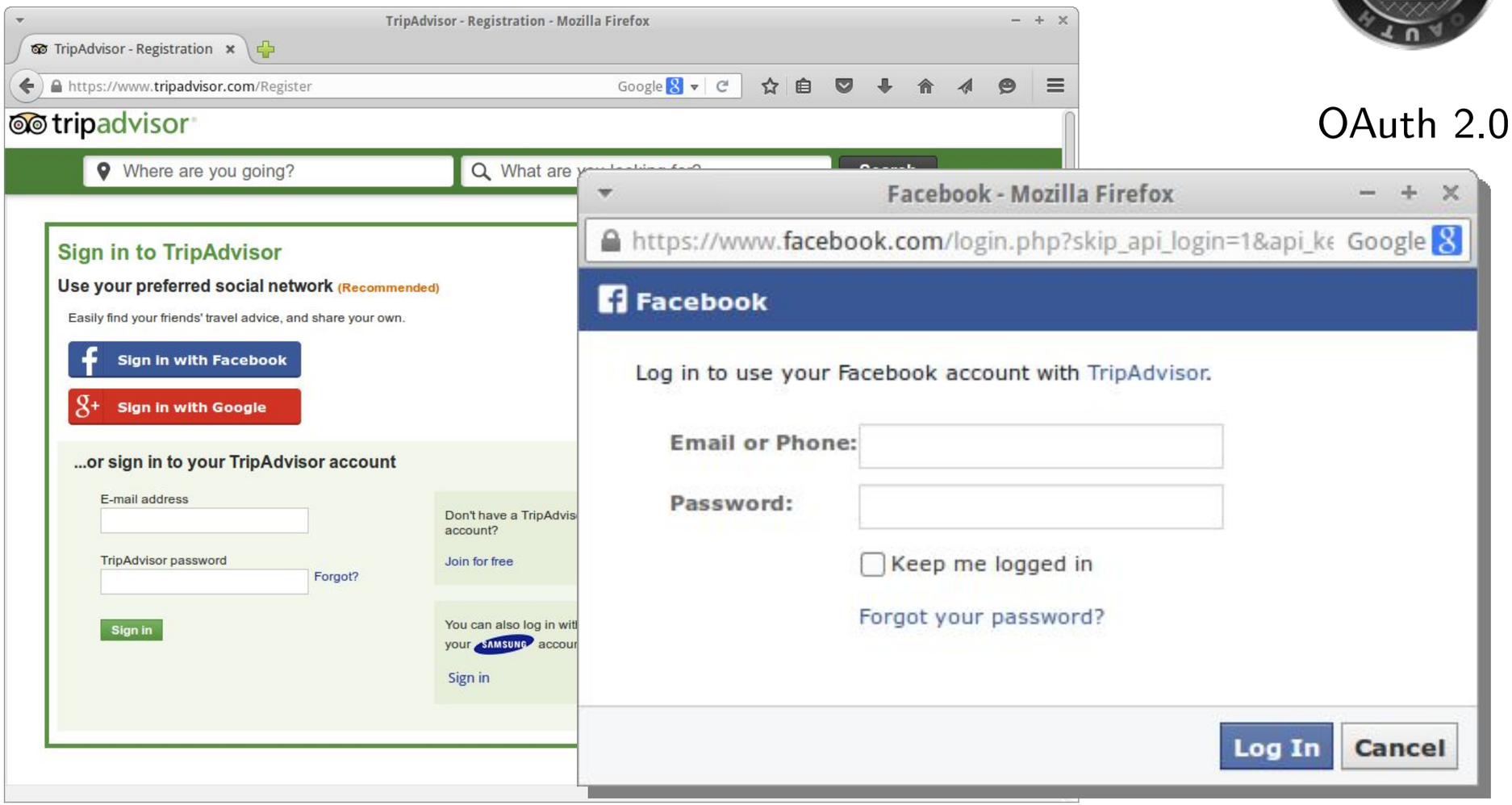

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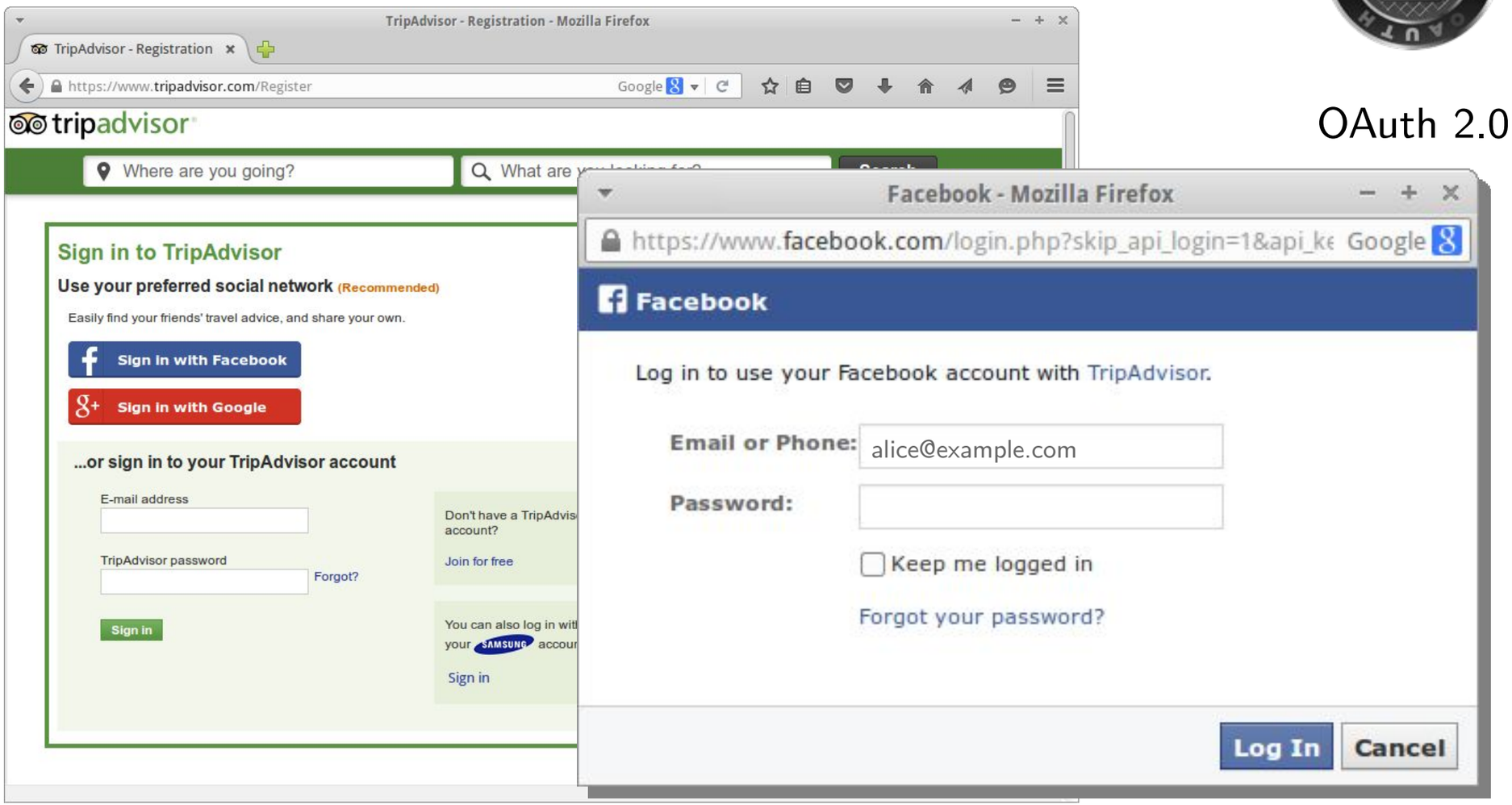

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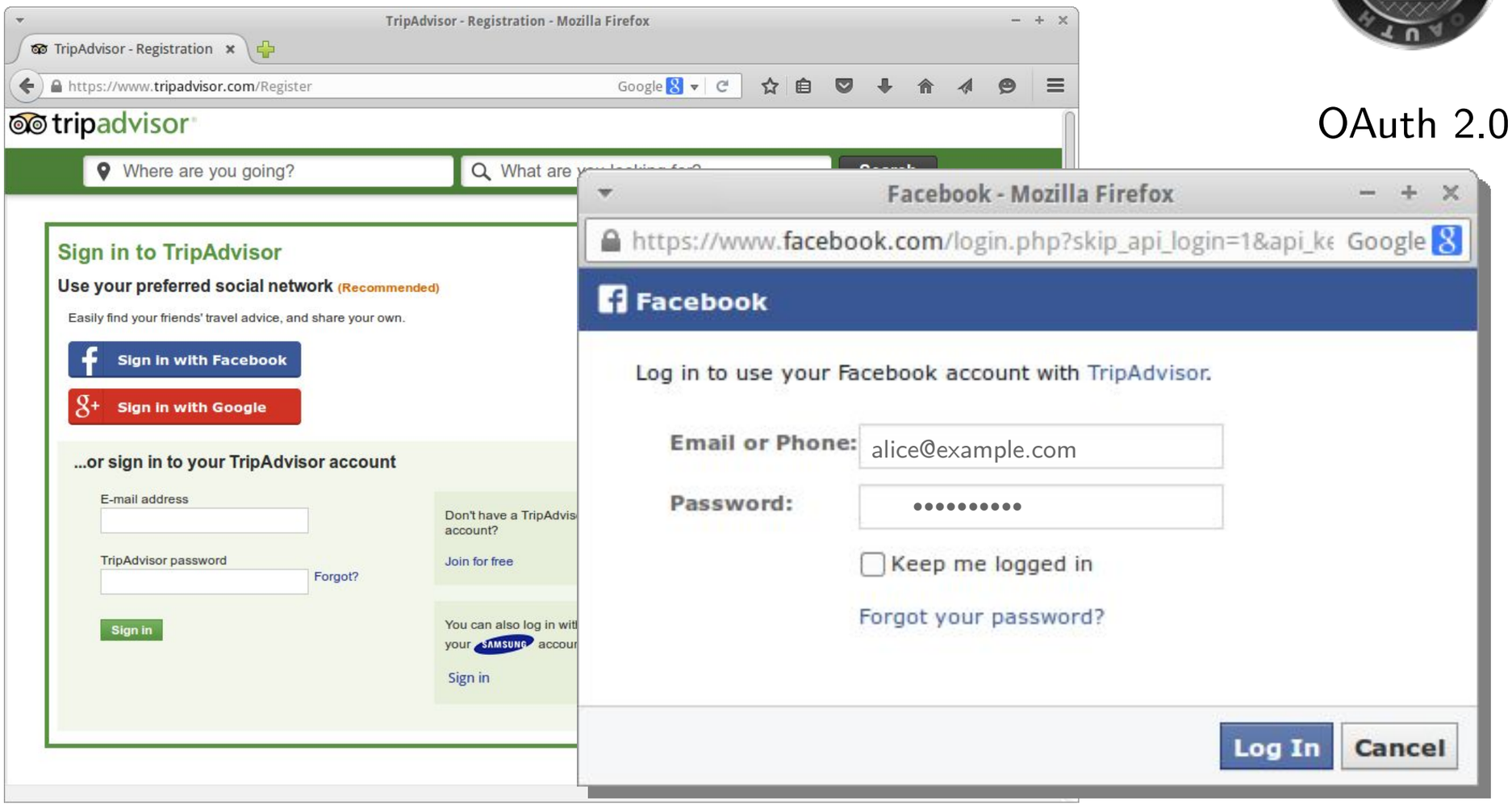

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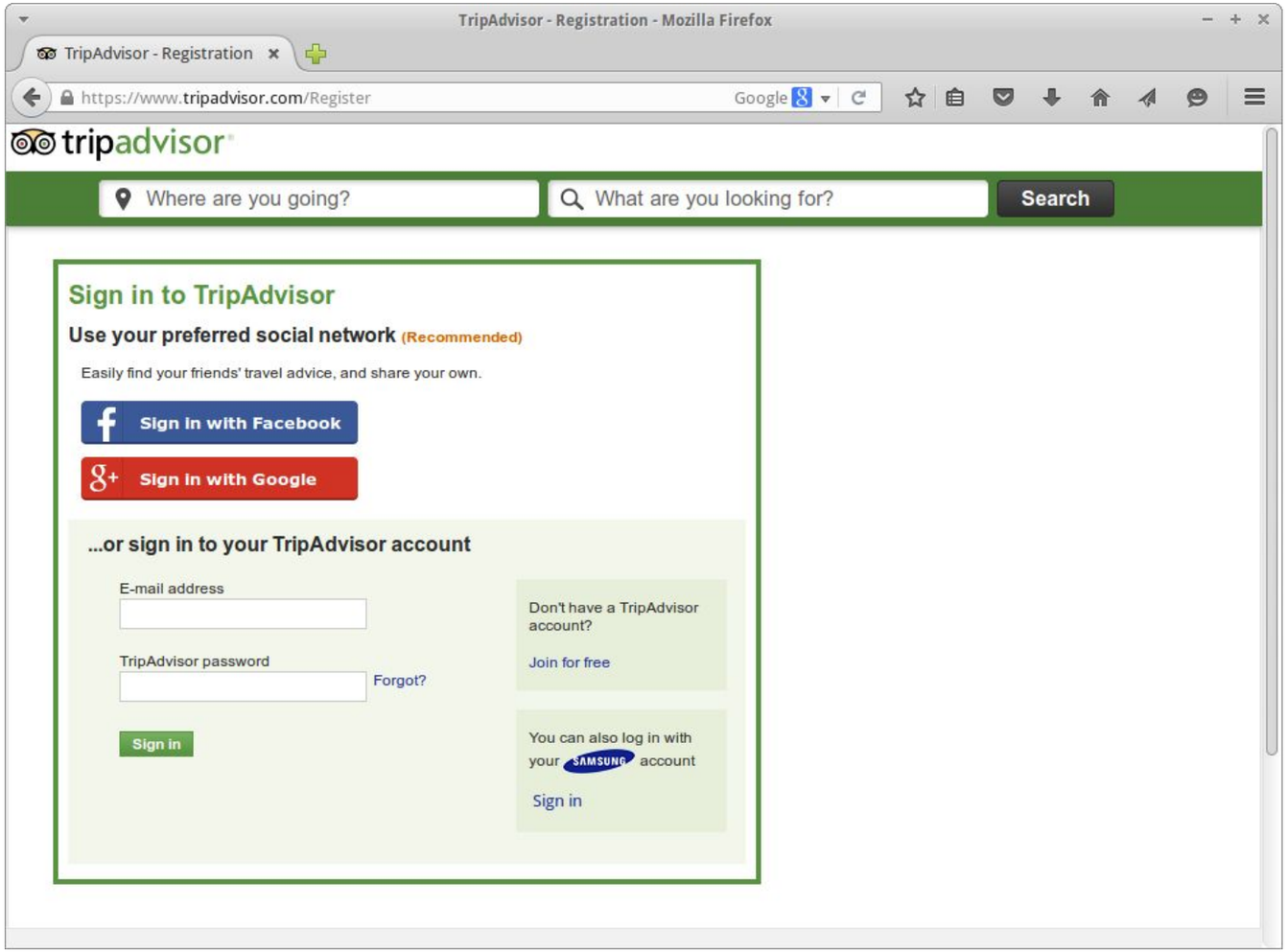

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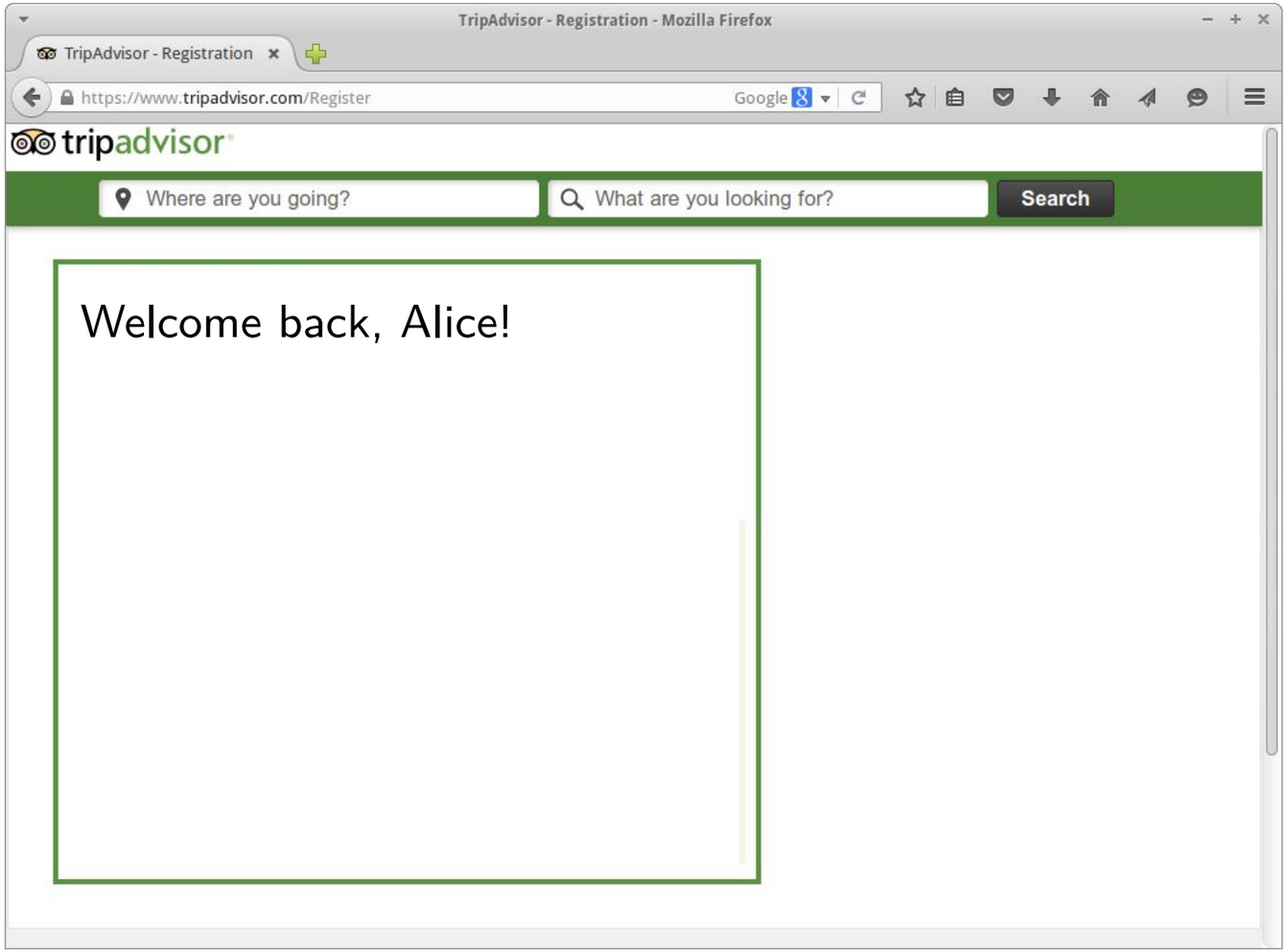

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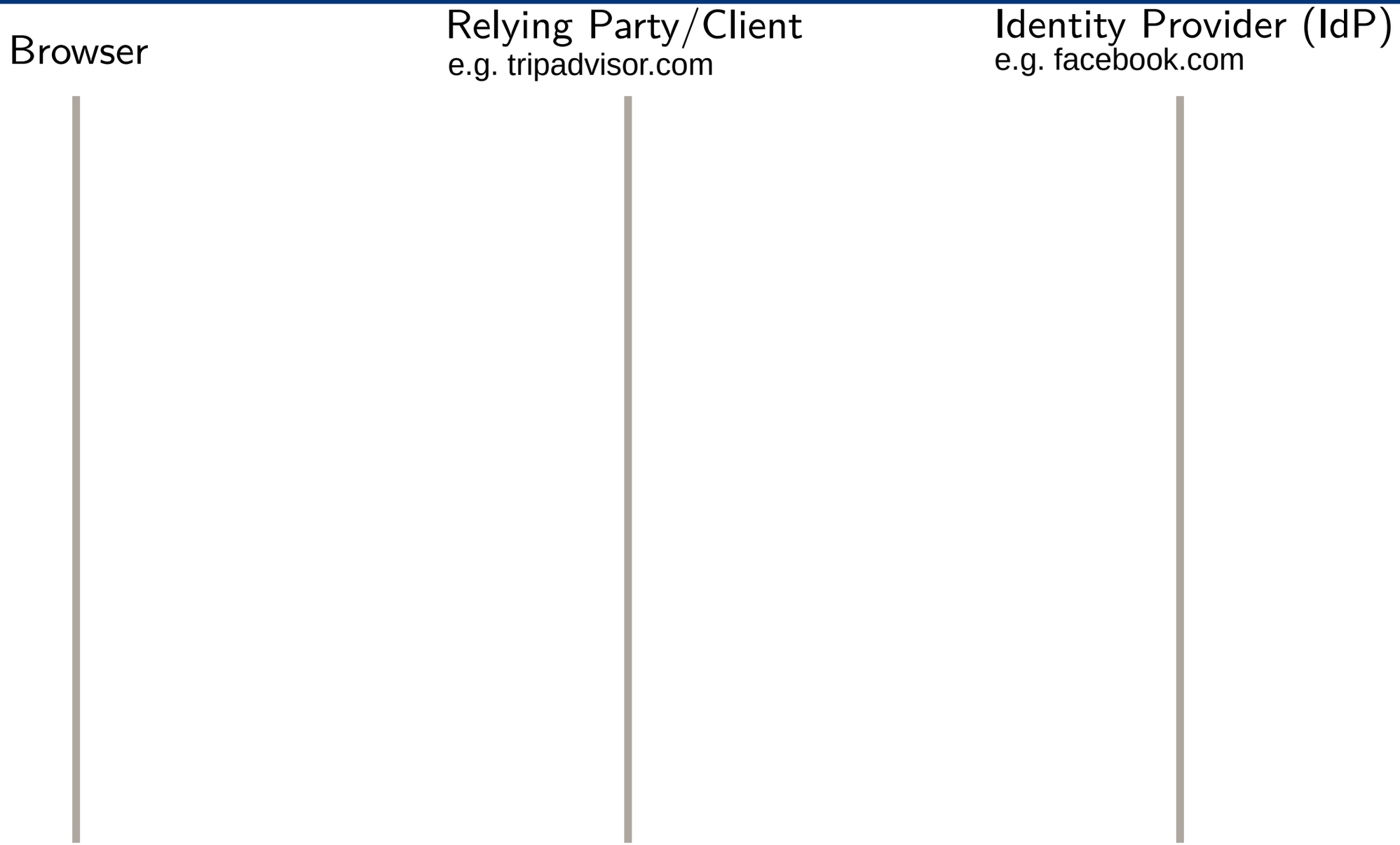
(Web) SSO: Basic Principle



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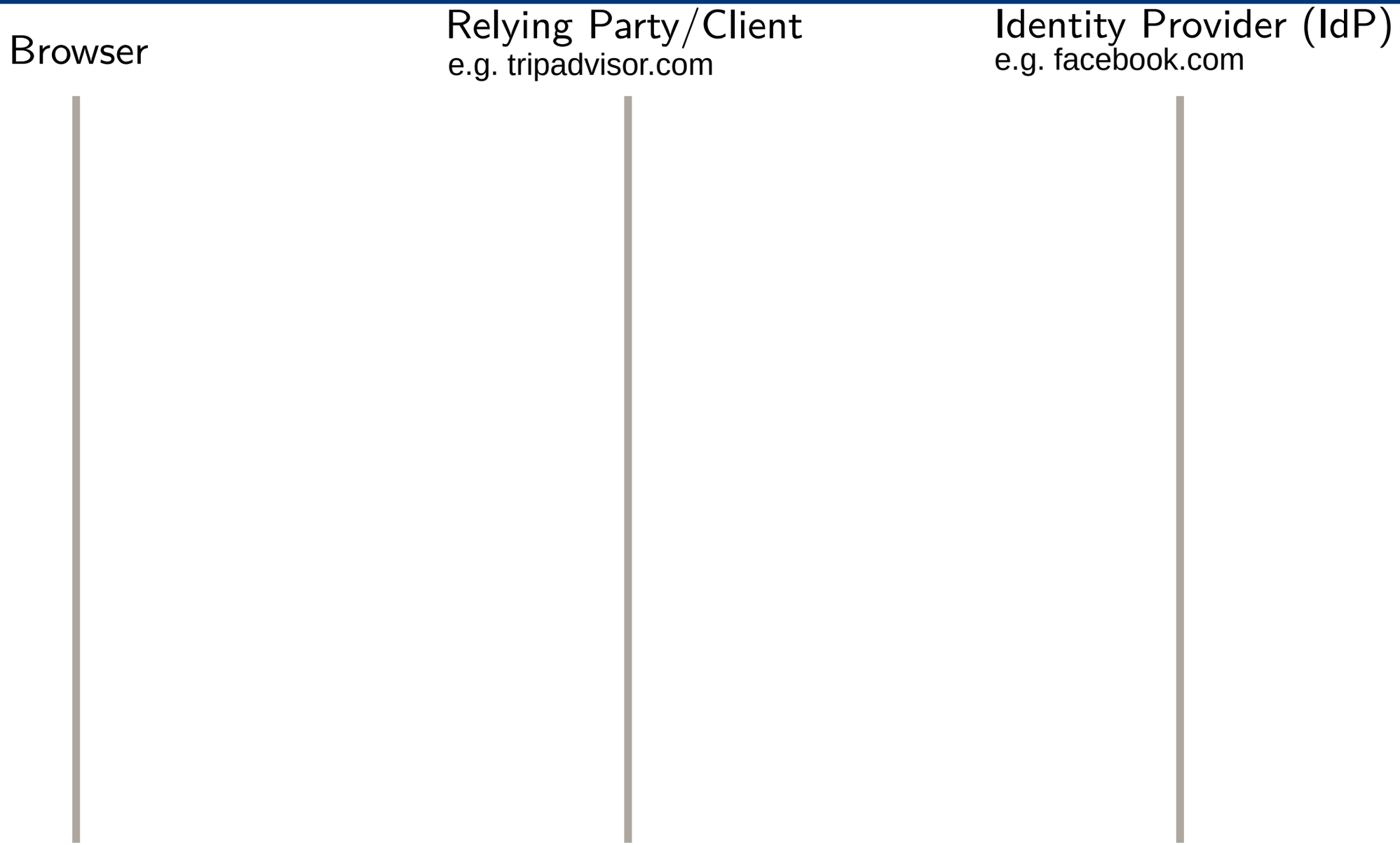

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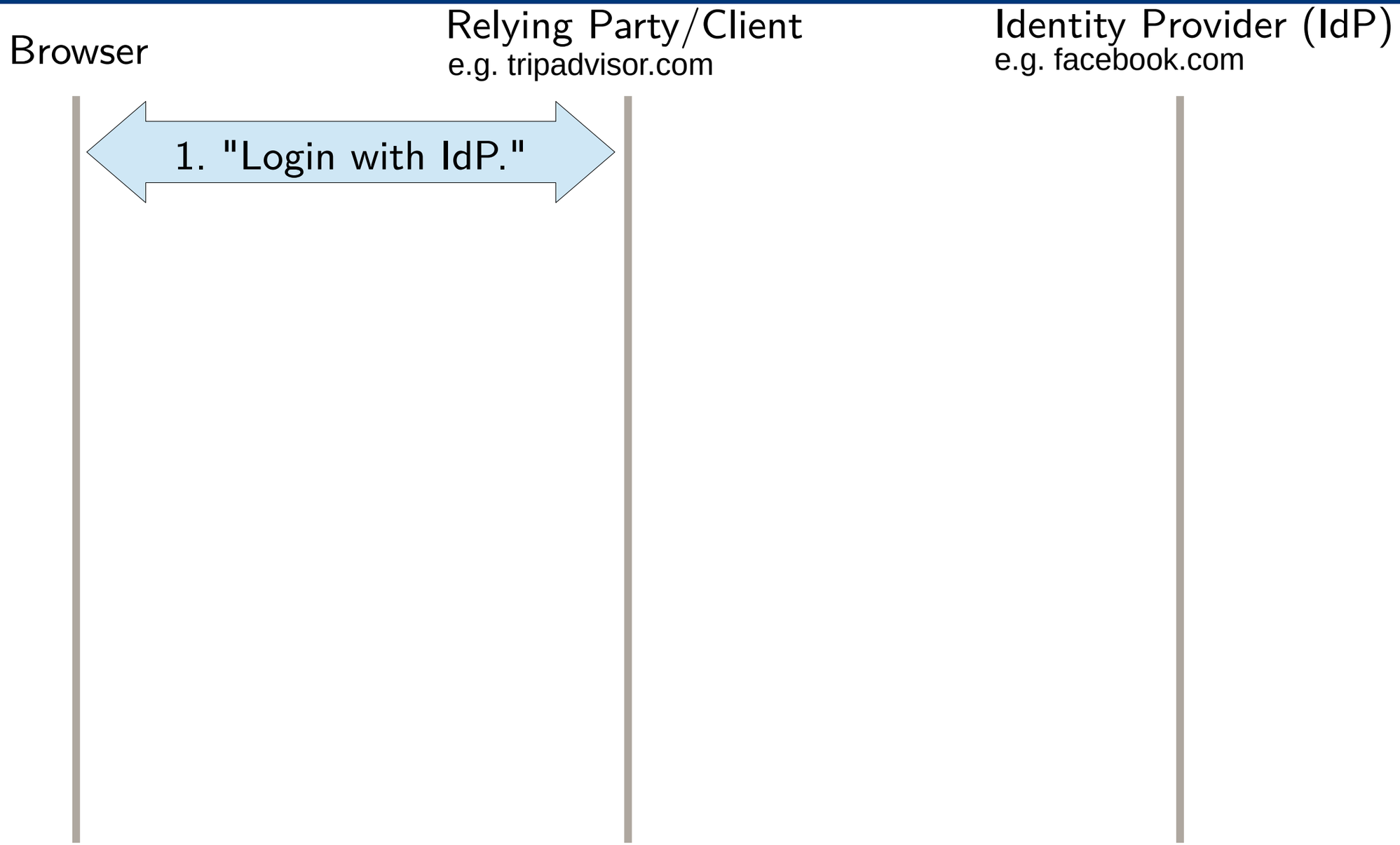

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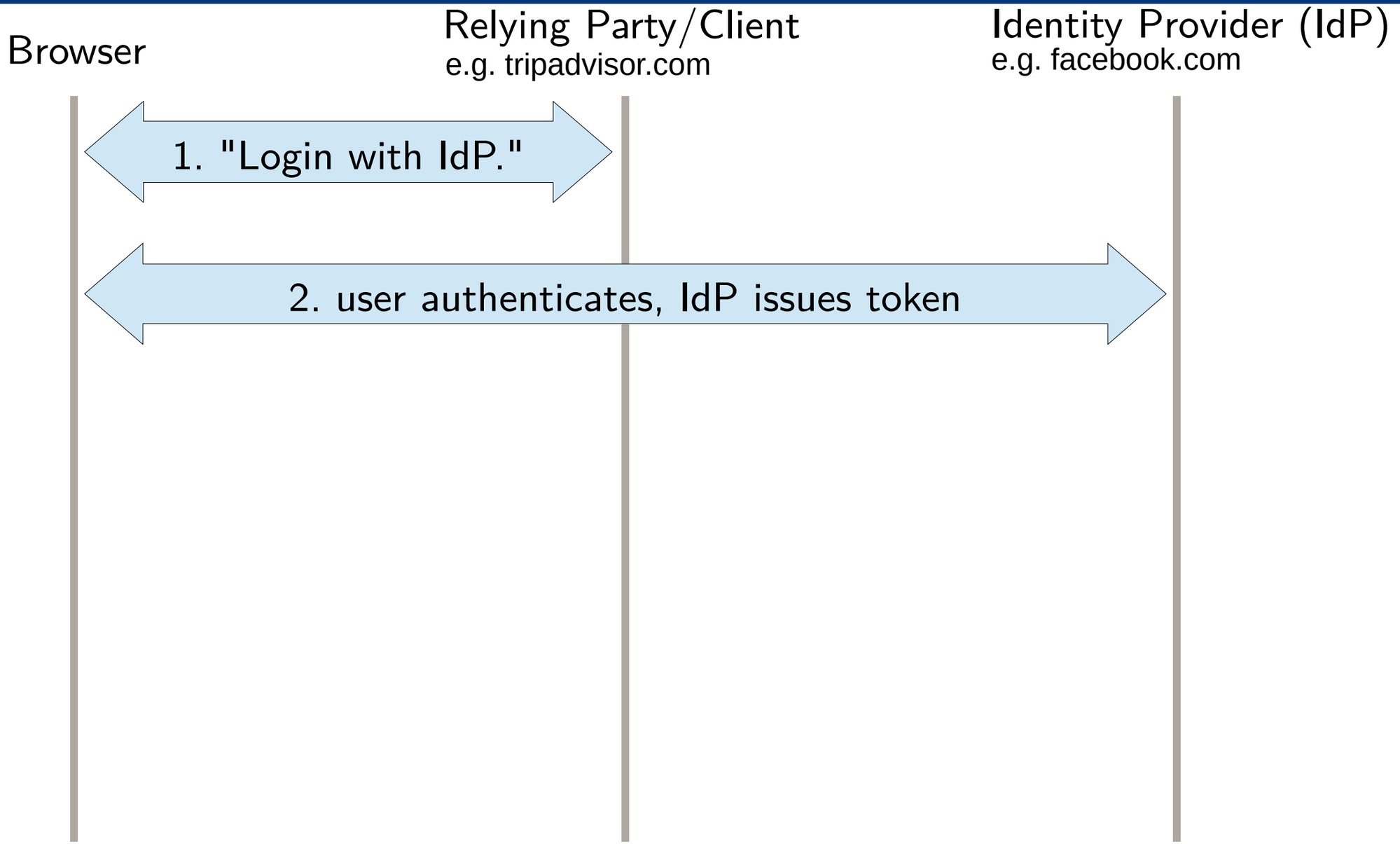

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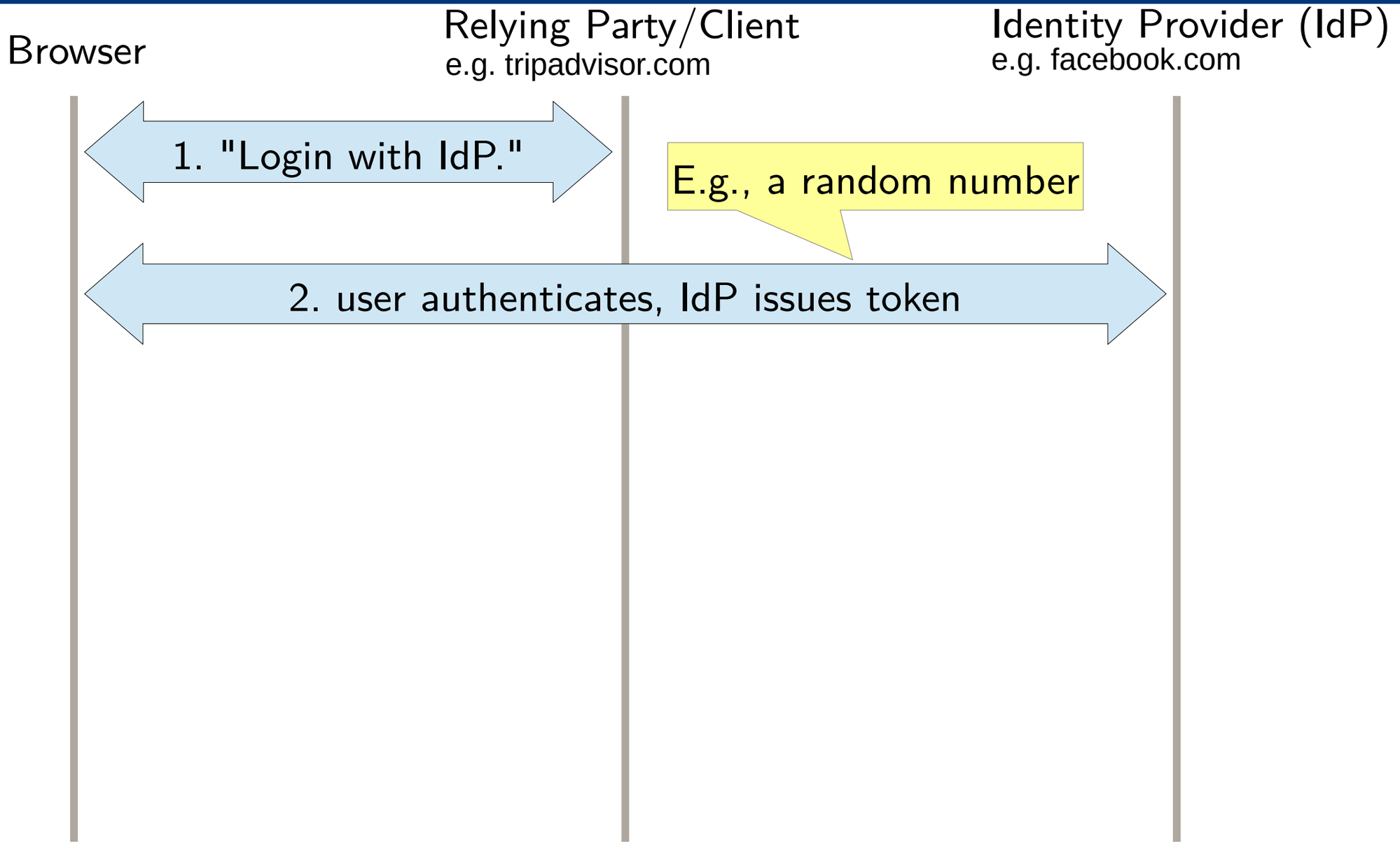

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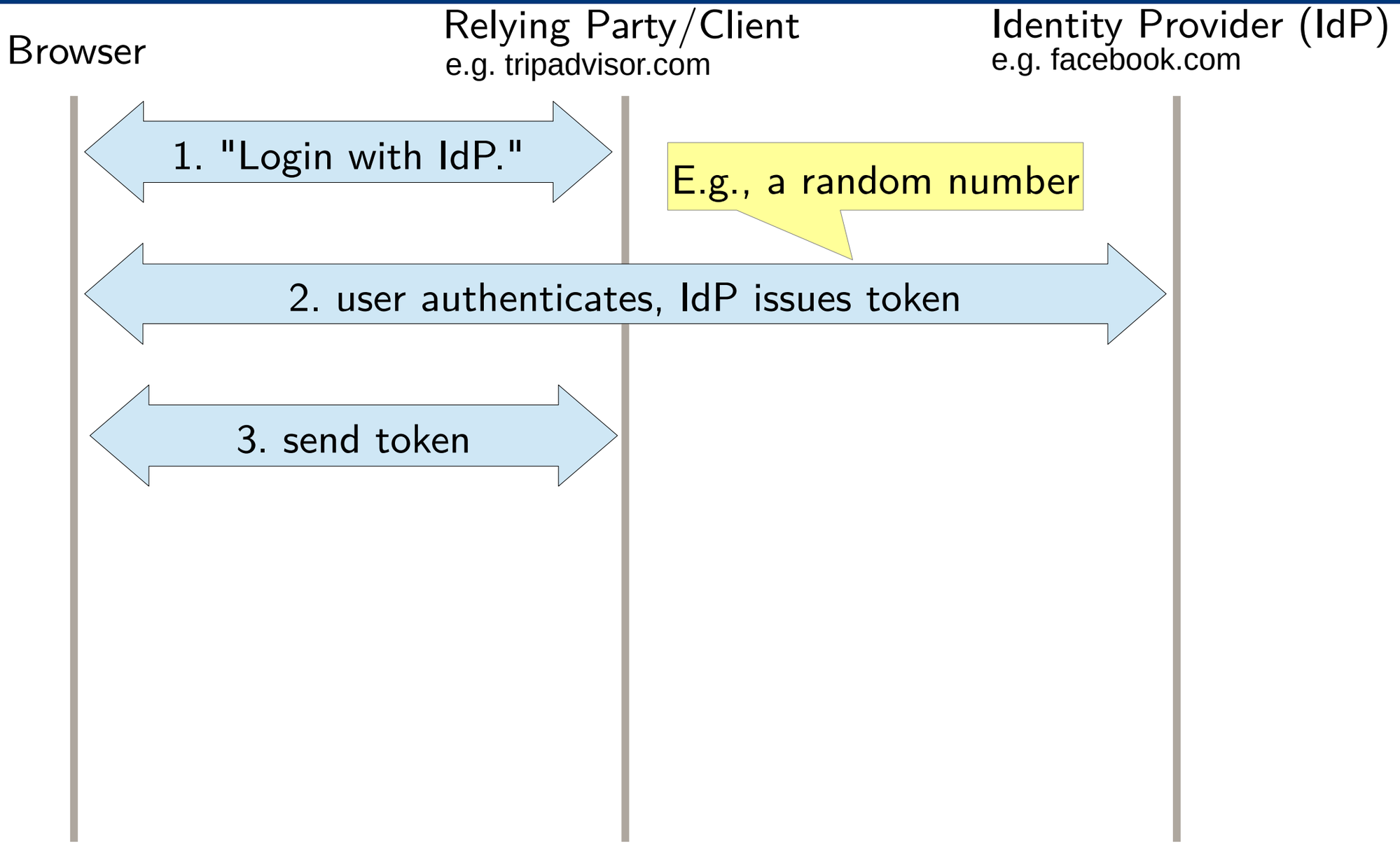

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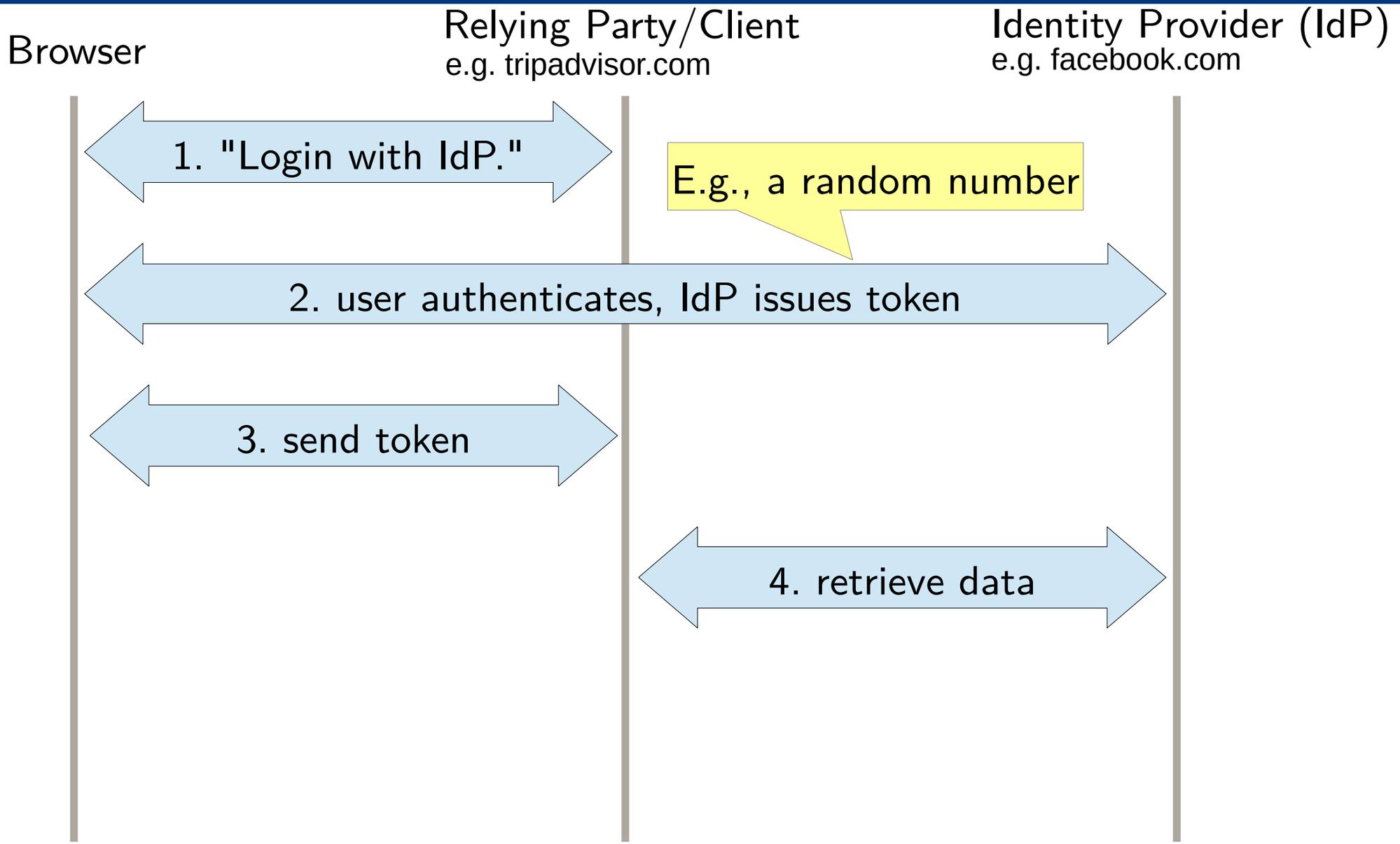

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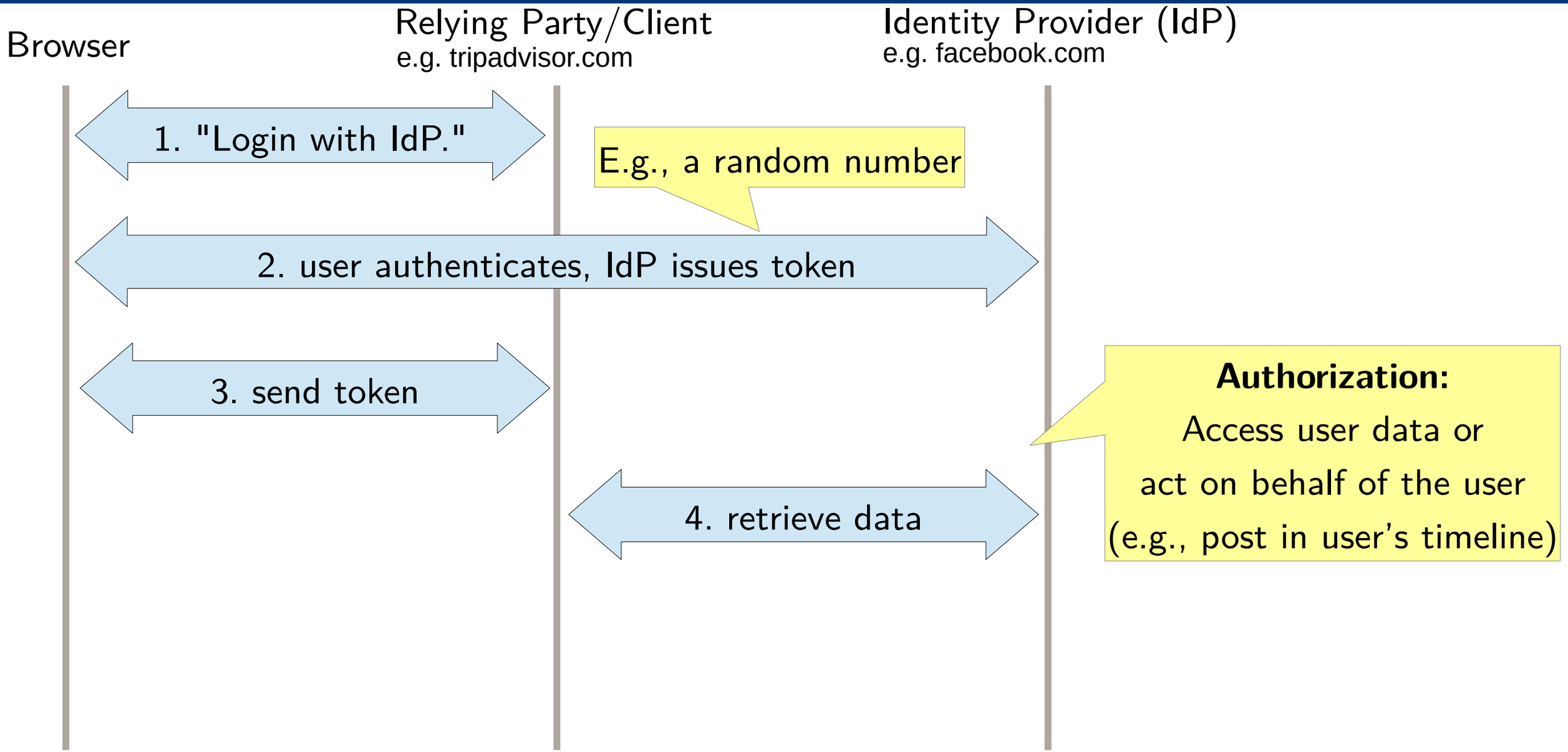

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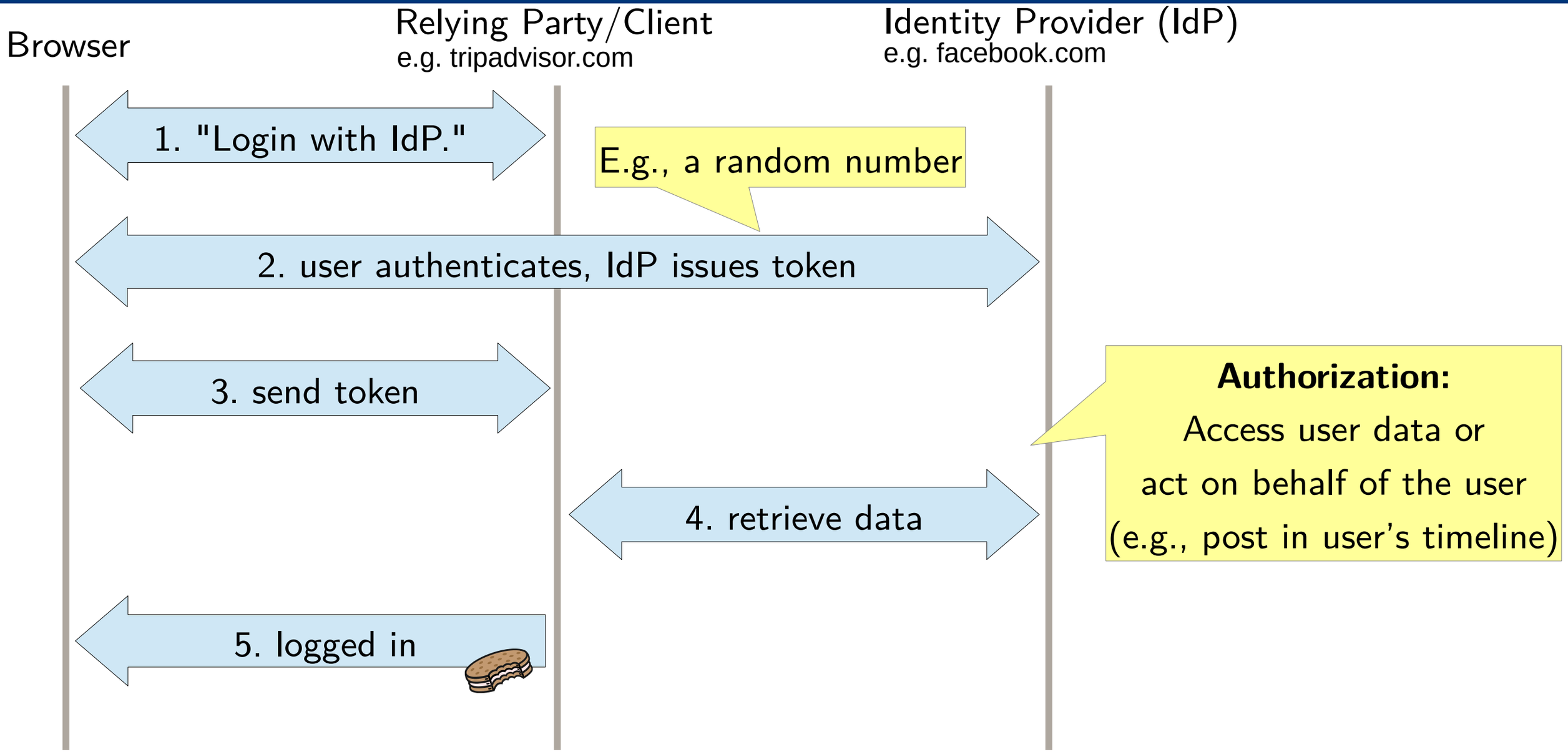

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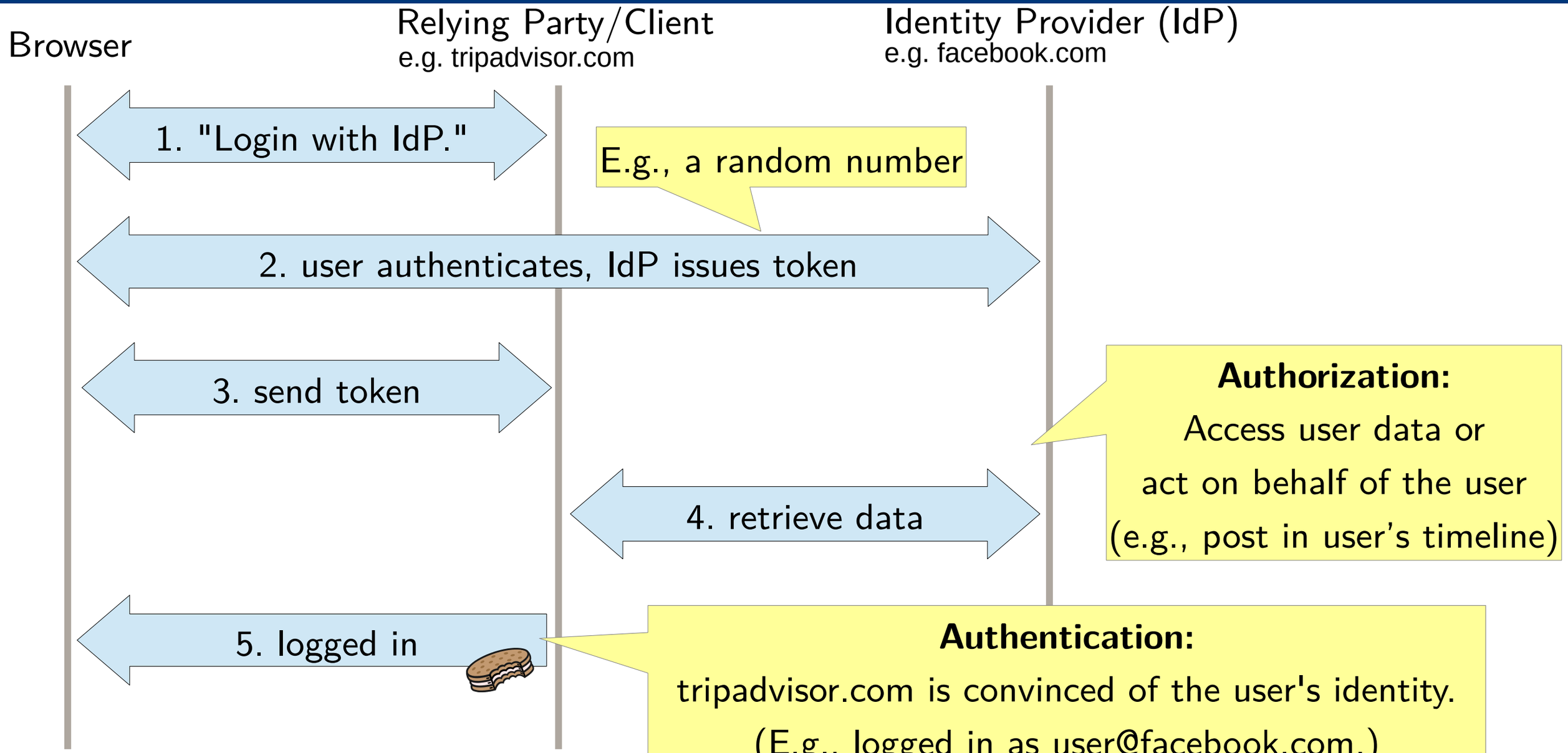

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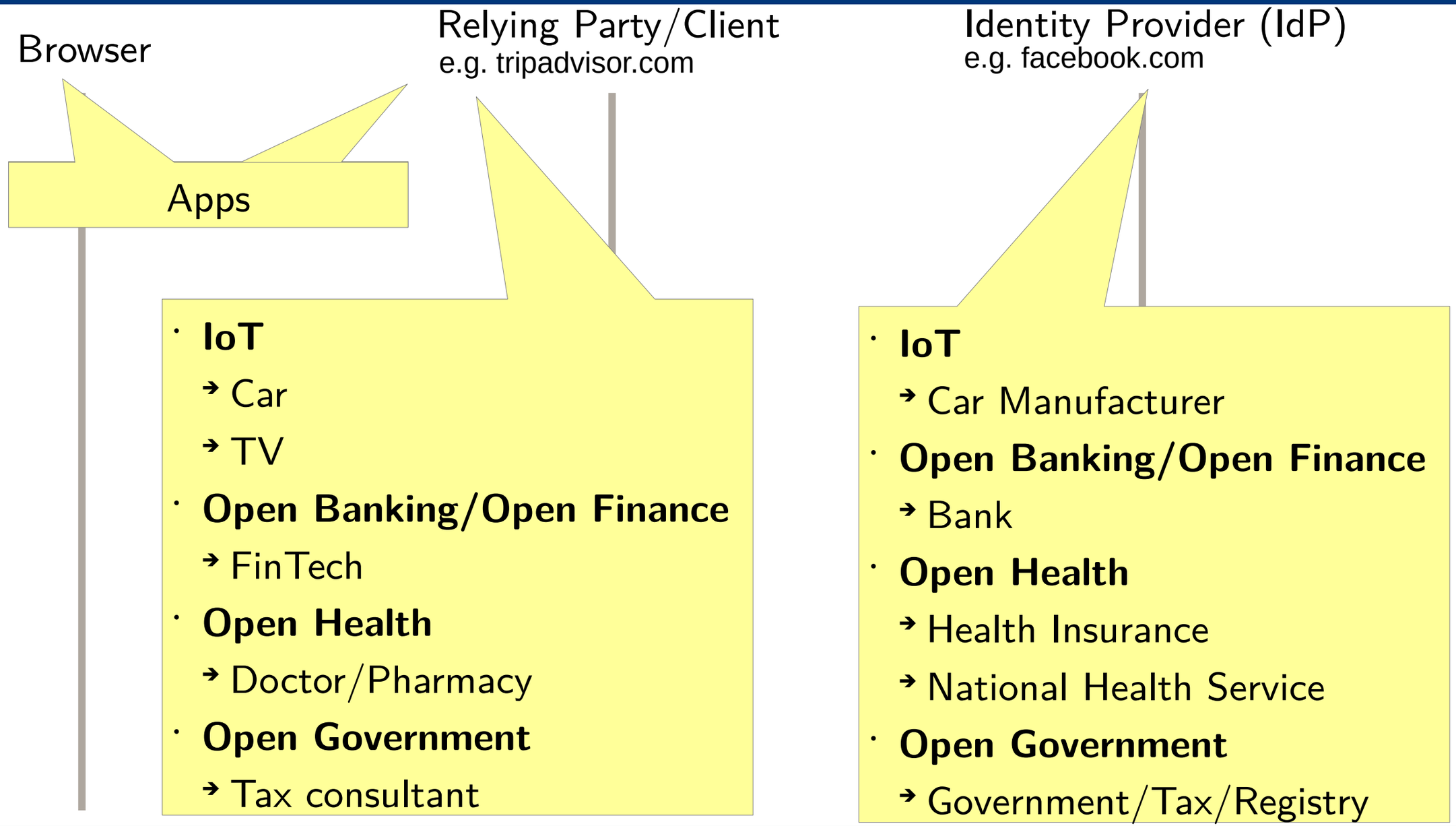
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Relying Party/Client
e.g. tripadvisor.com
Identity Provider (IdP)
e.g. facebook.com
Apps
• IoT
➔ Car
➔ TV
• Open Banking/Open Finance
➔ FinTech
• Open Health
➔ Doctor/Pharmacy
• Open Government
➔ Tax consultant
• IoT
➔ Car Manufacturer
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➔ National Health Service
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➔ Government/Tax/Registry

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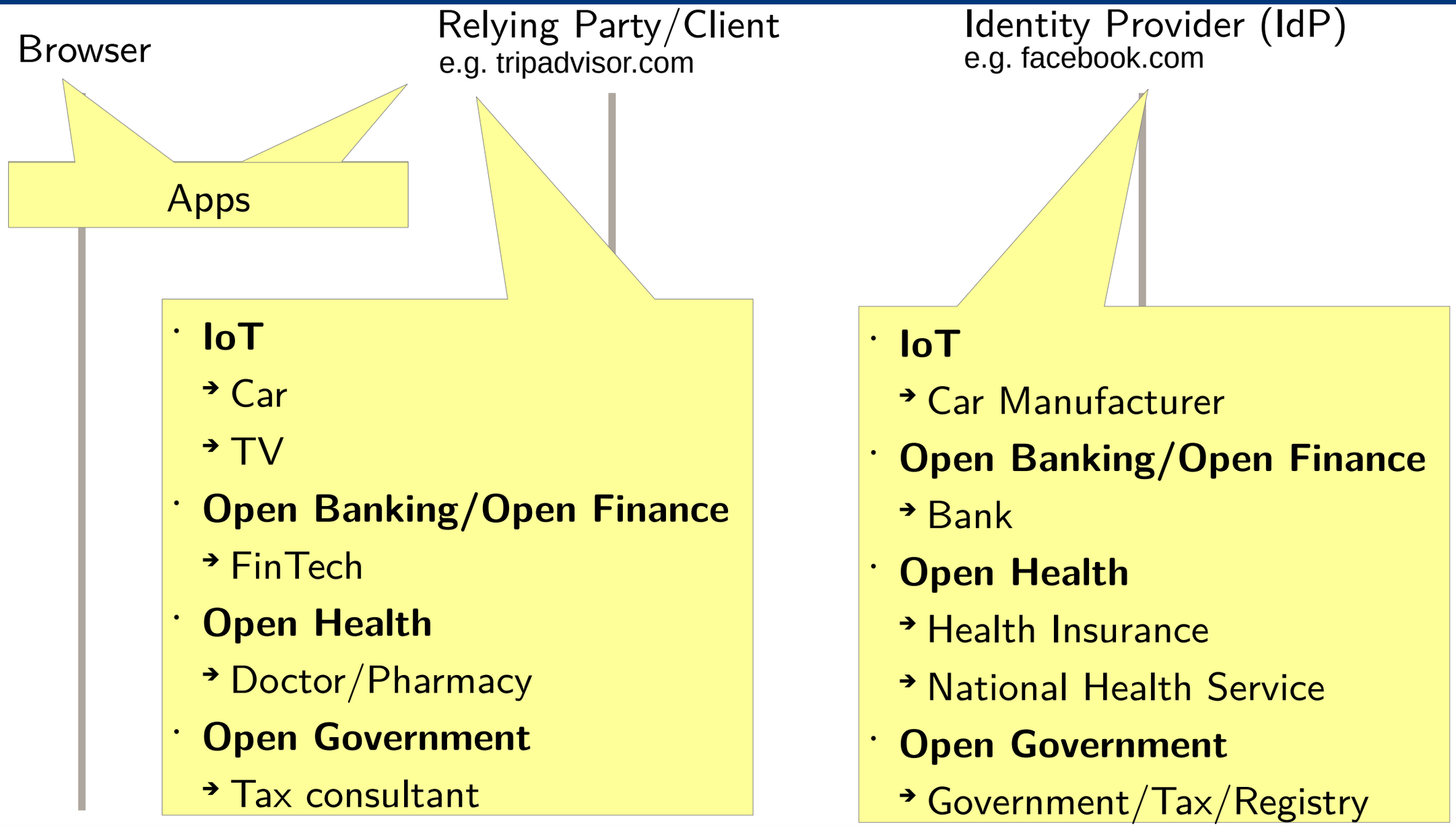
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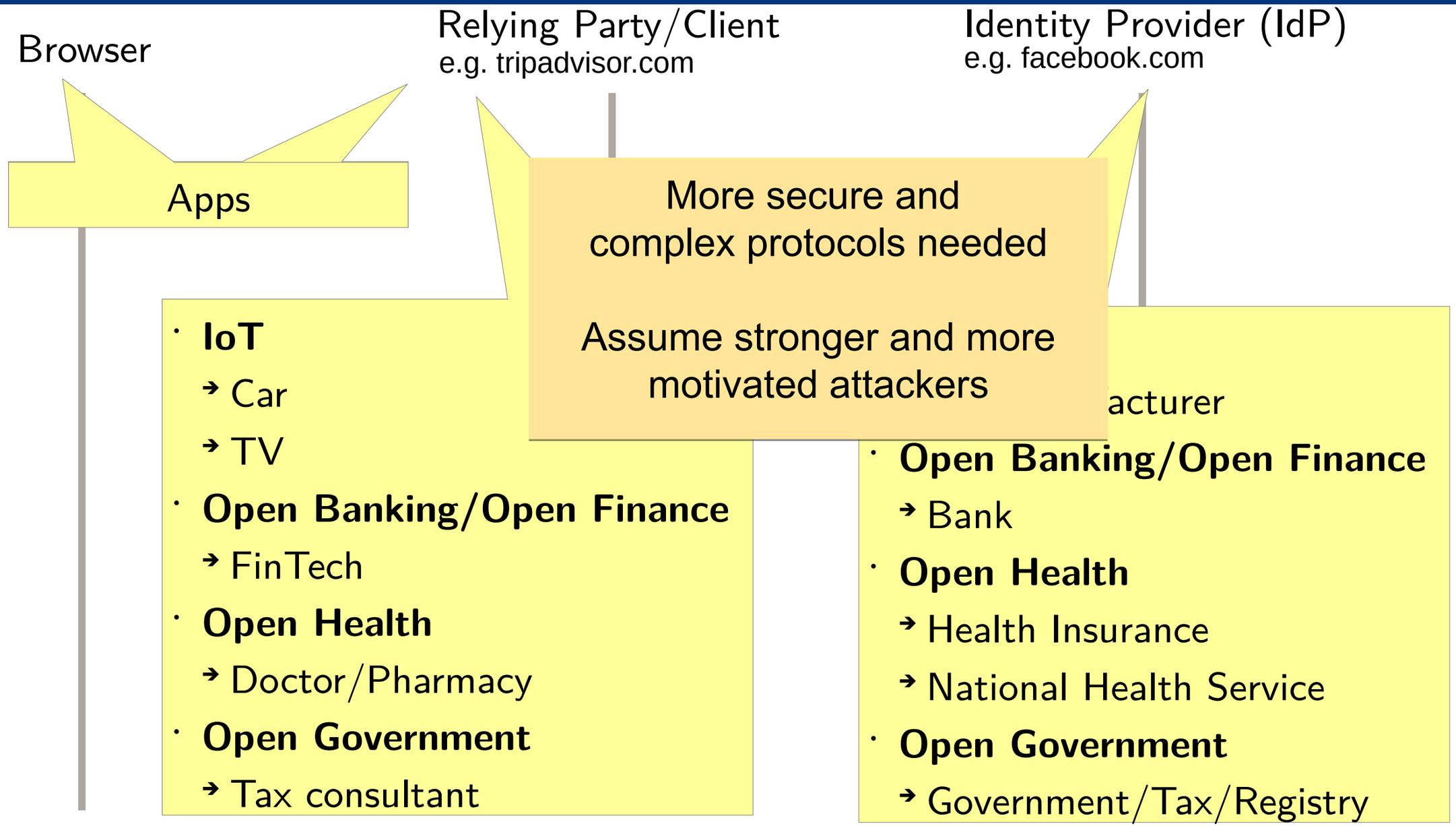
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Relying Party/Client
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Apps
More secure and
complex protocols needed
Assume stronger and more
motivated attackers
• IoT
➔ Car
➔ TV
• Open Banking/Open Finance
➔ FinTech
• Open Health
➔ Doctor/Pharmacy
• Open Government
➔ Tax consultant
acturer
• Open Banking/Open Finance
➔ Bank
• Open Health
➔ Health Insurance
➔ National Health Service
• Open Government
➔ Government/Tax/Registry

Our Goal About a Decade Ago

- A comprehensive model of the web infrastructure.
- To formally model and analyze web applications, protocols, and standards.

Result:

The Web Infrastructure Model (WIM) [S&P14]

- At that time only very limited models existed:
  - Kerschbaum as well as Akhawe et al. (Alloy models)
  - Bansal et al. (Proverif model)
- The WIM is still the by far most comprehensive model of the web infrastructure.

  [S&P14], [ESORICS15], [CCS15], [CCS16], [CSF17], [S&P19], [S&P22], [ESORICS23], [CSF24], [ACM TOPS24]

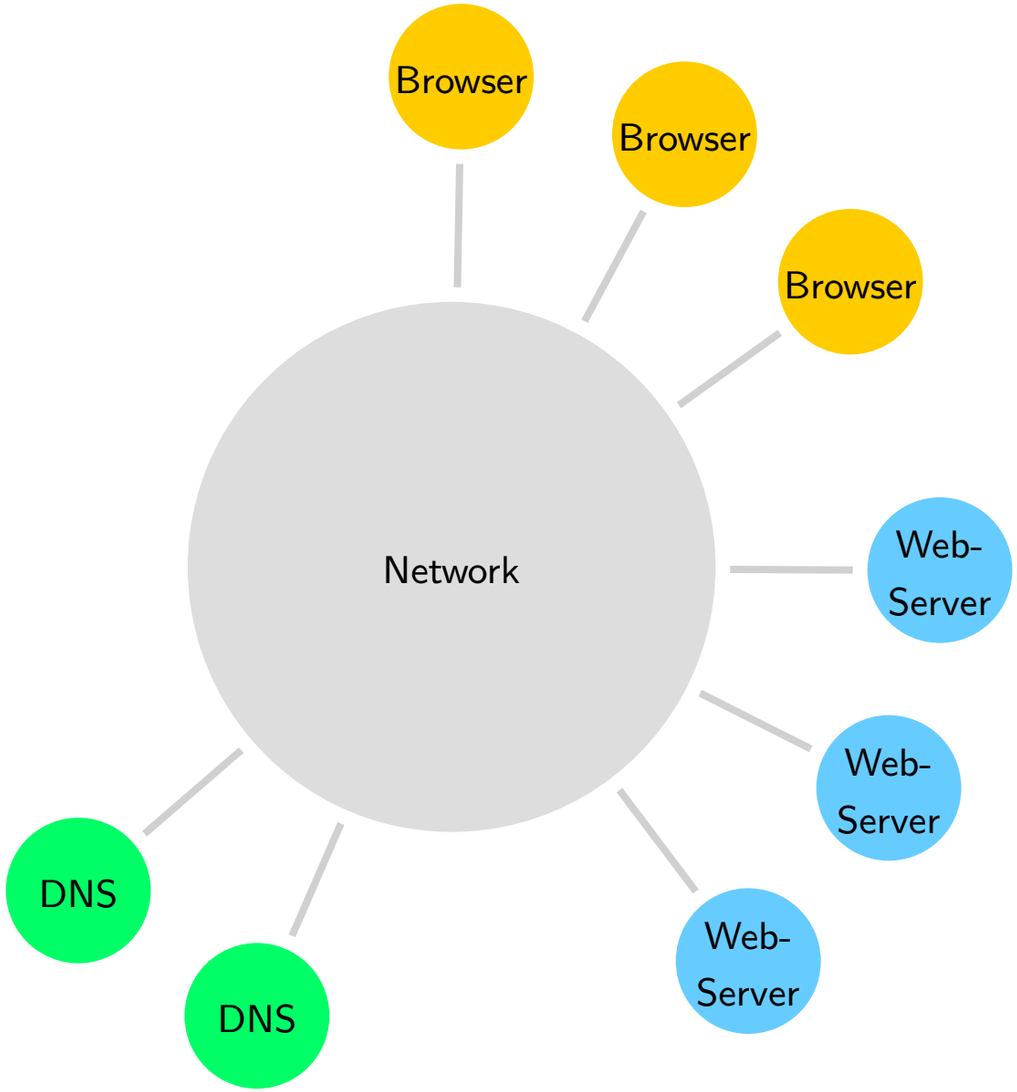
The Web Infrastructure Model (WIM)

Sources

Specifications for the web are spread across many sources with mutual dependencies:

- Standards and RFCs
  - HTTP/1.1, HTTP/2, HTTP/3 Standards
  - W3C HTML5
  - W3C Web Storage
  - WHATWG Fetch
  - W3C Cross-Origin Resource Sharing
  - RFCs (6265, 6797, 6454, 2616, ...)
- Browser implementations
  - Google Chrome
  - Mozilla Firefox
  - ...

WIM: Network Model and Attackers



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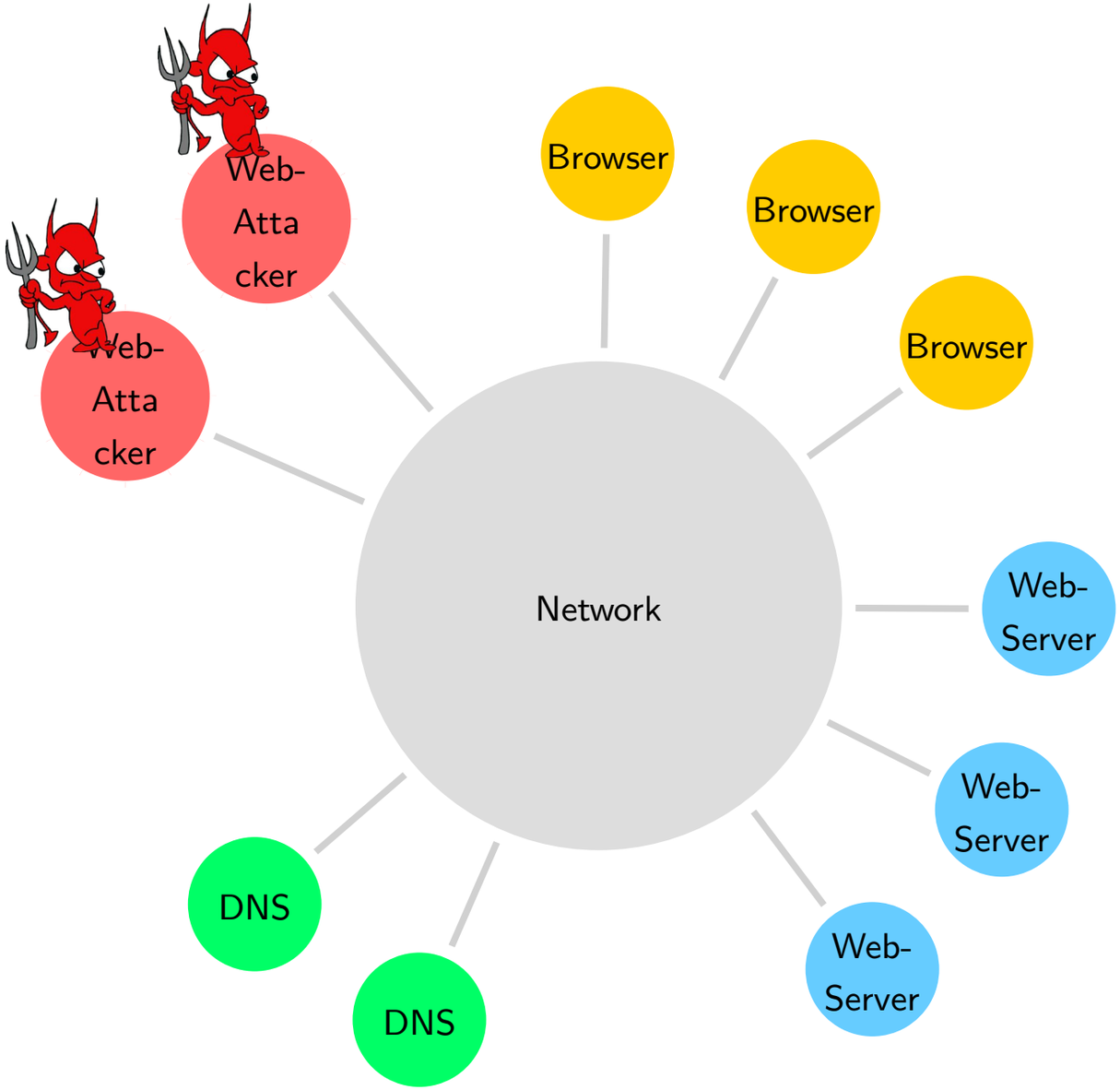

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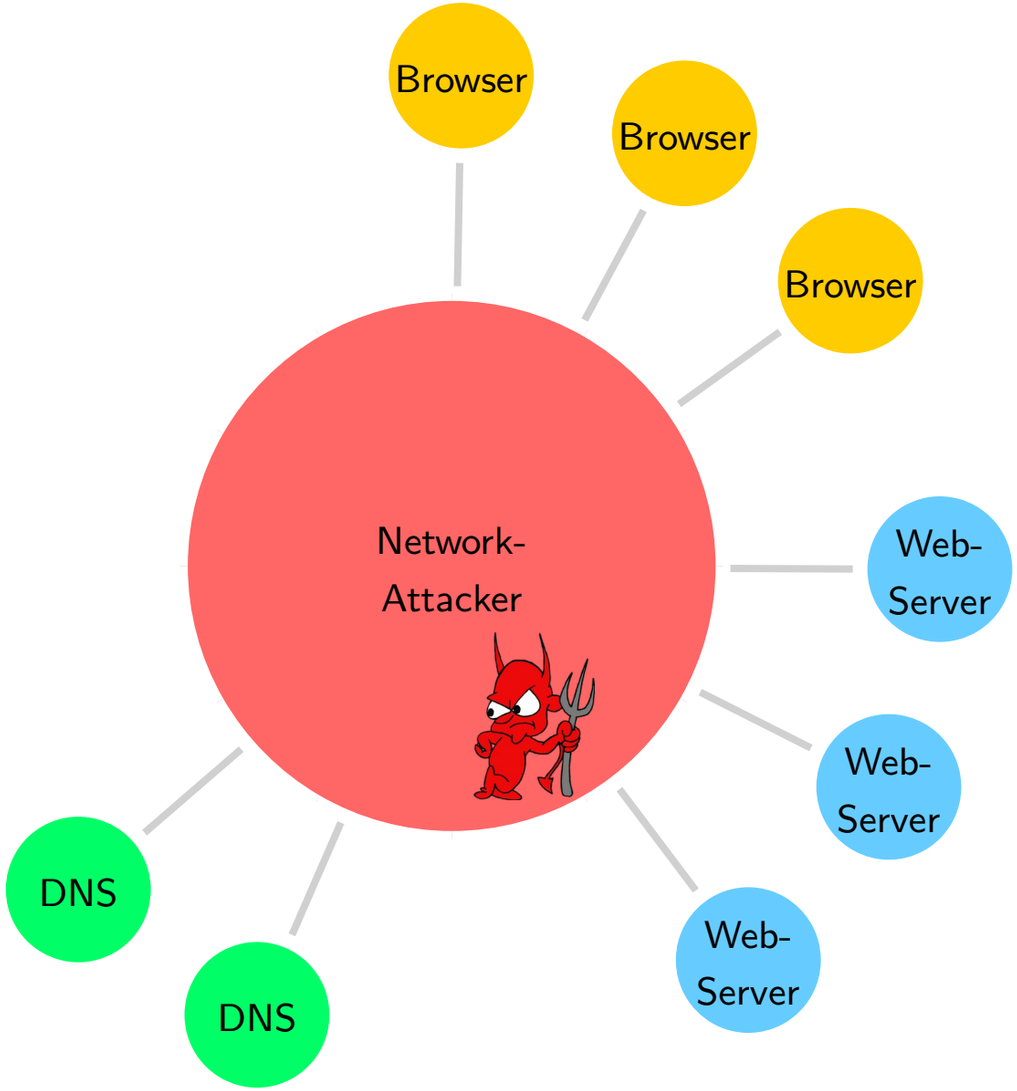

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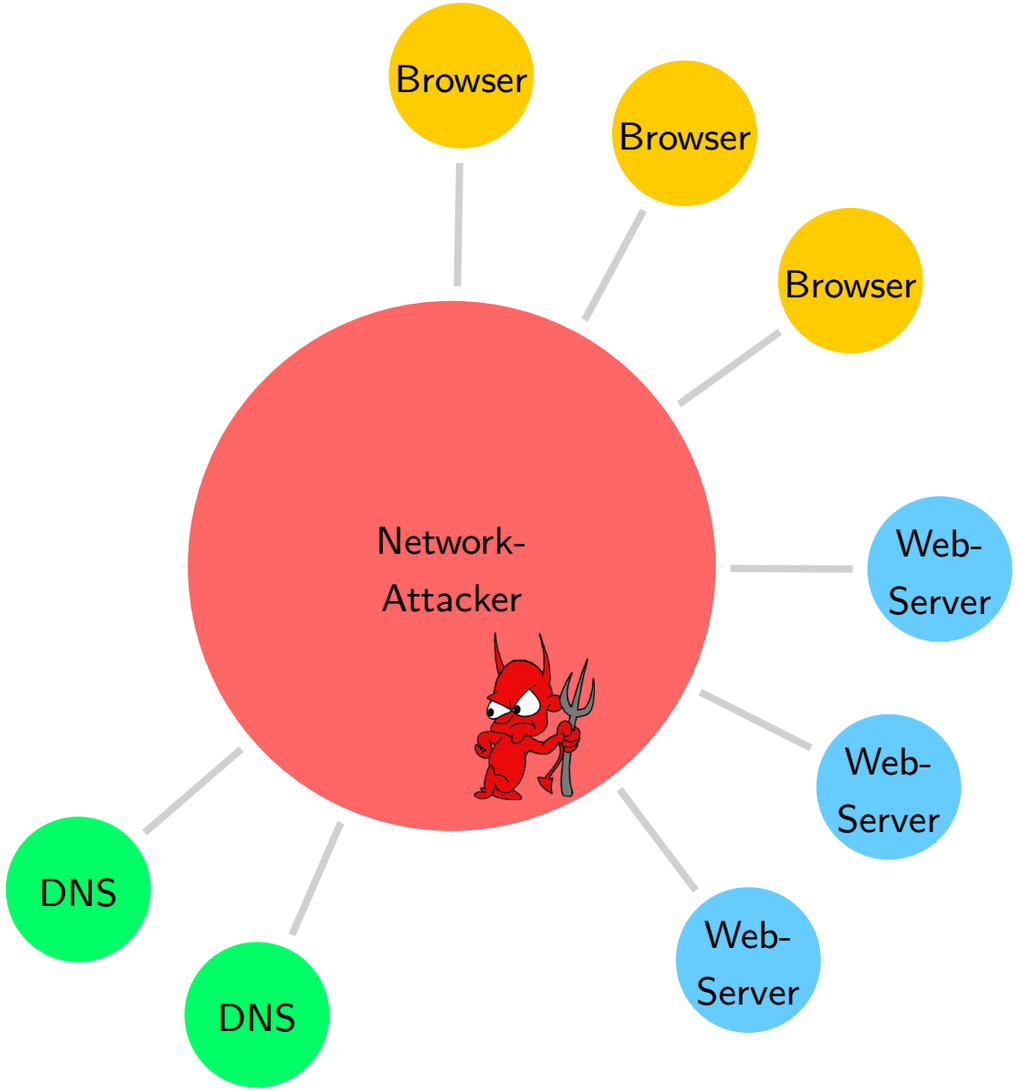

Dolev-Yao-Attacker

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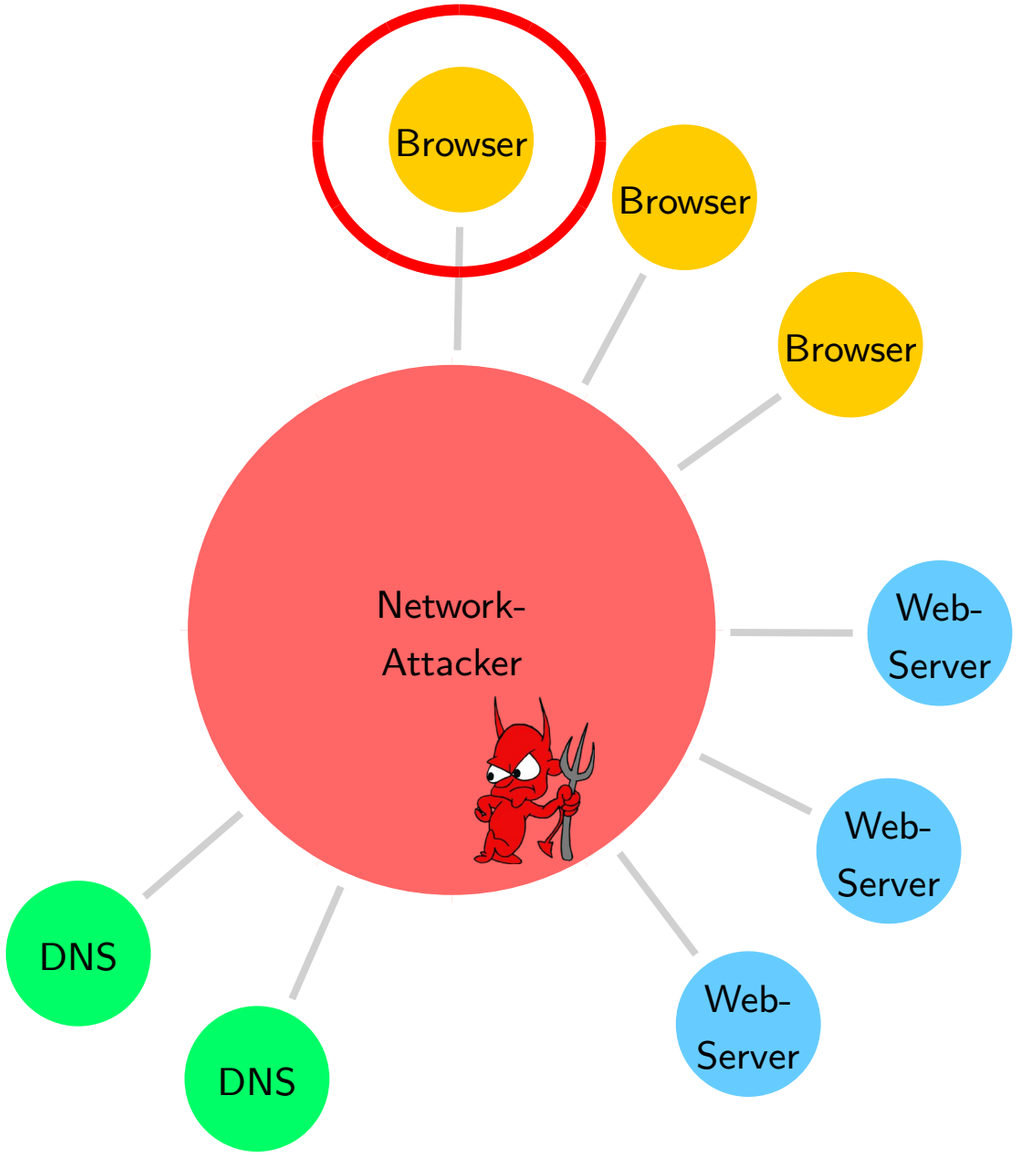

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Including ...

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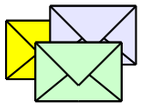

Including ...

- DNS, HTTP, HTTPS

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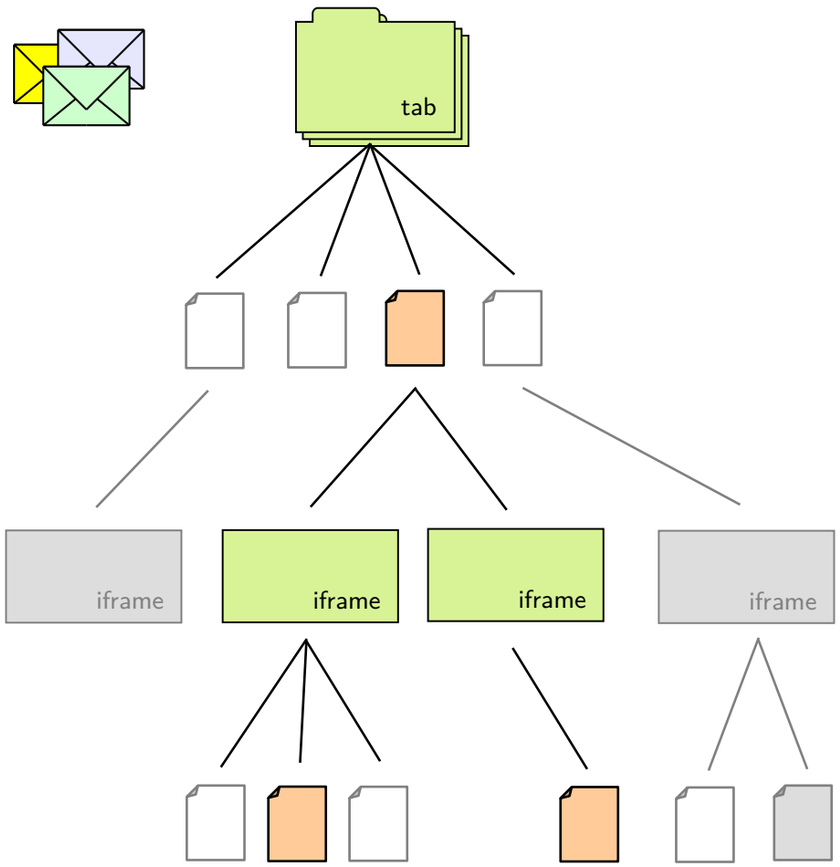
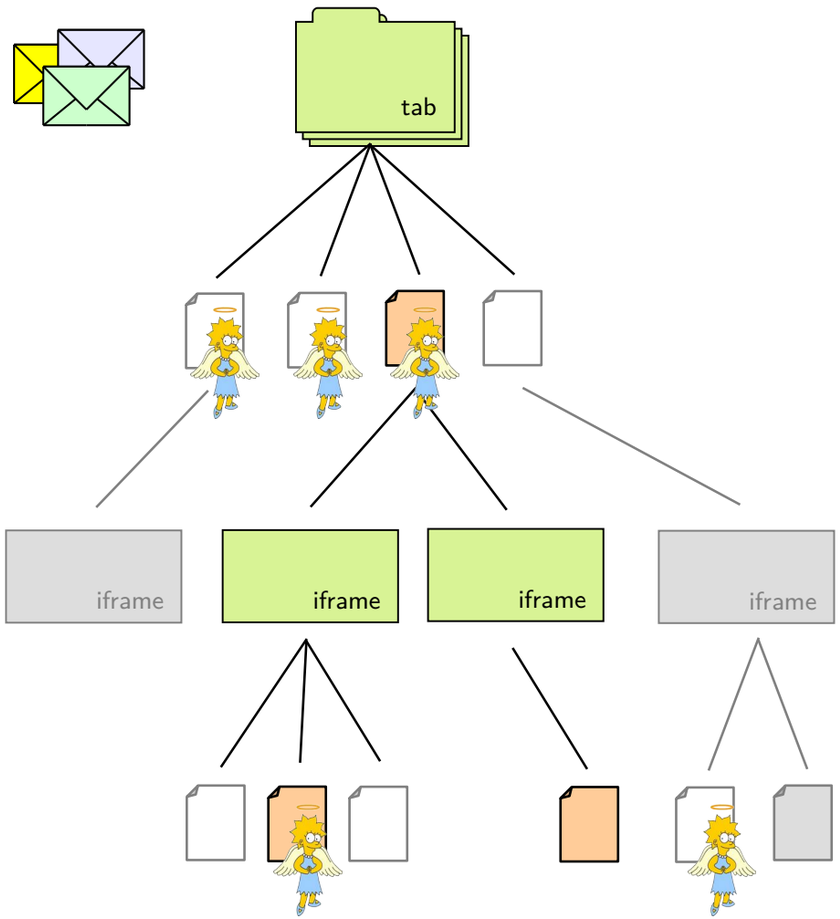


Including ...

- DNS, HTTP, HTTPS 
- window & document structure

WIM: Web Browser Model



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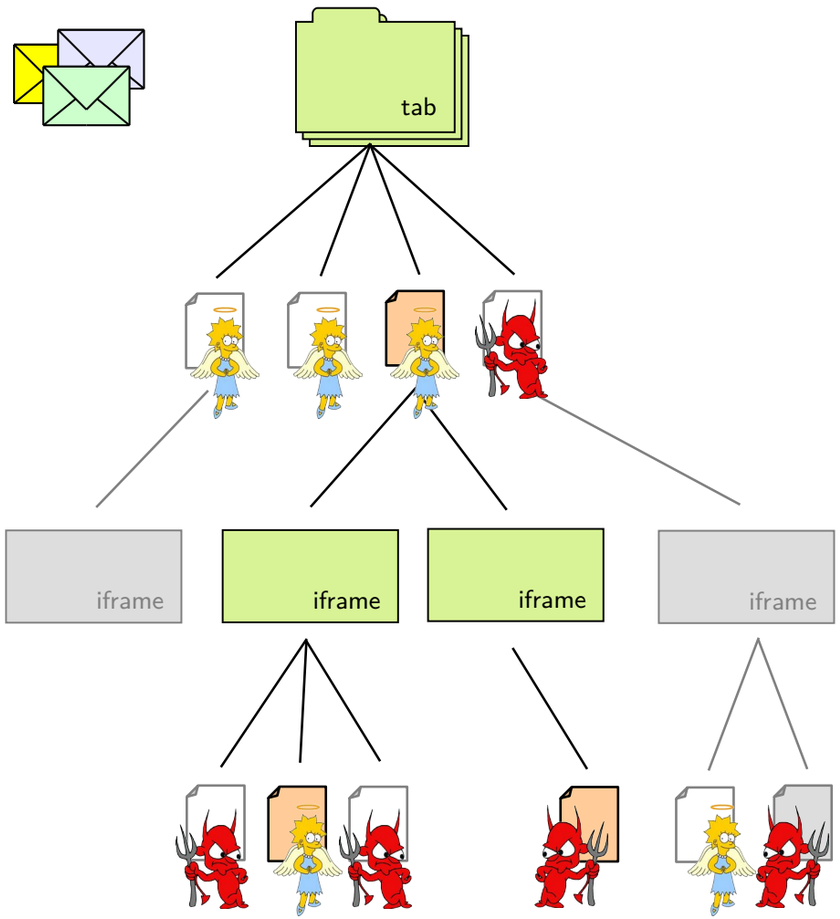

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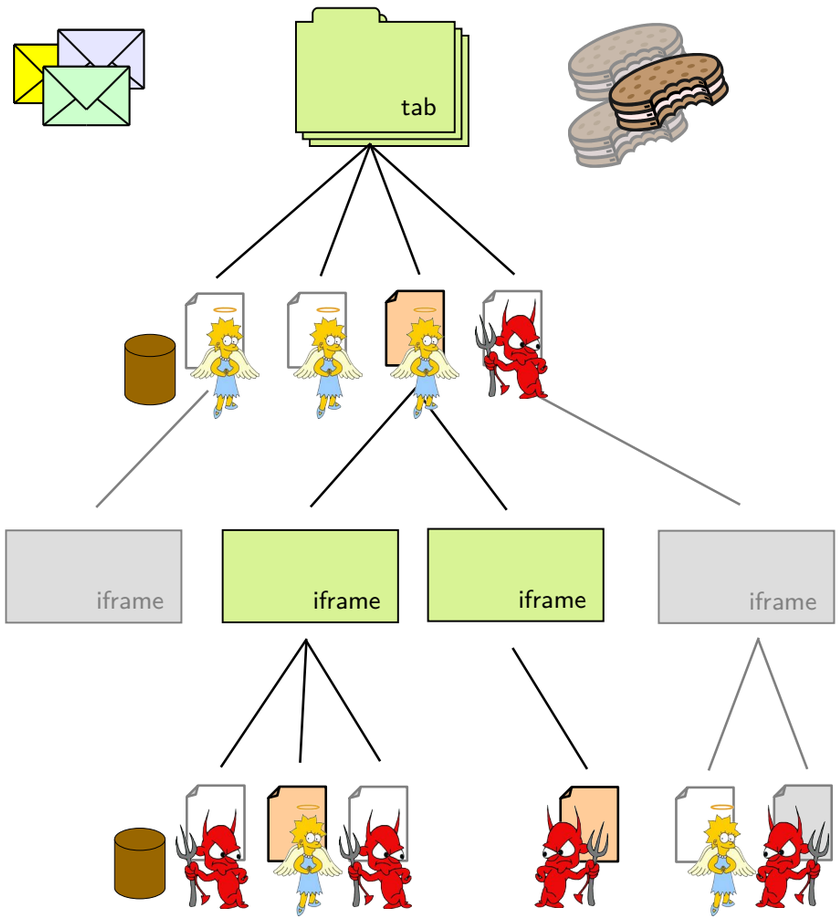

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- web storage & cookies

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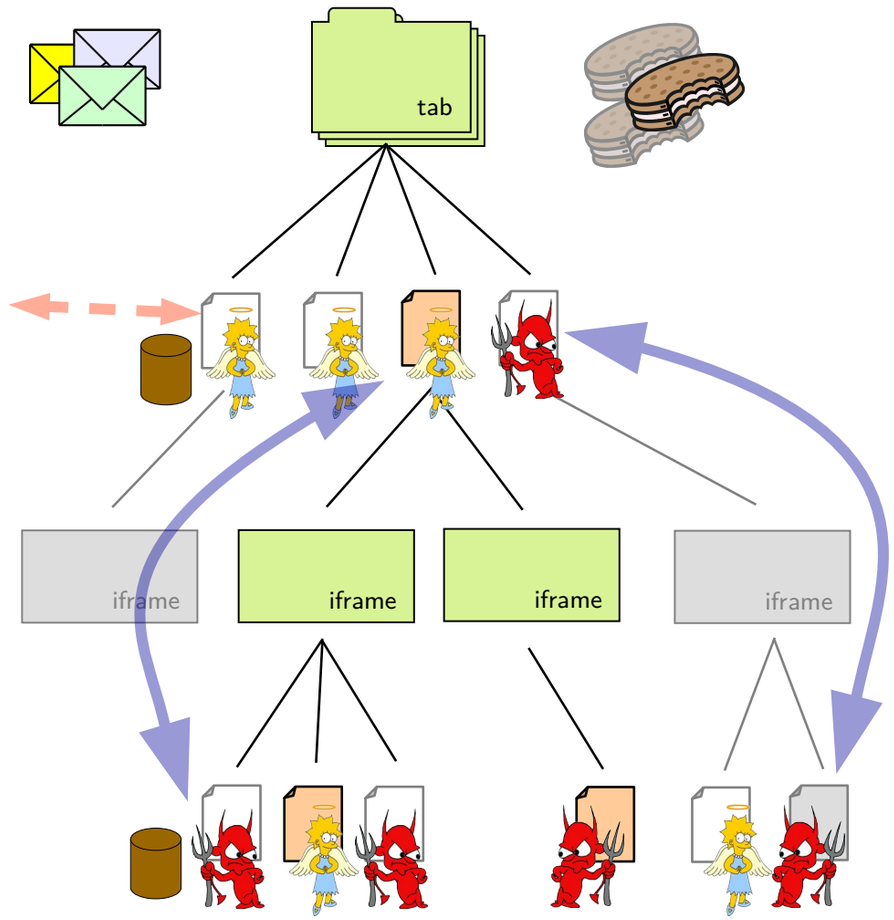

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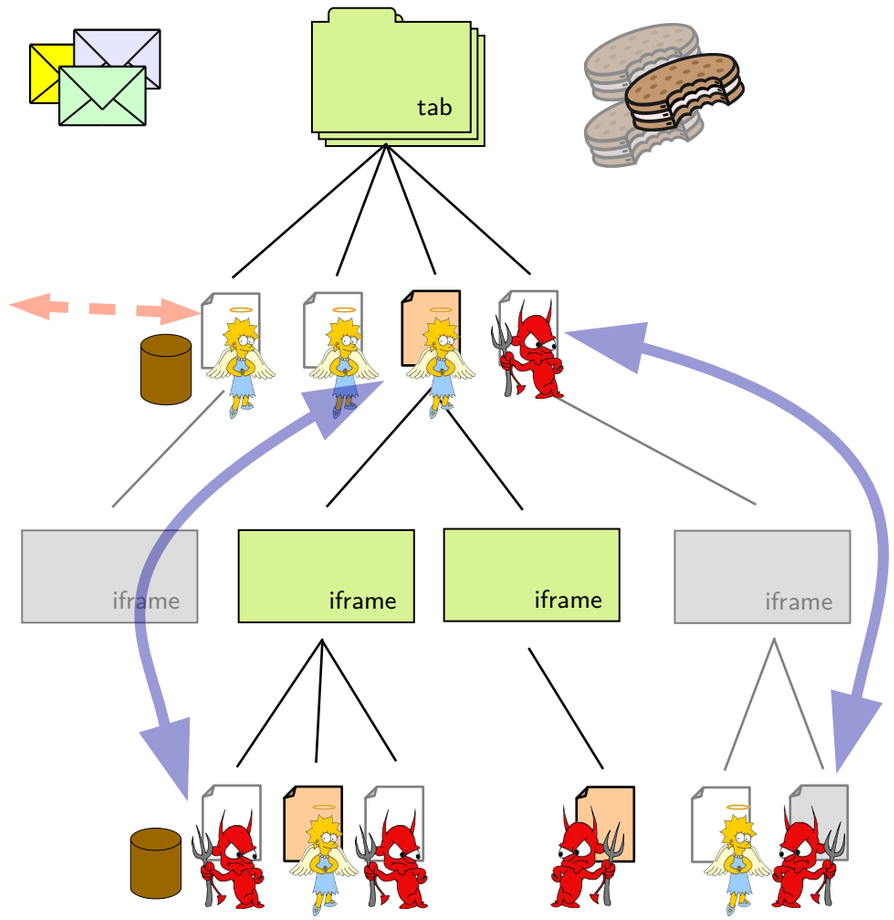

Including ...

- DNS, HTTP, HTTPS
- window & document structure
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- web storage & cookies
- web messaging & XHR
- message headers
- redirections
- security policies
- dynamic corruption
- WebRTC
- ...






WIM: Web Browser Model – Excerpt

Algorithm 8 Web Browser Model: Process an HTTP response.

```
1:  function PROCESSRESPONSE(response, reference, request, requestUrl, key, f, s')
2:      if Set-Cookie ∈ response.headers then
3:          for each c ∈⟨⟩ response.headers[Set-Cookie], c ∈ Cookies do
4:              let s'.cookies[request.host]
                ↪  := AddCookie(s'.cookies[request.host], c)
5:      if Strict-Transport-Security ∈ response.headers ∧ requestUrl.protocol ≡ S then
6:          let s'.sts := s'.sts +⟨⟩ request.host
7:      if Referer ∈ request.headers then
8:          let referrer := request.headers[Referer]
9:      else
10:         let referrer := ⊥
11:     if Location ∈ response.headers ∧ response.status ∈ {303, 307} then
12:         let url := response.headers[Location]
13:         if url.fragment ≡ ⊥ then
14:             let url.fragment := requestUrl.fragment
15:         let method' := request.method
16:         let body' := request.body
17:         if Origin ∈ request.headers then
18:             let origin := ⟨request.headers[Origin], ⟨request.host, url.protocol⟩⟩
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6:          let s'.sts := s'.sts +⟨⟩ request.host
7:      if Referer ∈ request.headers then
8:          let referrer := request.headers[Referer]
9:      else
10:         let referrer := ⊥
11:     if Location ∈ response.headers ∧ response.status ∈ {303, 307} then
12:         let url := response.headers[Location]
13:         if url.fragment ≡ ⊥ then
14:             let url.fragment := requestUrl.fragment
15:         let method' := request.method
16:         let body' := request.body
17:         if Origin ∈ request.headers then
18:             let origin := ⟨request.headers[Origin], ⟨request.host, url.protocol⟩⟩
19:         else
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22:             let method' := GET
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```

WIM: Web Browser Model – Excerpt

Algorithm 8 Web Browser Model: Process an HTTP response.

```
1:  function PROCESSRESPONSE(response, reference, request, requestUrl, key, f, s')
2:      if Set-Cookie ∈ response.headers then
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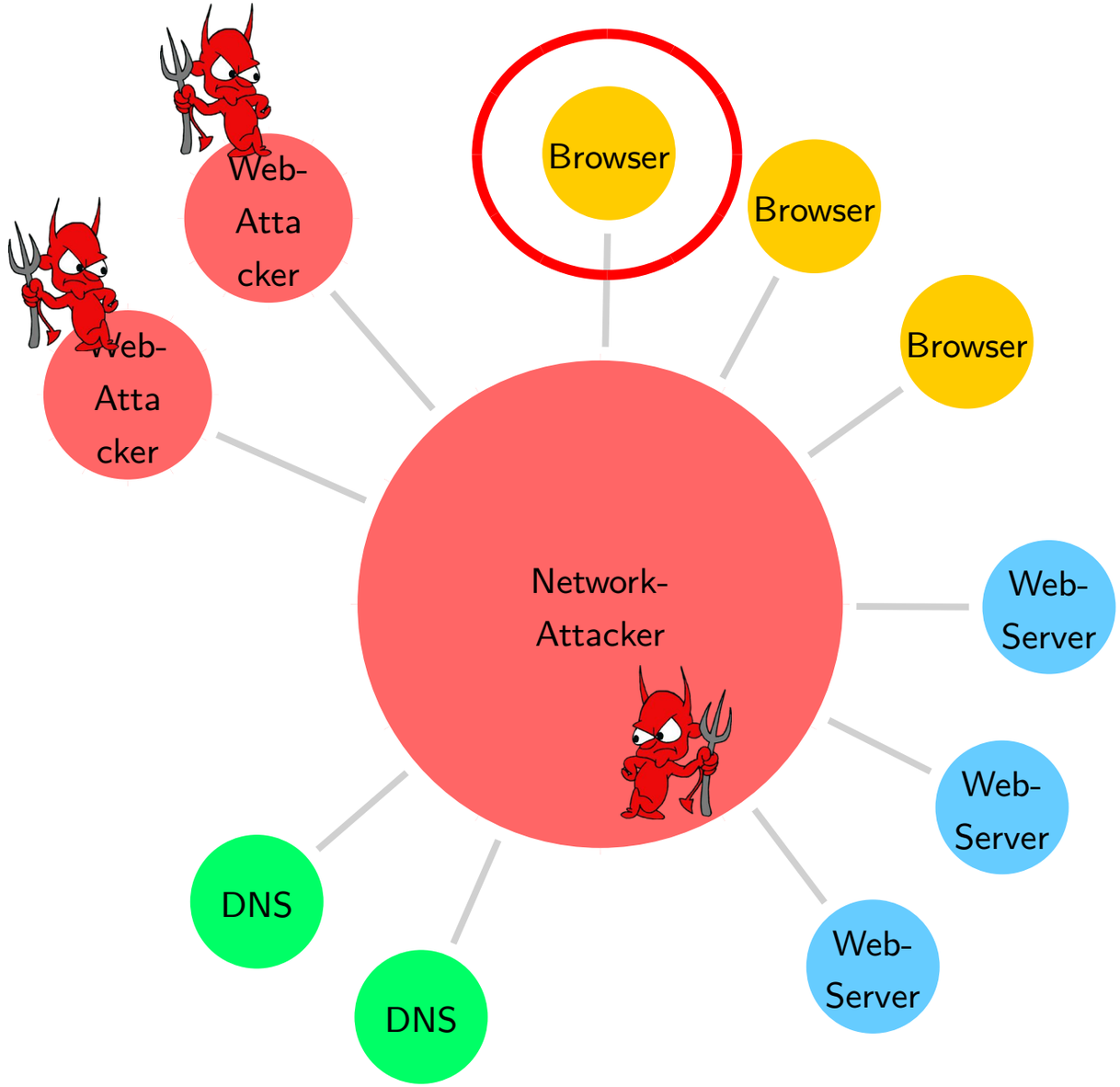
WIM: Network Model and Attackers



Dolev-Yao-Attacker

Limitations

- No language details
- No user interface details (e.g., no clickjacking attacks)
- No byte-level attacks (e.g., buffer overflows)
- Abstract view on cryptography and TLS

Model can in principle be extended to capture these aspects as well.

Trade-off: comprehensiveness vs. simplicity

Limitations

- No language details

main focus in other work,
e.g., Calzavara, Foccardi et al.

- No user interface details (e.g., no clickjacking attacks)
- No byte-level attacks (e.g., buffer overflows)
- Abstract view on cryptography and TLS

Model can in principle be extended to capture these aspects as well.

Trade-off: comprehensiveness vs. simplicity

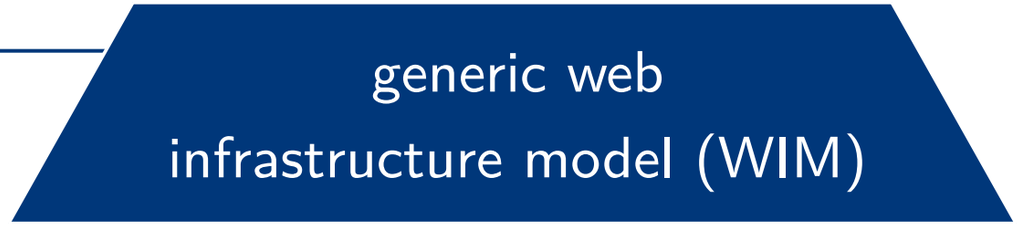
How to use the WIM?

How to use the WIM?

Foundation:

Formal description

of the web



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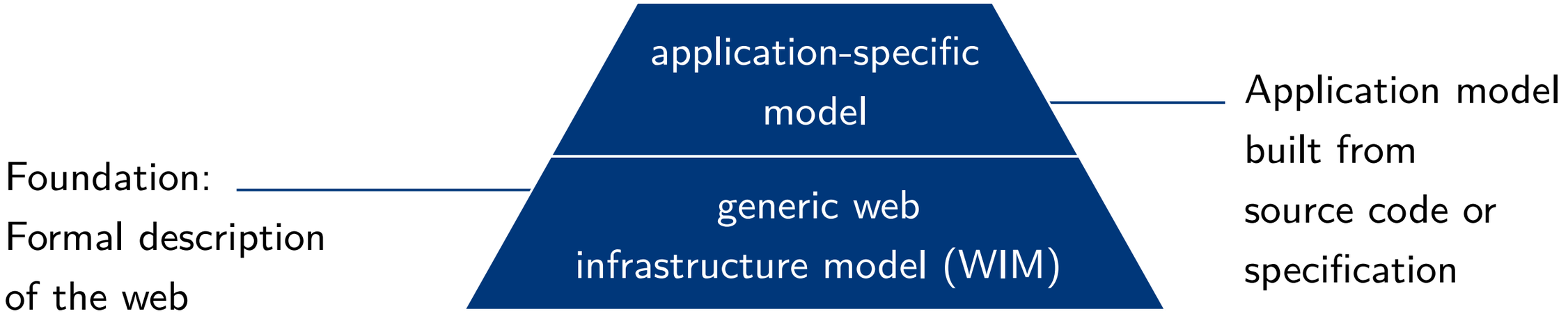

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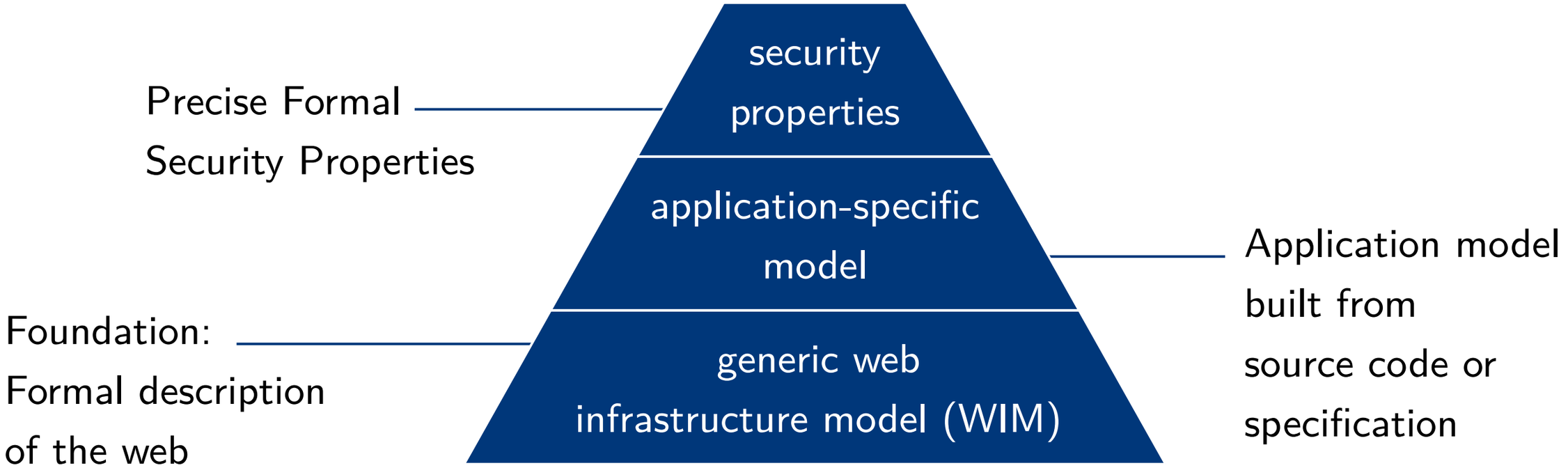

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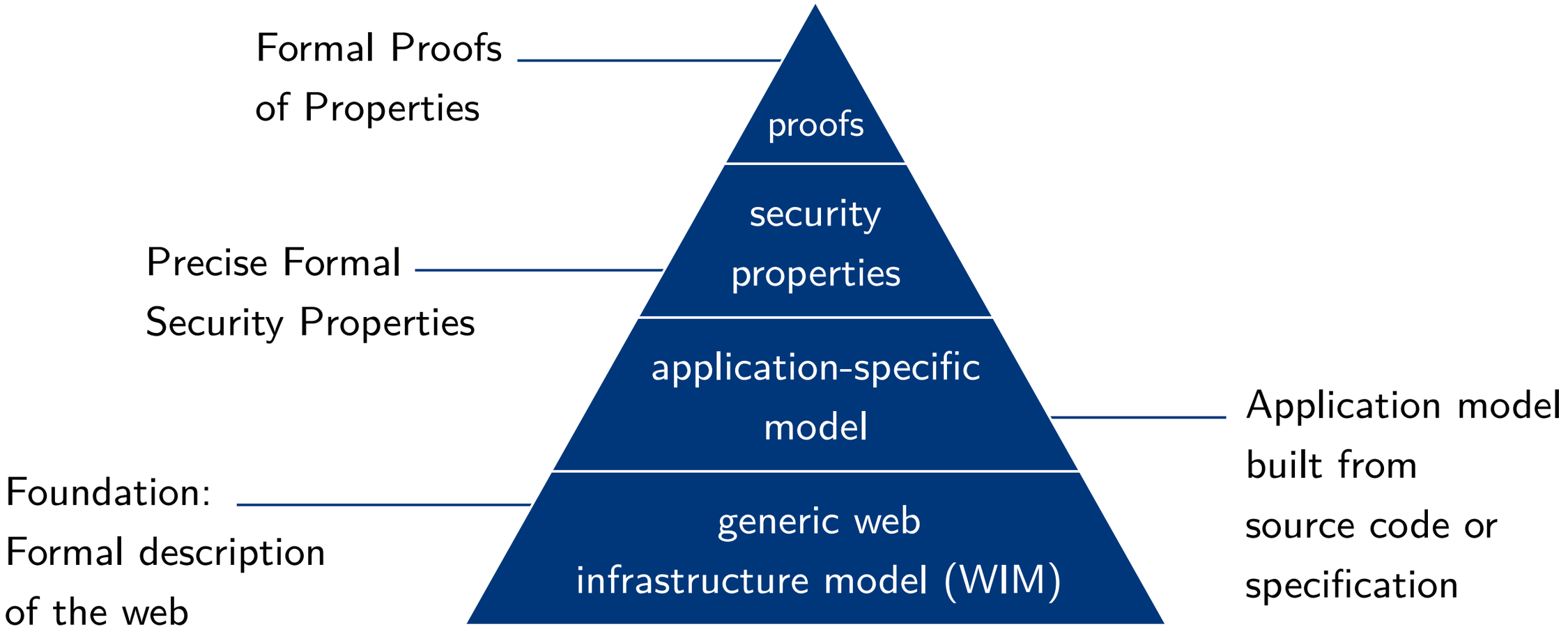

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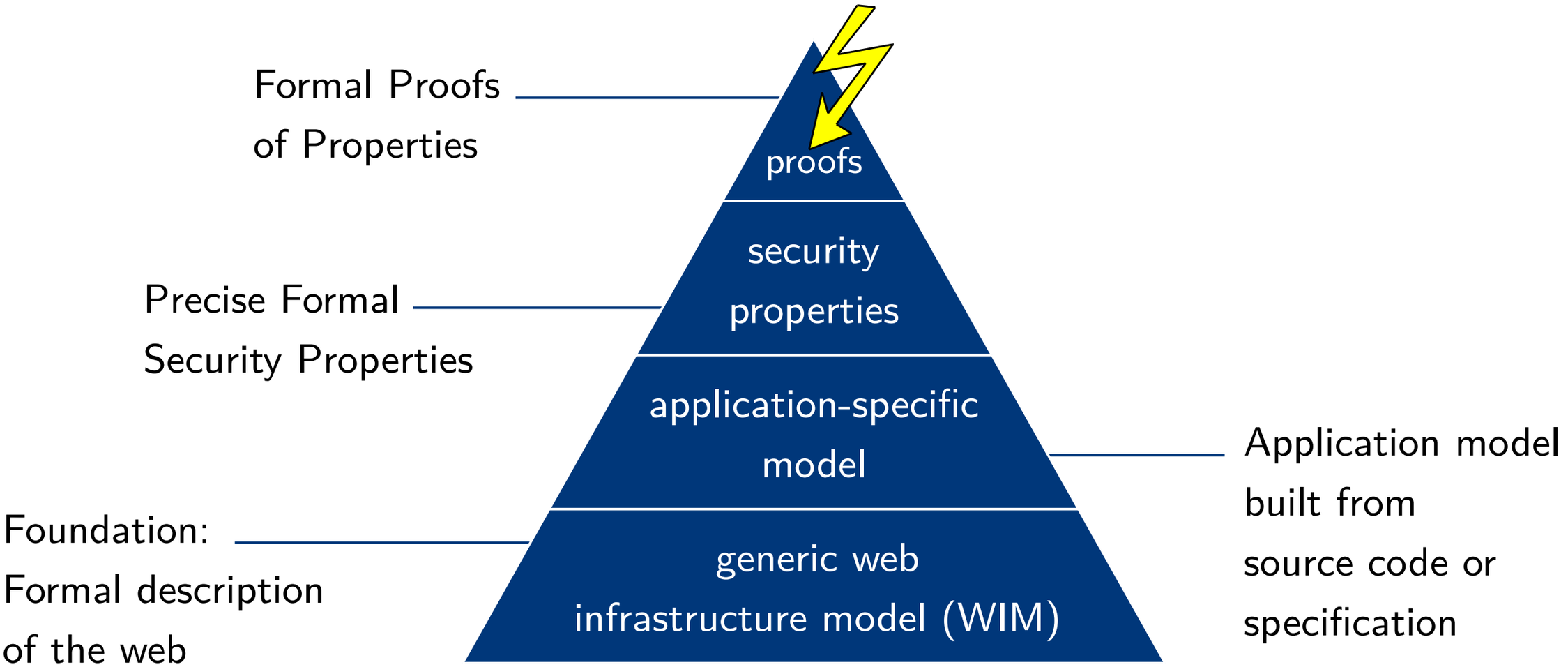

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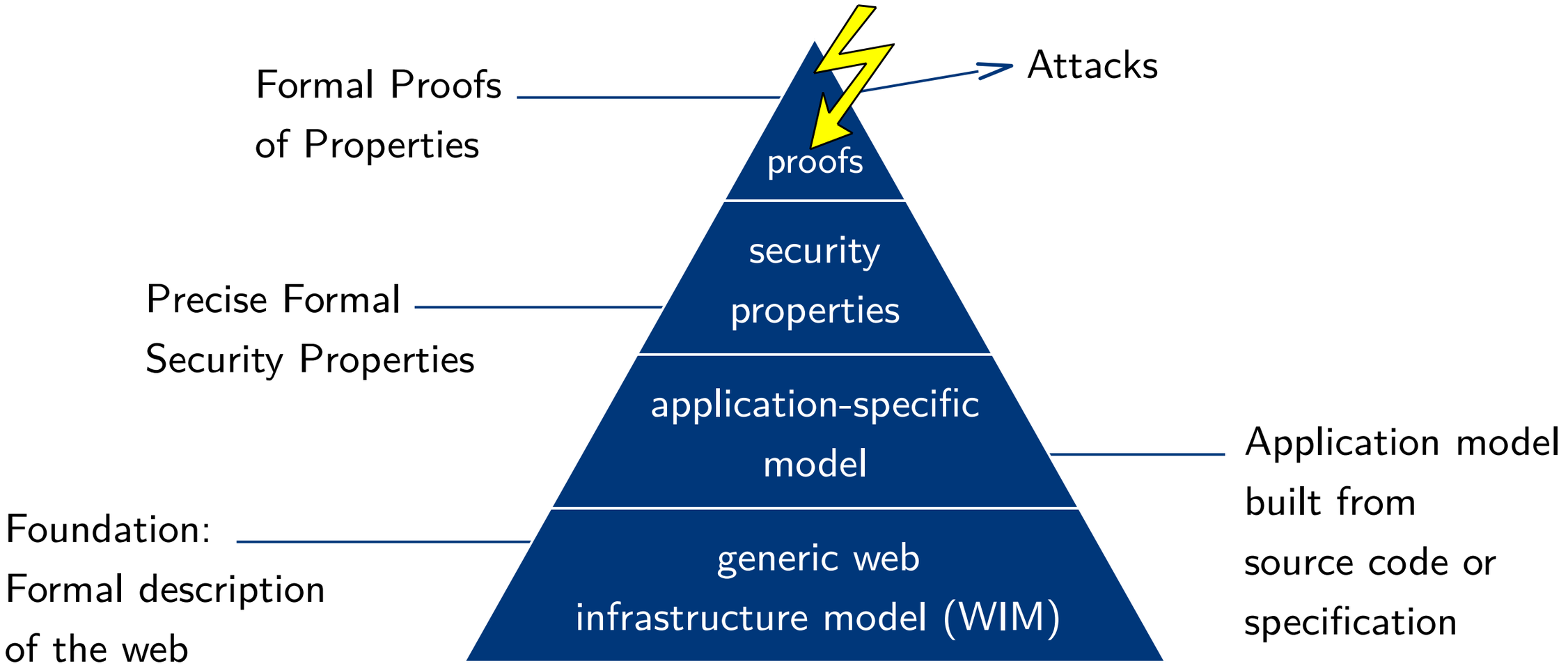

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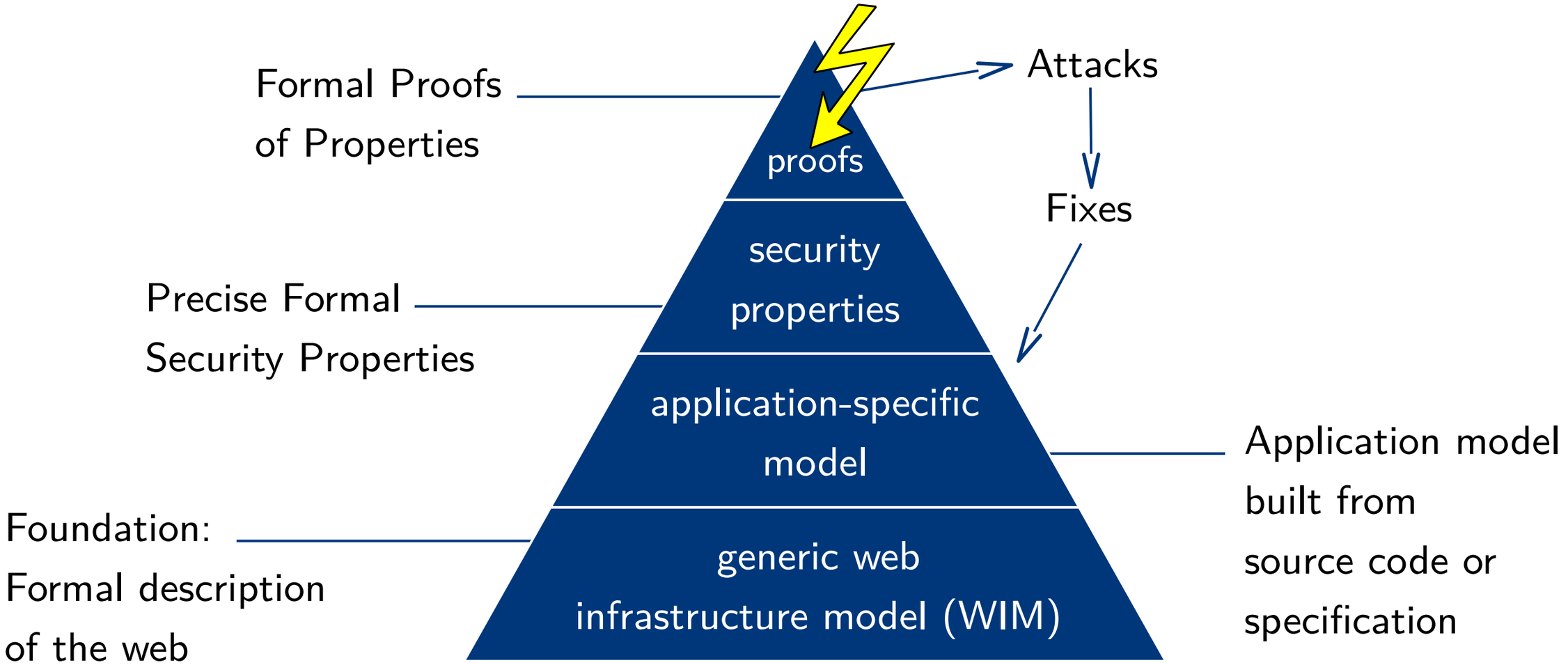
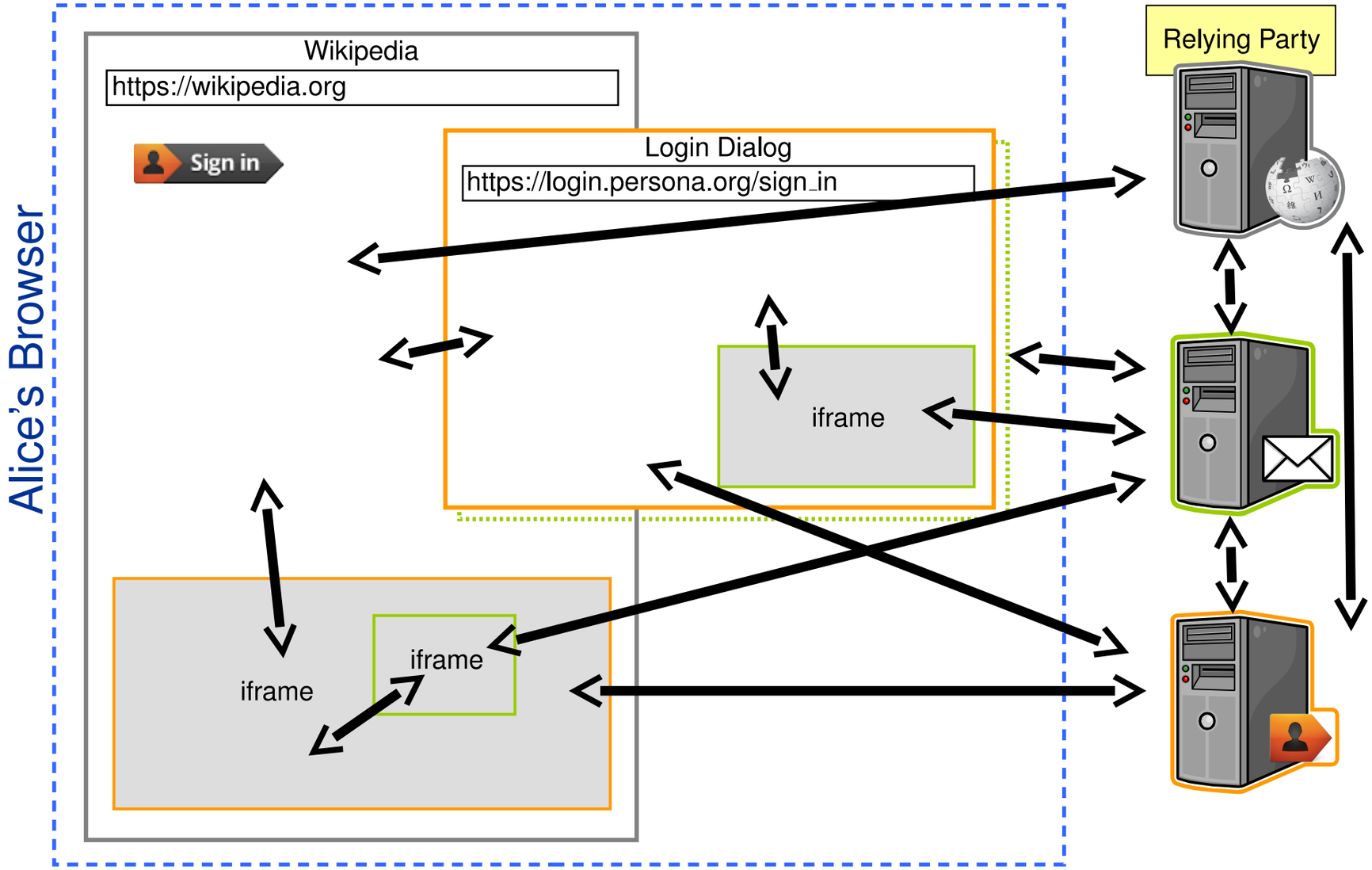


Case Studies

Case Studies – Our Very First WIM Case Study [S&P14]

Example: Mozilla BrowserID

- BrowserID used a complex system of iframes and windows to transfer data between the RP, the IdP, and Mozilla's own servers.
- This was supposed to hide RP's identity from the IdP (but not from Mozilla).



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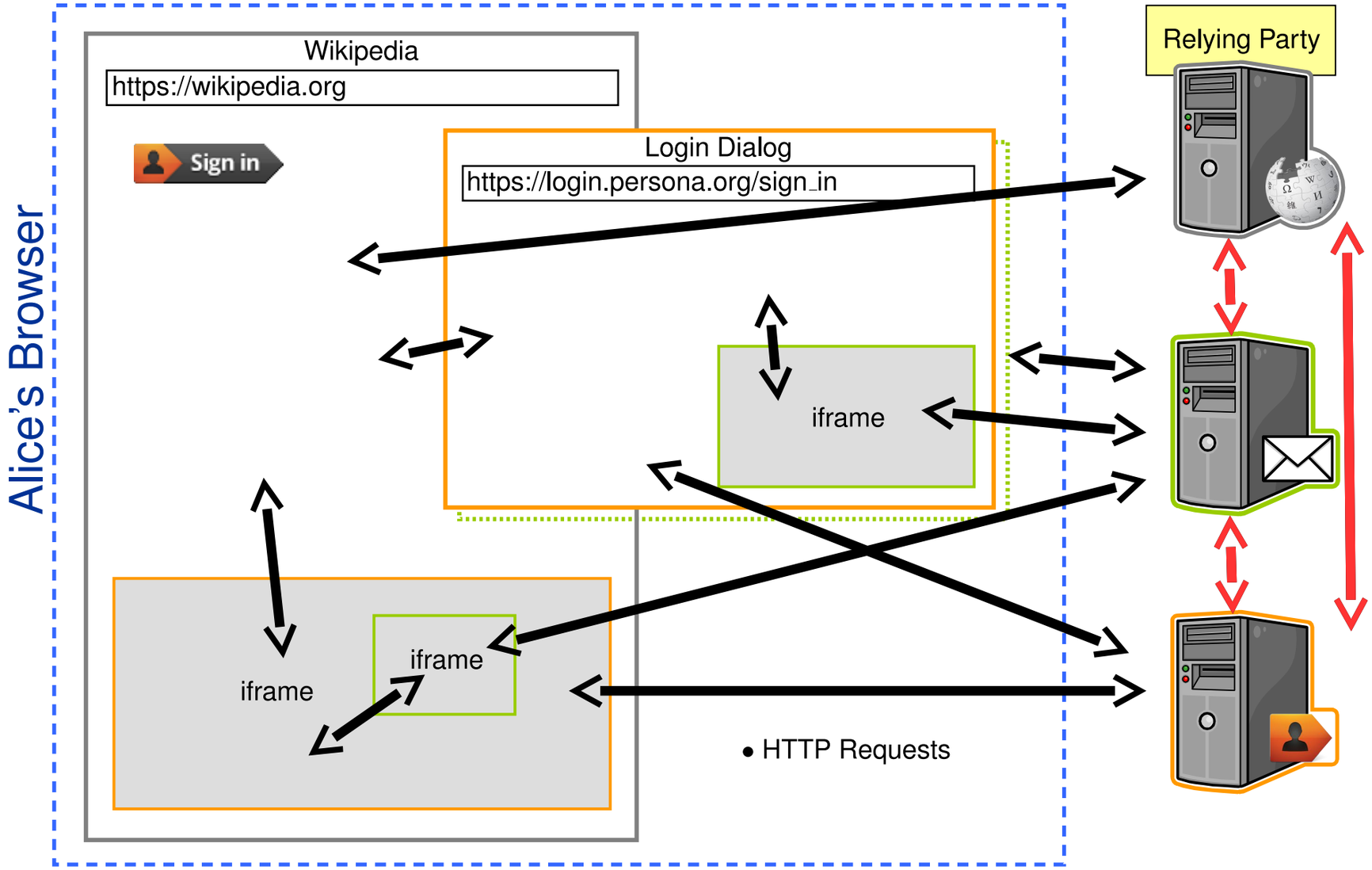

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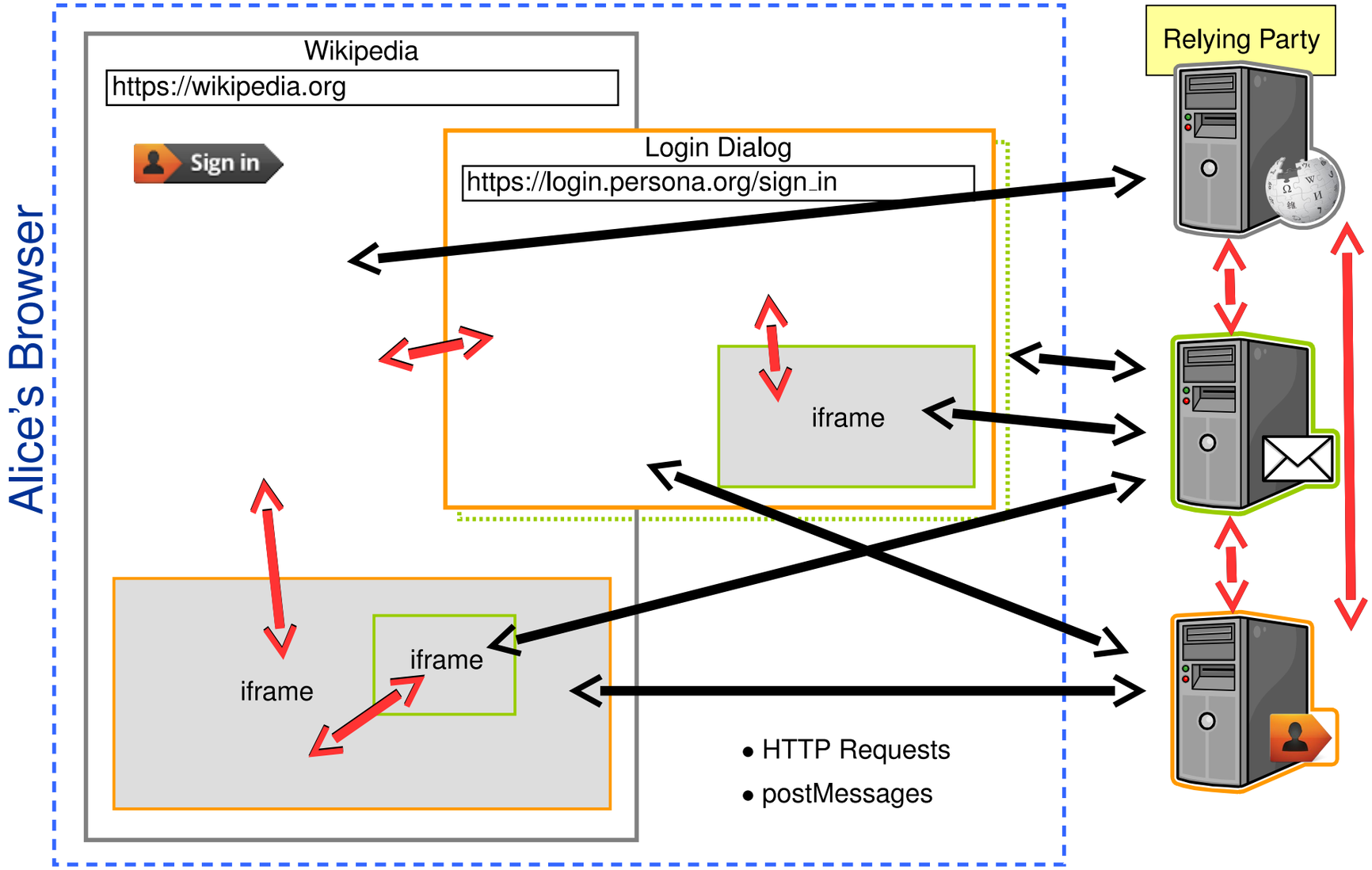

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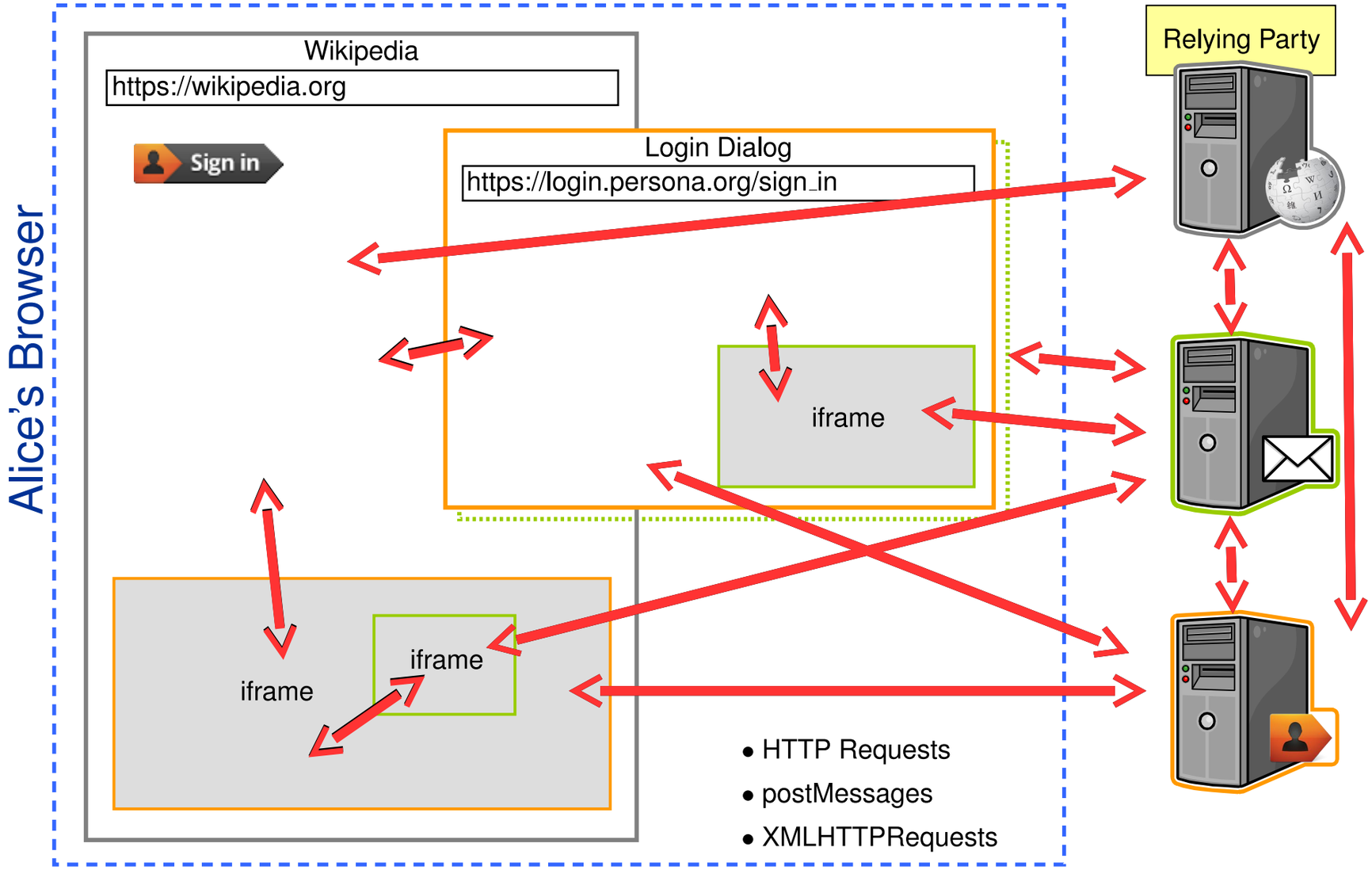


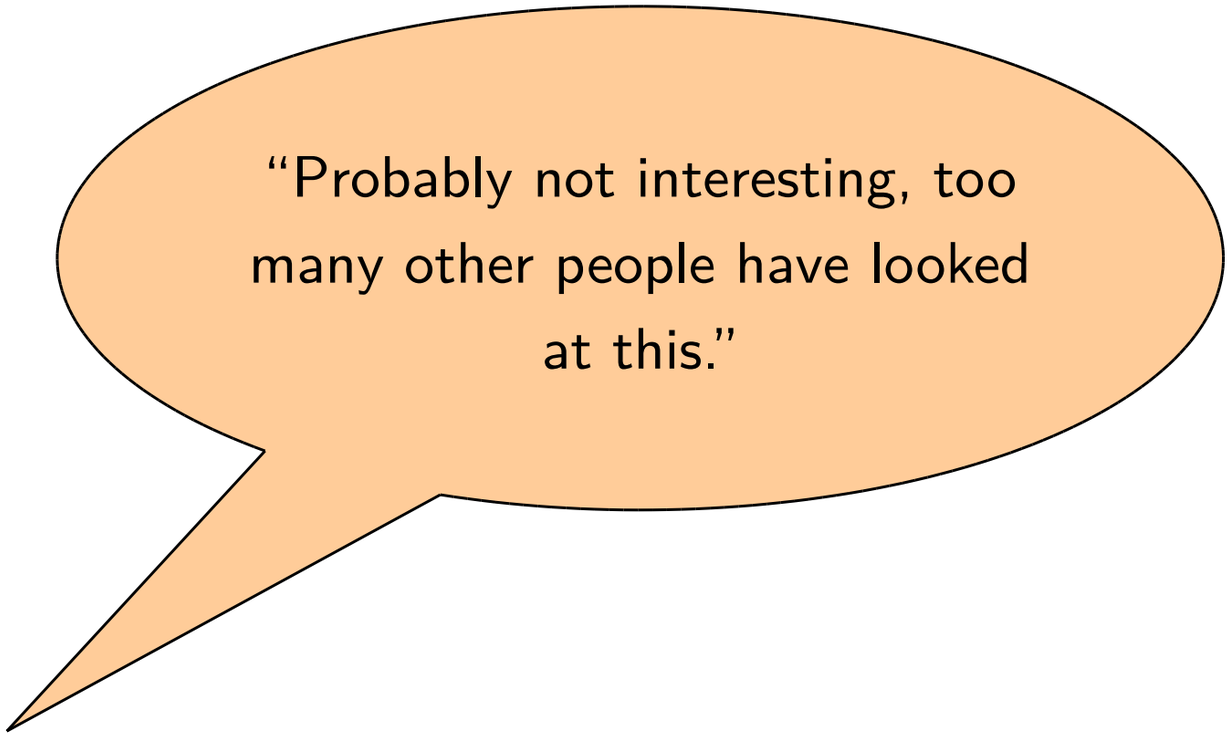
Results

- Analysis of Mozilla's BrowserID (a.k.a. Mozilla Persona) [SP2014, ESORICS2015]
  Main design goal: privacy
  - Found severe attacks: Identity Injection Attack, PostMessage-Based Attack,
  - Proposed fixes for authentication and proved security
  - Privacy broken beyond repair

- Designed our own new SSO system: SPRESSO (https://spresso.me) [CCS2015]
  First provably secure SSO system that provides strong authentication and privacy properties.

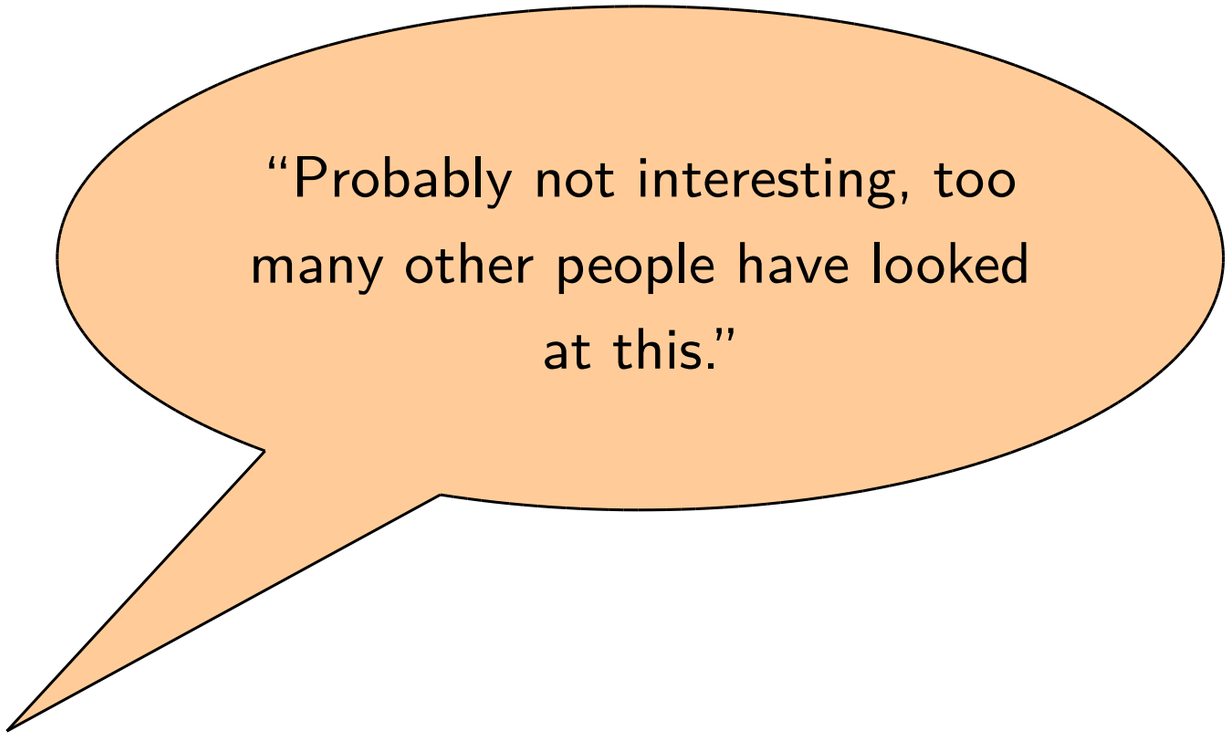
Case Studies – The Obvious Next Targets:
OAuth and OpenID Connect [CCS 2016; CSF 2017]

OAuth 2.0 and OpenID Connect WIM Analyses



PhD Students

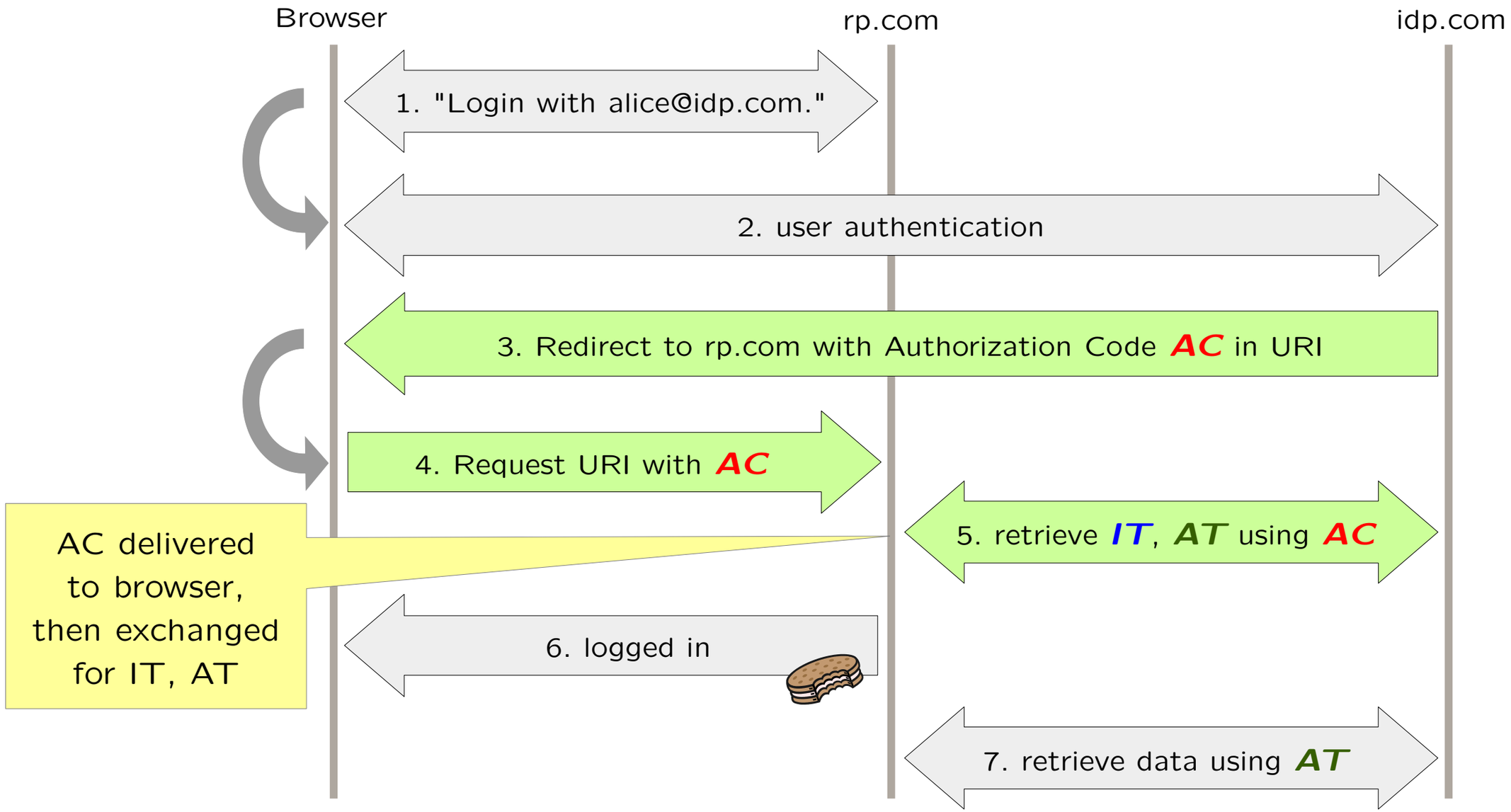
OAuth 2.0 and OpenID Connect WIM Analyses



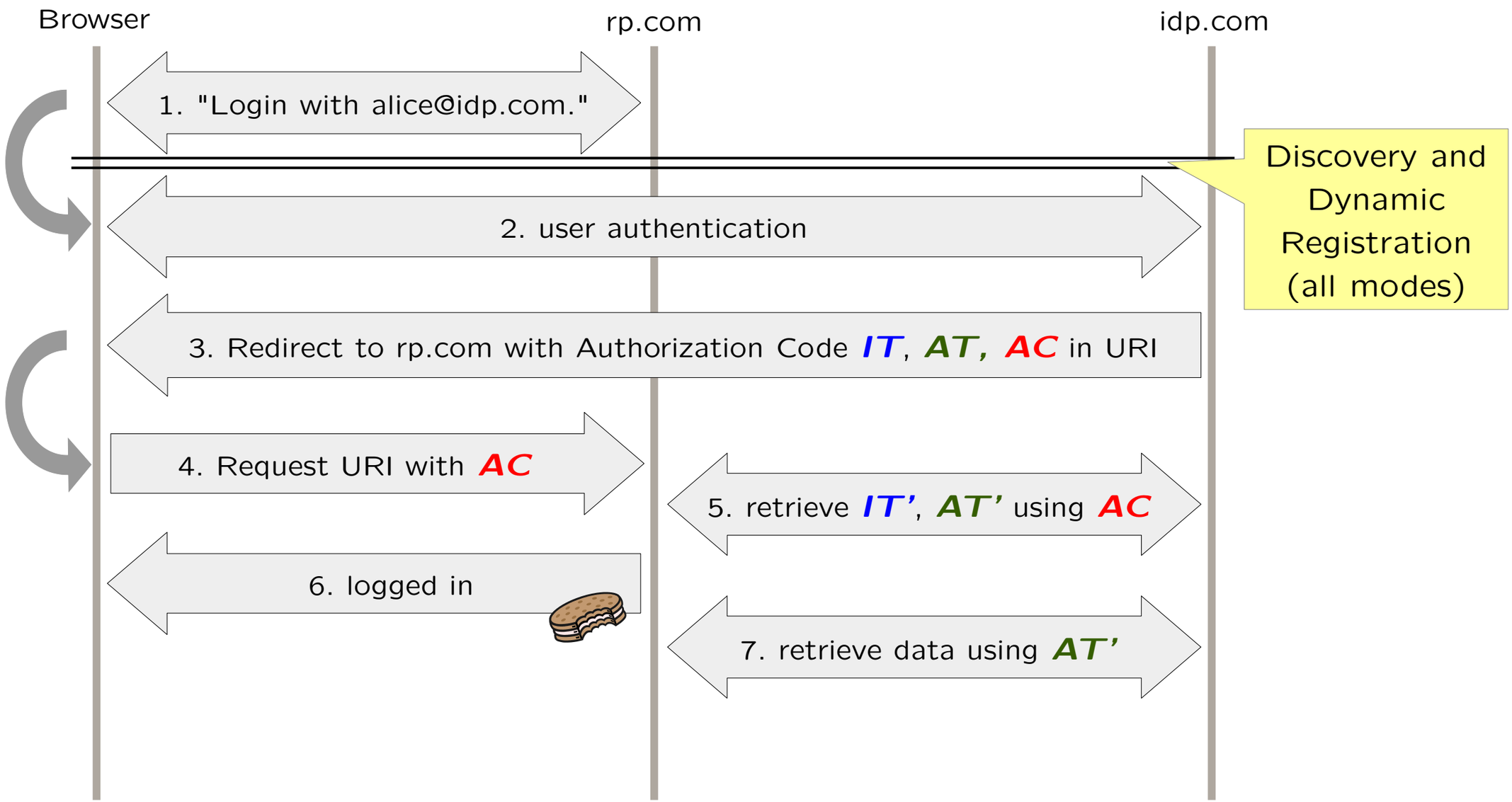
PhD Students

... I insisted

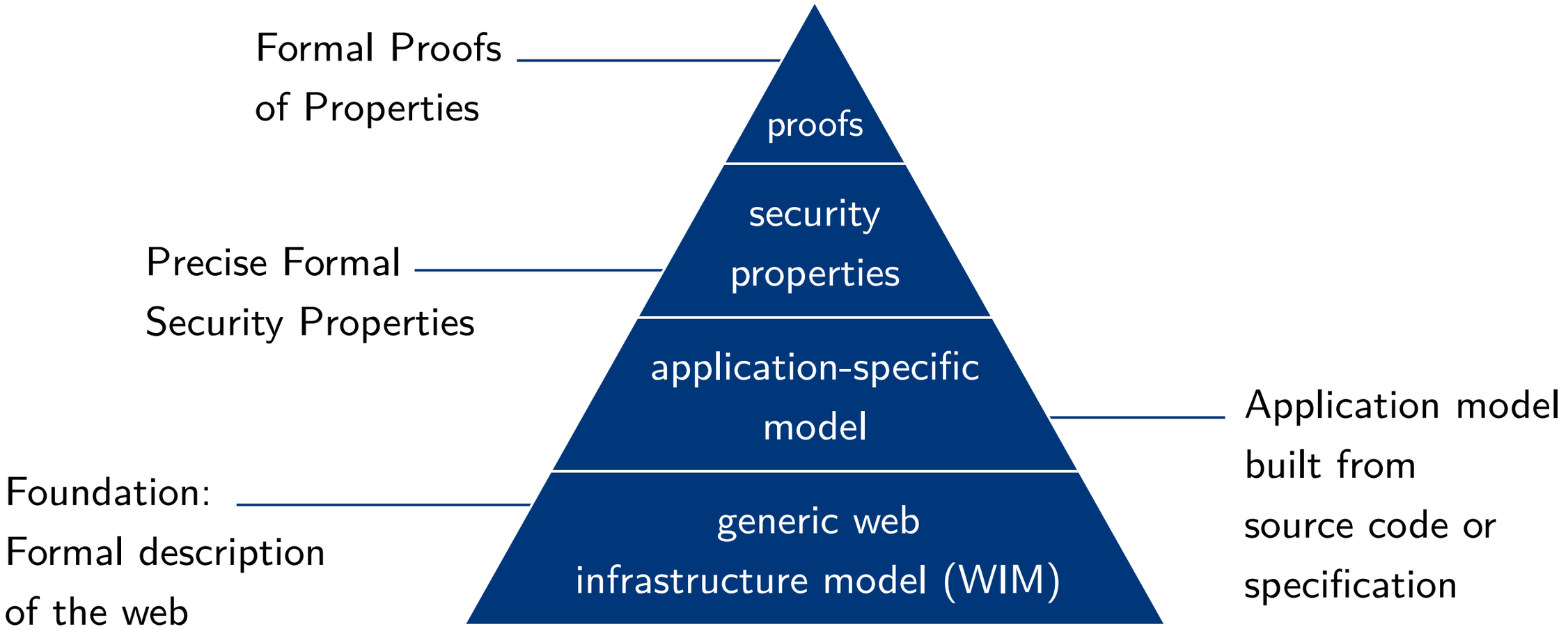
Authorization Code Mode



Discovery and Dynamic Registration



How to use the WIM?



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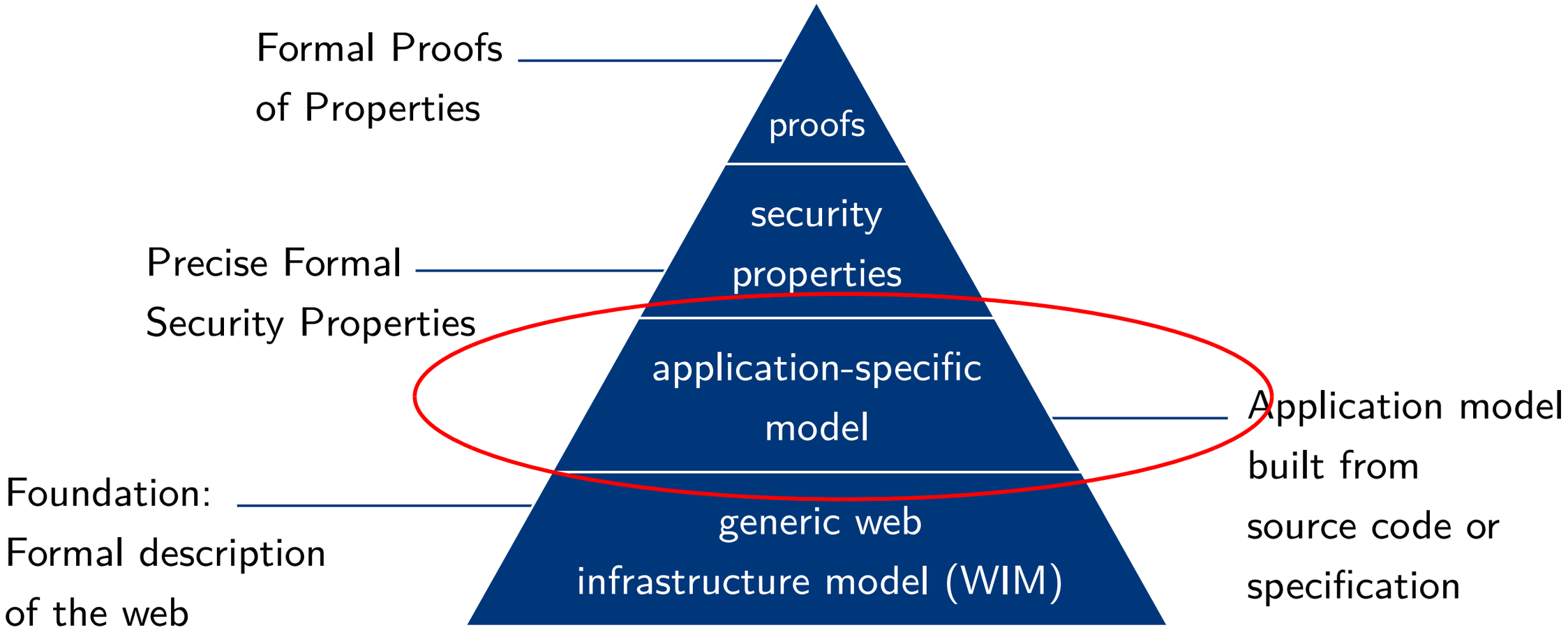

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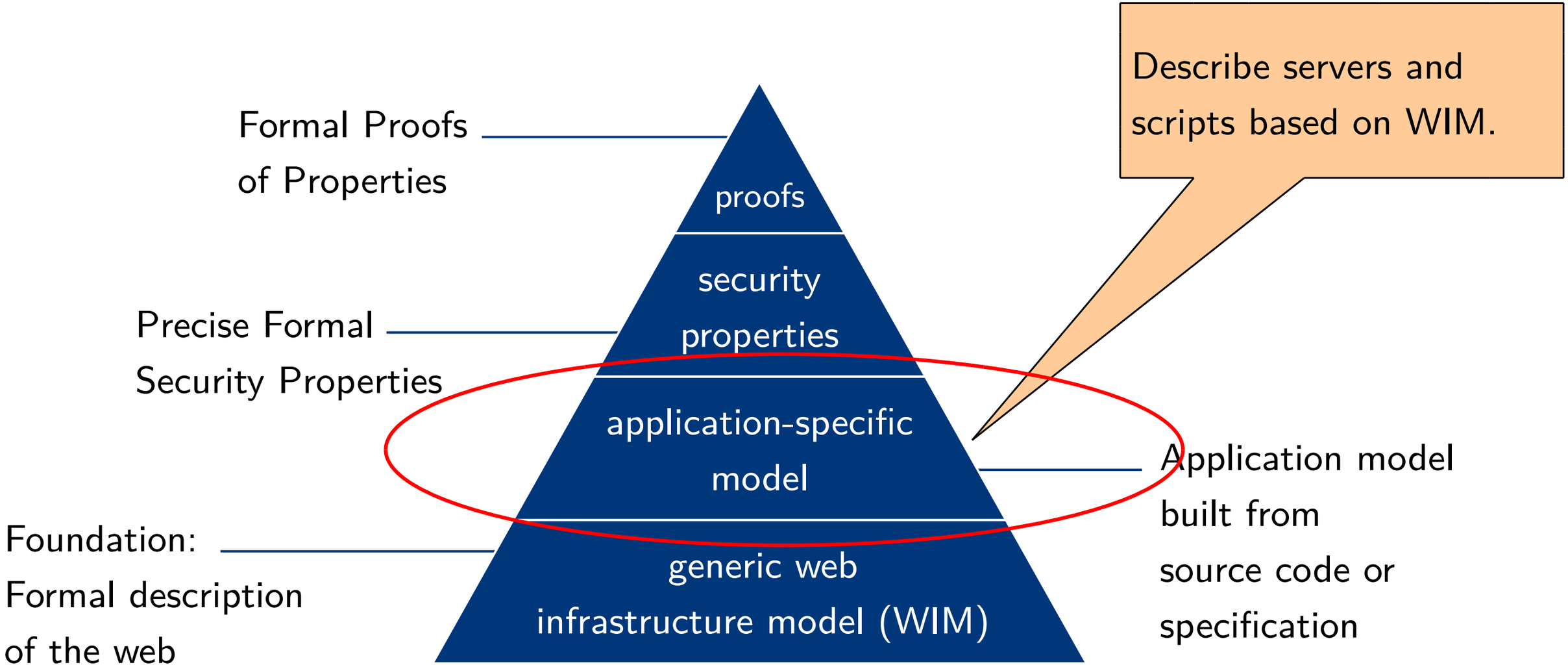
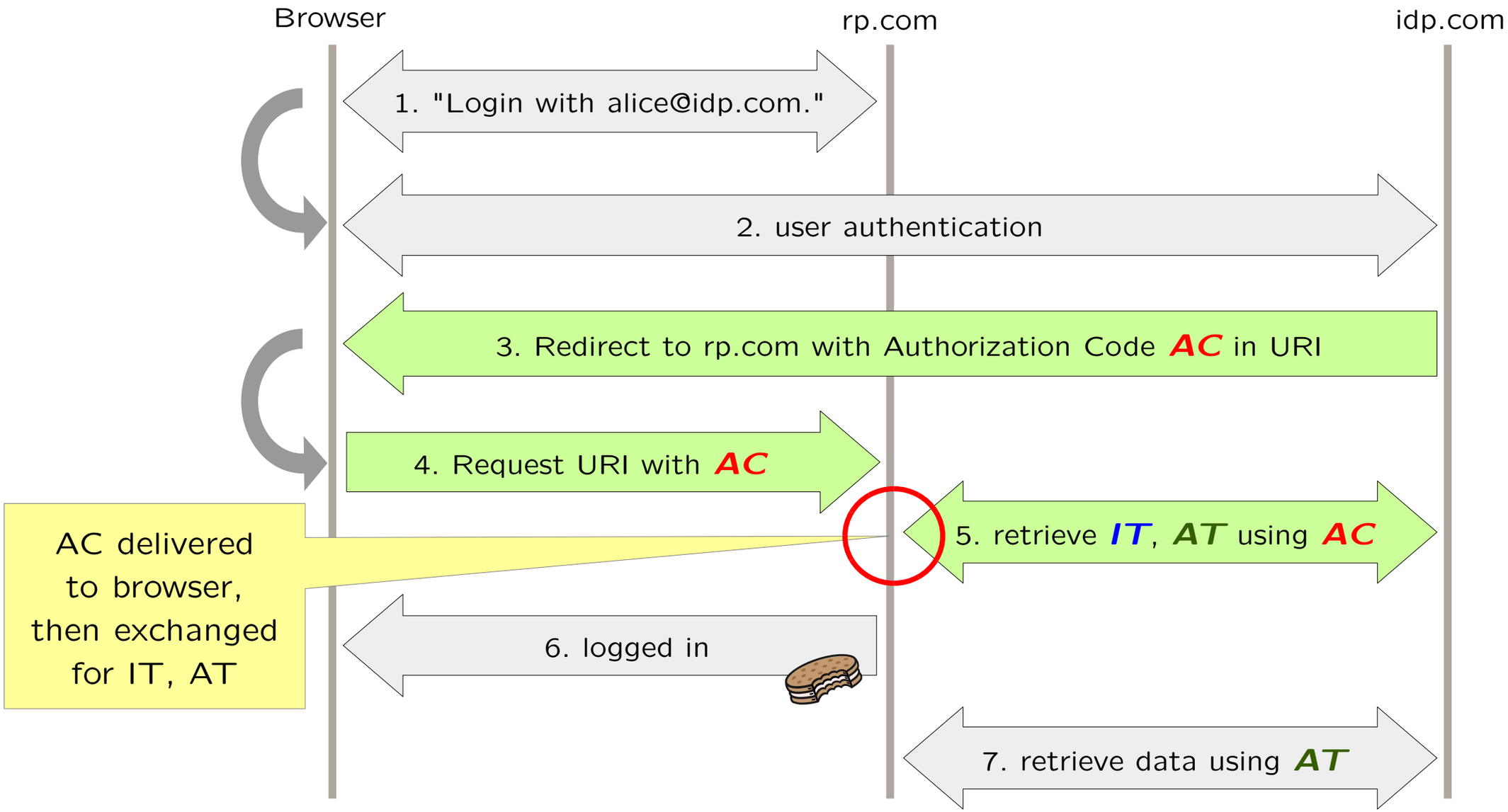


Authorization Code Mode

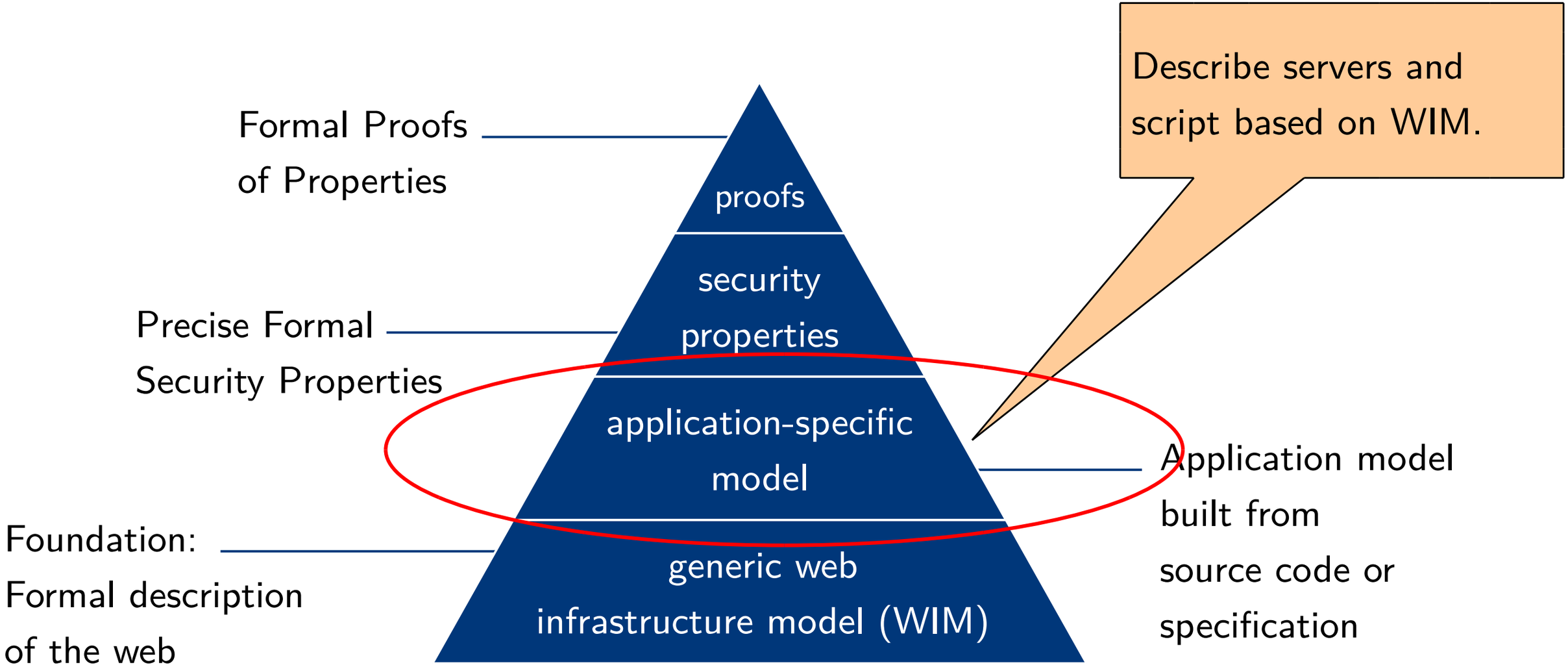


Example: RP Checks an ID Token

Algorithm 20 Relying Party R^r: Check id token.

1: **function** CHECK_ID_TOKEN(*sessionId*, *id_token*, s') → **Check id token validity and create service session.**
2: **let** *session* := s'.`sessions`[*sessionId*] → Retrieve session data.
3: **let** *identity* := *session*[`identity`]
4: **let** *issuer* := s'.`issuerCache`[*identity*] → Retrieve issuer.
5: **let** *oidcConfig* := s'.`oidcConfigCache`[*issuer*] → Retrieve OIDC configuration for that issuer.
6: **let** *credentials* := s'.`clientCredentialsCache`[*issuer*] → Retrieve OIDC credentials for issuer.
7: **let** *jwks* := s'.`jwksCache`[*issuer*] → Retrieve signing keys for issuer.
8: **let** *data* := extractmsg(*id_token*) → Extract contents of signed id token.
9: **if** *data*[`iss`] $\not\equiv$ *issuer* **then**
10: **stop** → Check the issuer.
11: **if** *data*[`aud`] $\not\equiv$ *credentials*[`client_id`] **then**
12: **stop** → Check the audience against own client id.
13: **if** checksig(*id_token*,*jwks*) $\not\equiv \top$ **then**
14: **stop** → Check the signature of the id token.
15: **if** *nonce* $\in$ *session* $\wedge$ *data*[`nonce`] $\not\equiv$ *session*[`nonce`] **then**
16: **stop** → If a nonce was used, check its value.
17: **let** s'.`sessions`[*sessionId*][`loggedInAs`] := $\langle$*issuer*,*data*[`sub`]$\rangle$ → User is now logged in. Store user identity and issuer.
18: **let** s'.`sessions`[*sessionId*][`serviceSessionId`] := v_4 → Choose a new service session id.
19: **let** *request* := *session*[`redirectEpRequest`] → Retrieve stored meta data of the request from the browser to the redir. endpoint in order to respond to it now. The request's meta data was stored in PROCESS_HTTPS_REQUEST (Algorithm 17).
20: **let** *headers* := [`ReferrerPolicy:origin`]
21: **let** *headers*[`Set-Cookie`] := [`serviceSessionId:`$\langle v_4, \top, \top, \top\rangle$] → Create a cookie containing the service session id.
22: **let** m' := $\mathsf{enc_s}(\langle$`HTTPResp`,*request*[`message`].`nonce`,200,*headers*,`ok`$\rangle$,*request*[`key`]) → Respond to browser's request to the redirection endpoint.
23: **stop** $\langle\langle$*request*[`sender`],*request*[`receiver`],$m'\rangle\rangle$, s'

How to use the WIM?



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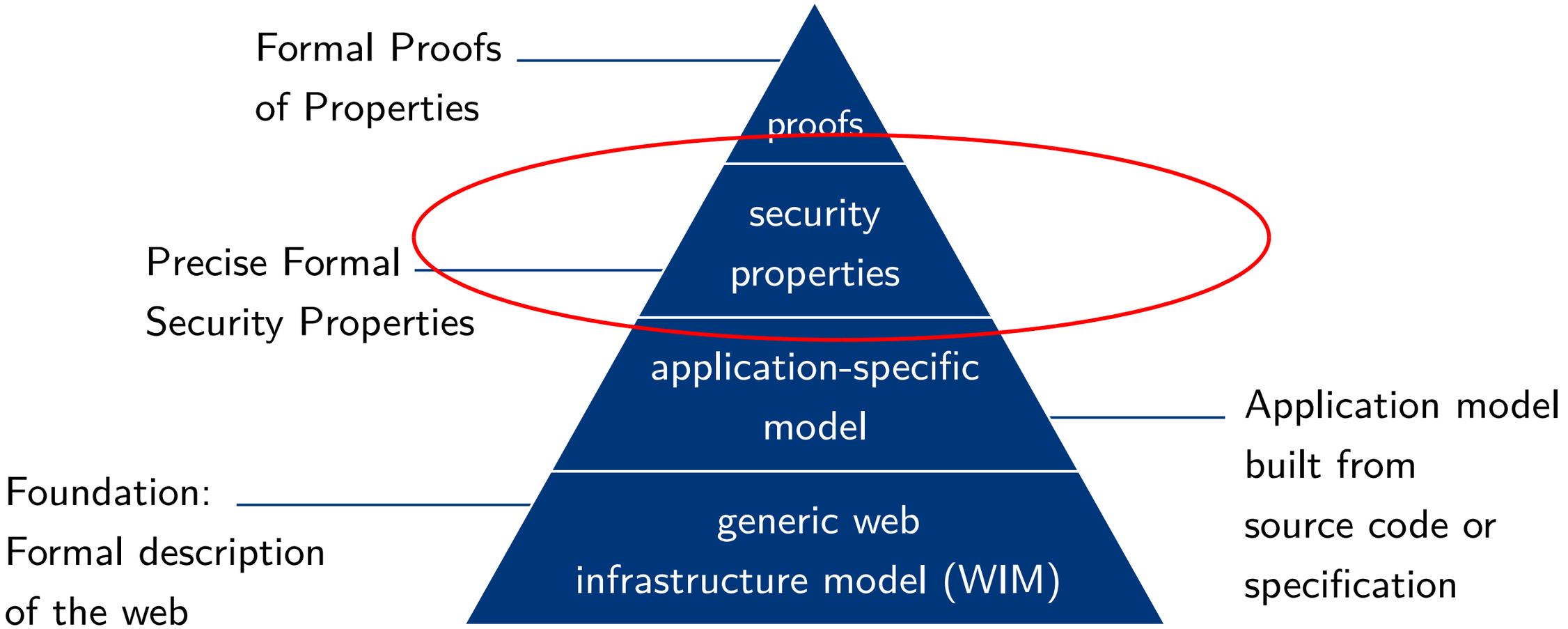
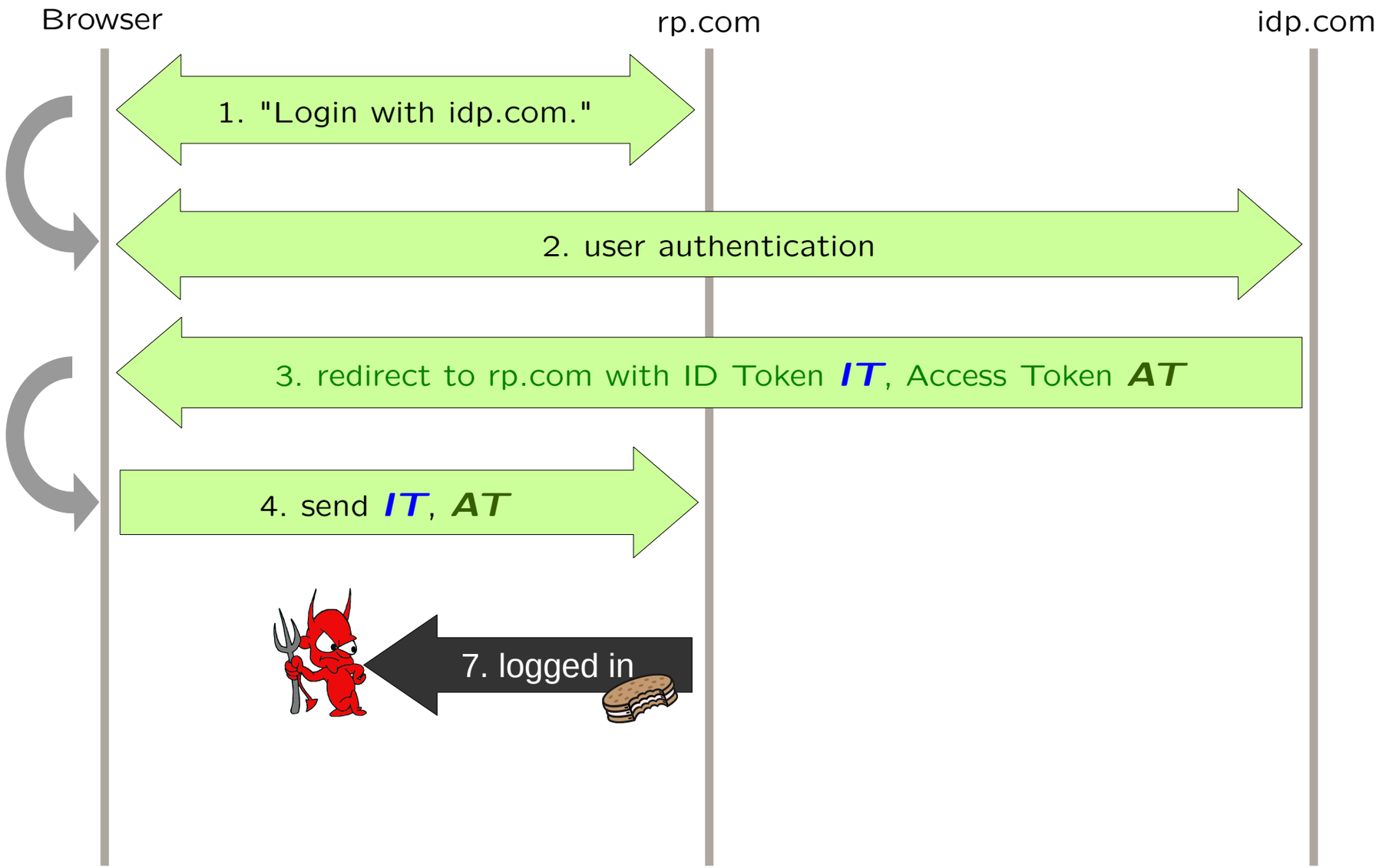


Authentication Property



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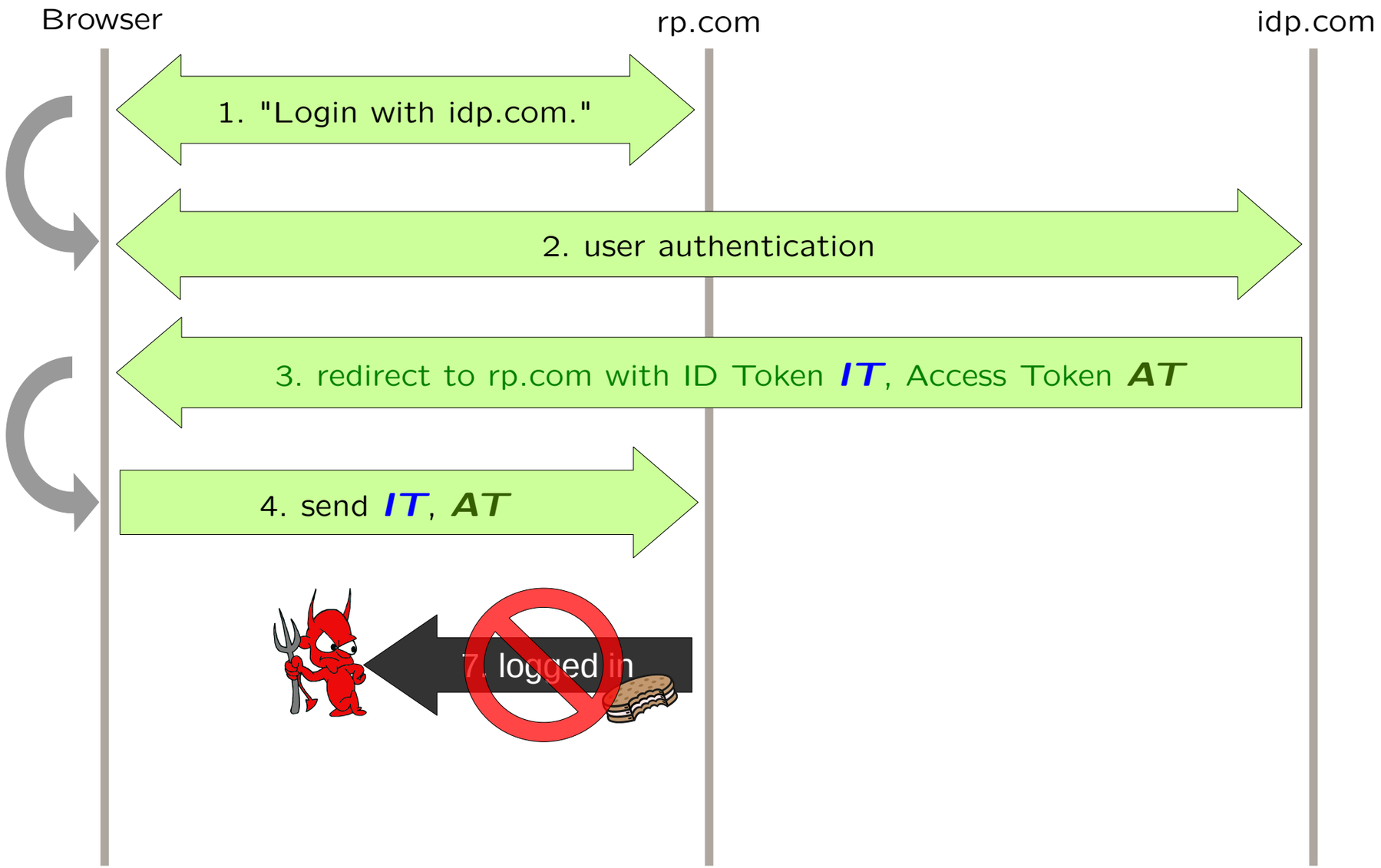
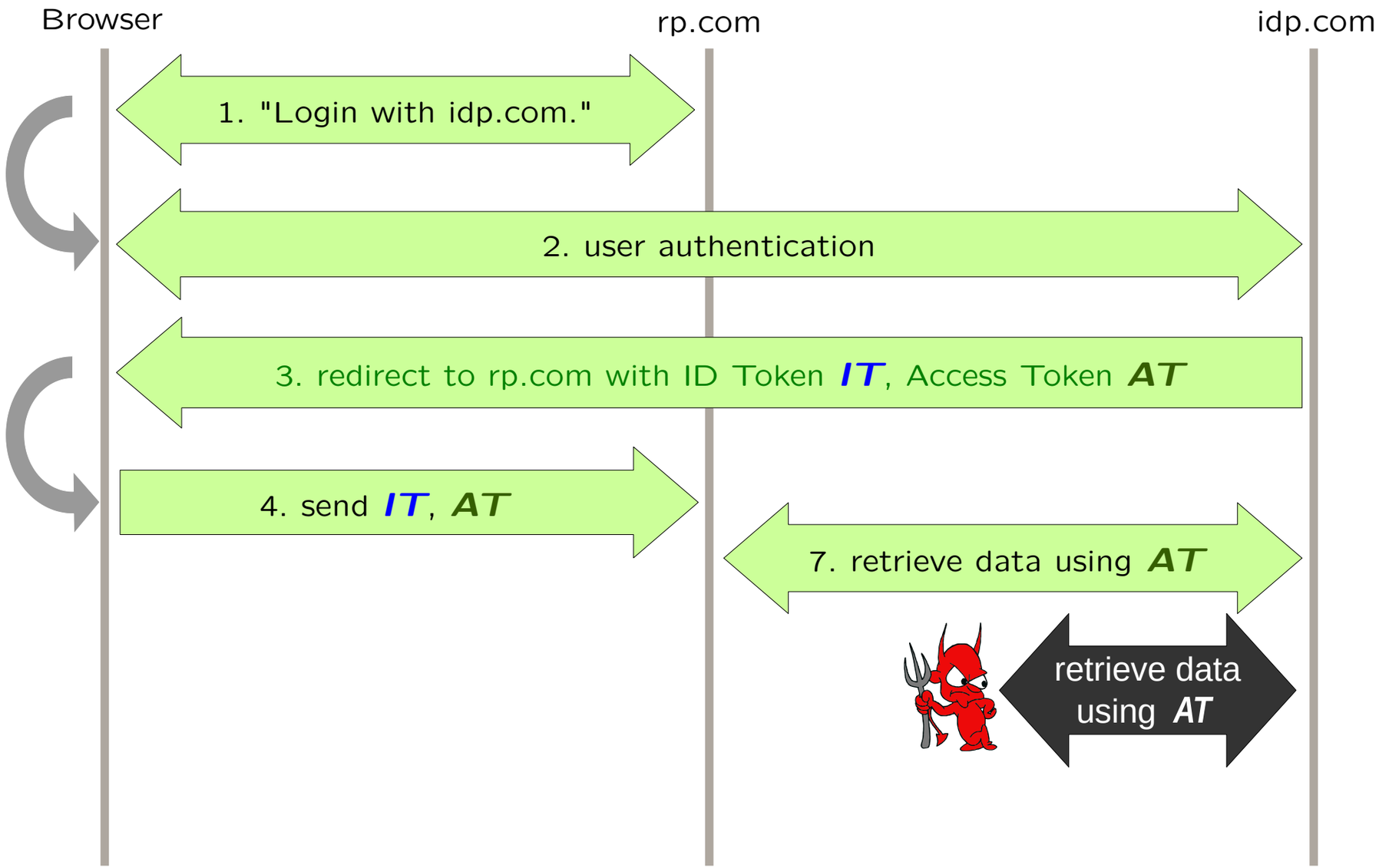


Authentication Property

Definition 46 (Authentication Property). Let $\mathcal{OWS}^n$ be an OAuth web system with a network attacker. We say that $\mathcal{OWS}^n$ *is secure w.r.t. authentication* iff for every run ρ of $\mathcal{OWS}^n$, every state (S^j, E^j, N^j) in ρ, every $r \in \mathsf{RP}$ that is honest in S^j, every $i \in \mathsf{IDP}$, every $g \in \mathsf{dom}(i)$, every $u \in \mathbb{S}$, every RP service token of the form $\langle n, \langle u, g \rangle \rangle$ recorded in $S^j(r).\texttt{serviceTokens}$, and n being derivable from the attackers knowledge in S^j (i.e., $n \in d_\emptyset(S^j(\mathsf{attacker}))$), then the browser b owning u is fully corrupted in S^j (i.e., the value of *isCorrupted* is `FULLCORRUPT`), some $r' \in \mathsf{trustedRPs}(\mathsf{secretOfID}(\langle u, g \rangle))$ is corrupted in S^j, or i is corrupted in S^j.

Authorization Property



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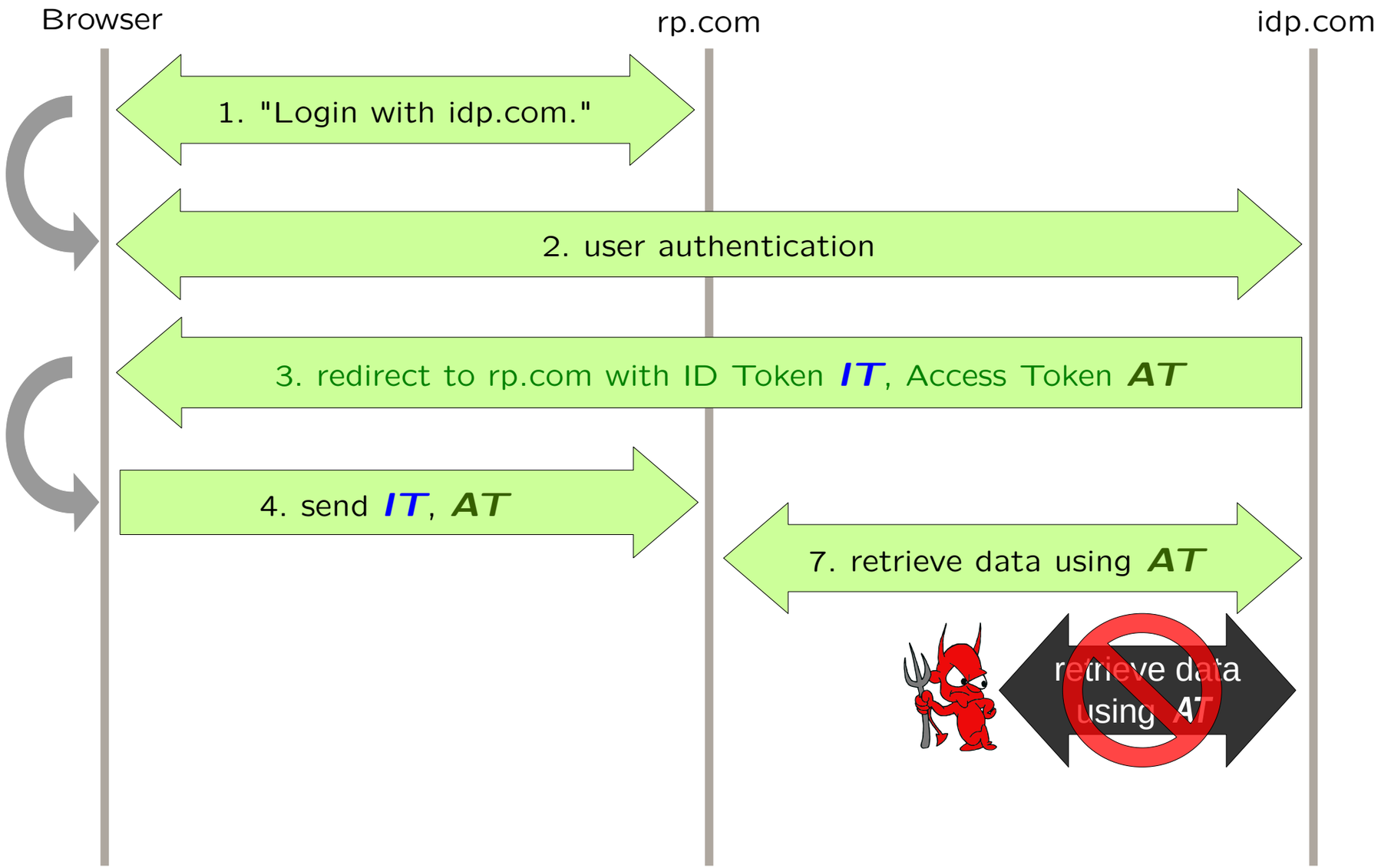
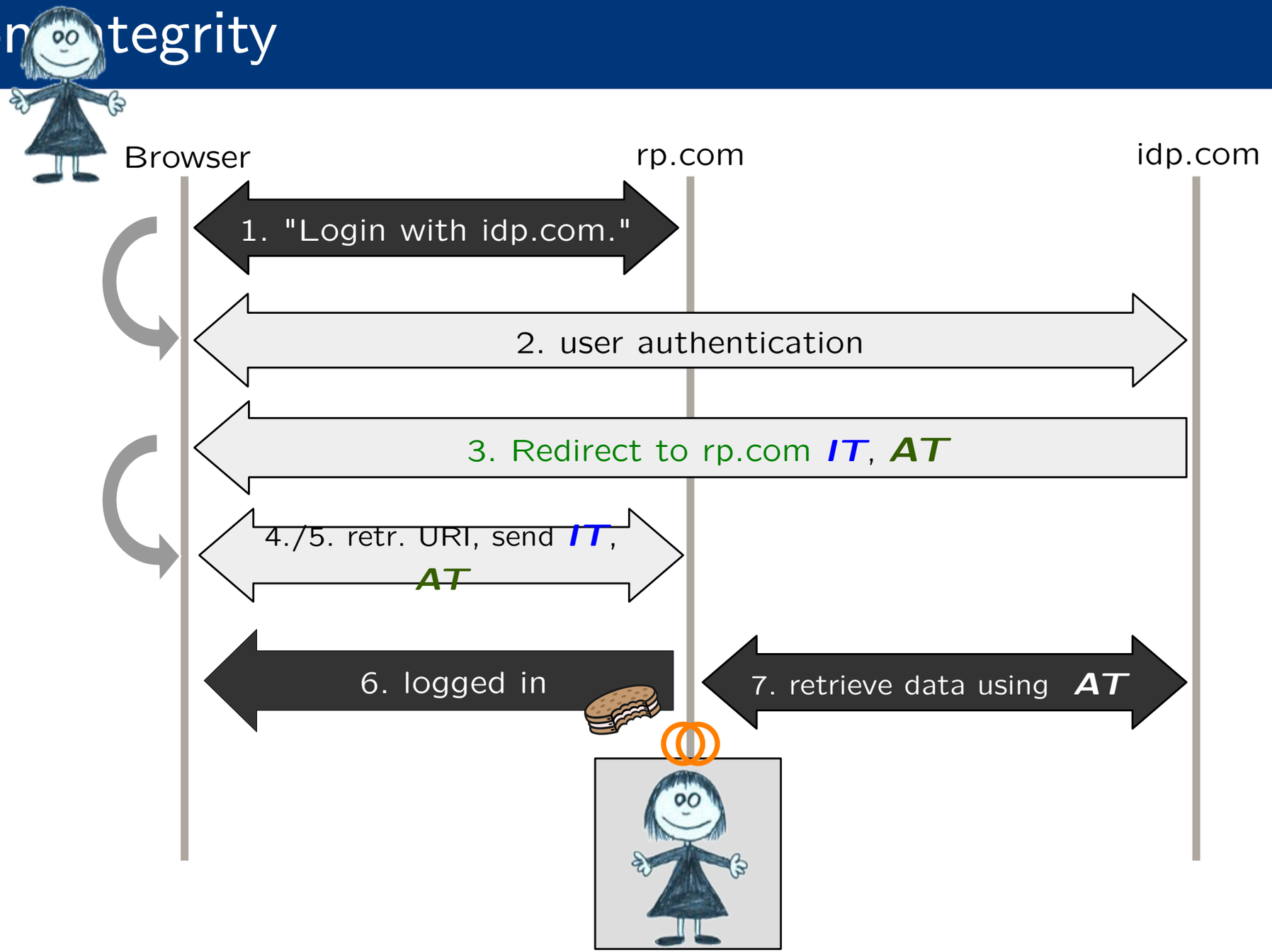


Session Integrity



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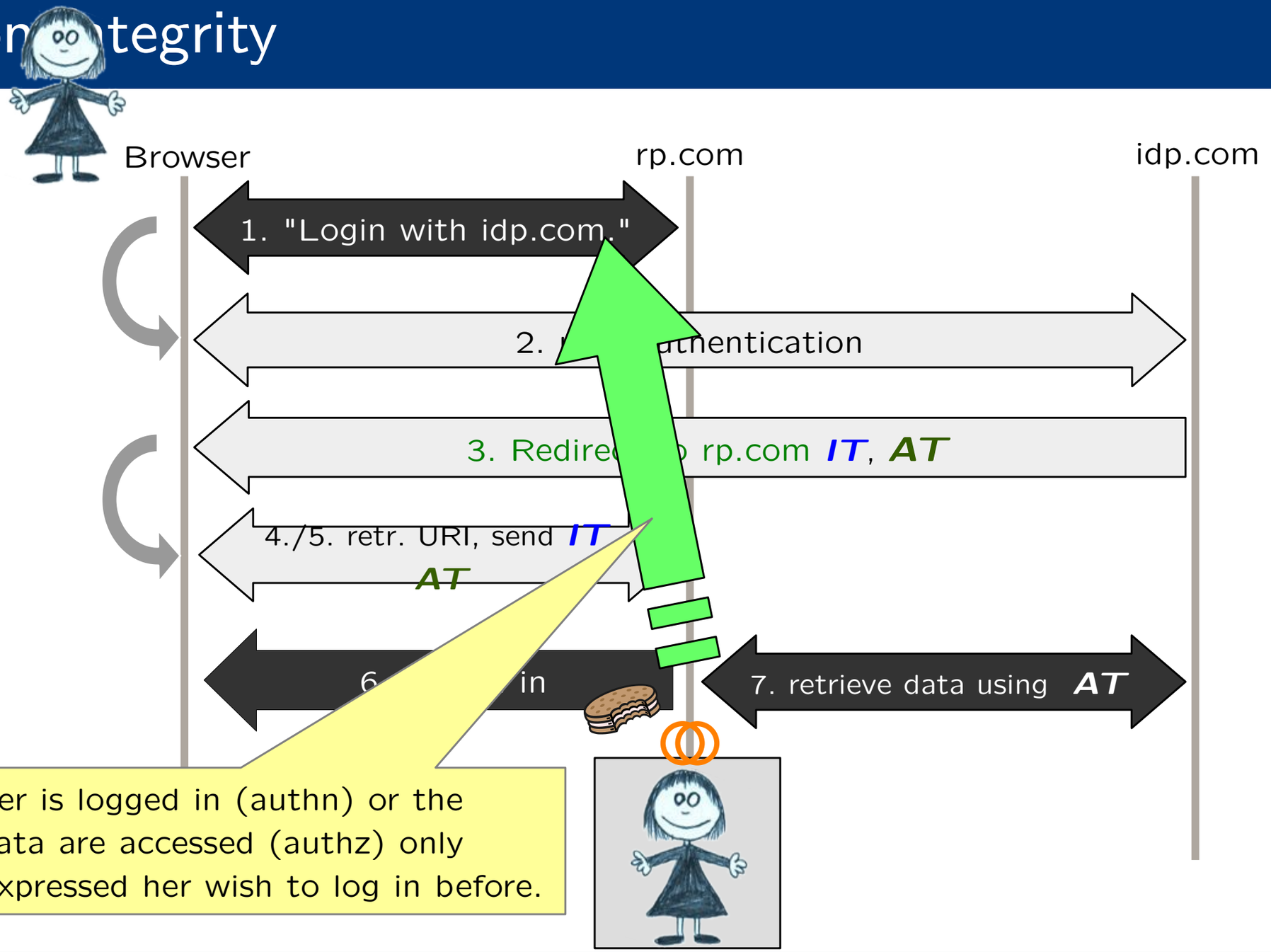
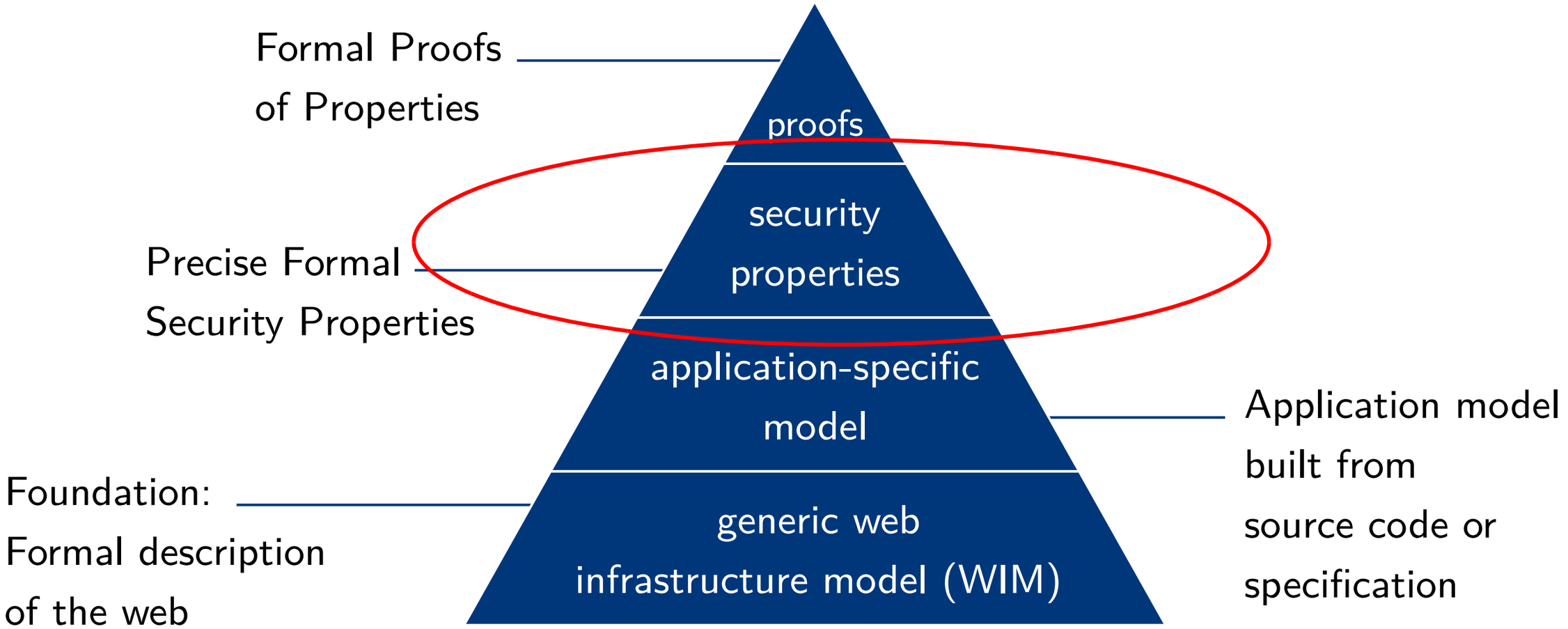


The user is logged in (authn) or the
user's data are accessed (authz) only
if the user expressed her wish to log in before.

How to use the WIM?



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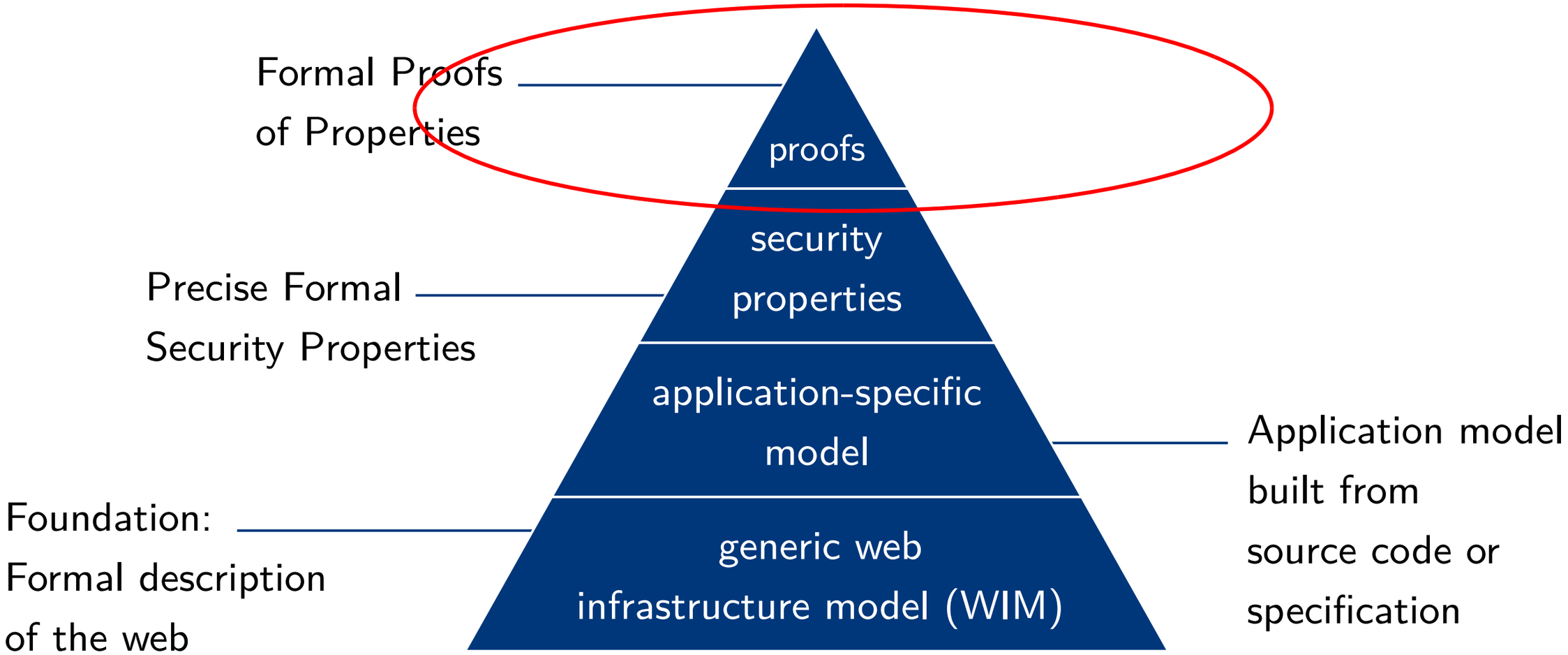
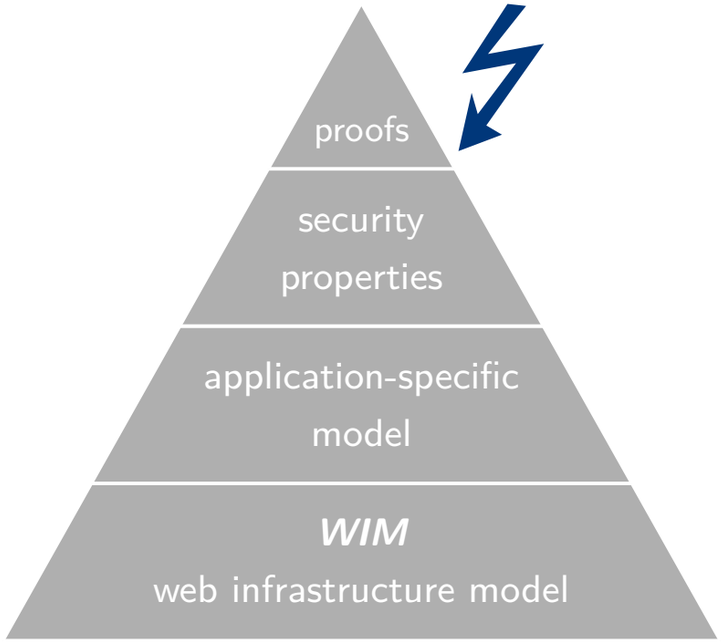


OAuth 2.0: New Attacks

OAuth 2.0 had been analyzed many times before,
but not in a comprehensive formal model.

New attacks:

- 307 Redirect Attack
- Identity Provider Mix-Up Attack (new class of attacks)
- State Leak Attack
- Naïve Client Session Integrity Attack
- Across Identity Provider State Reuse Attack



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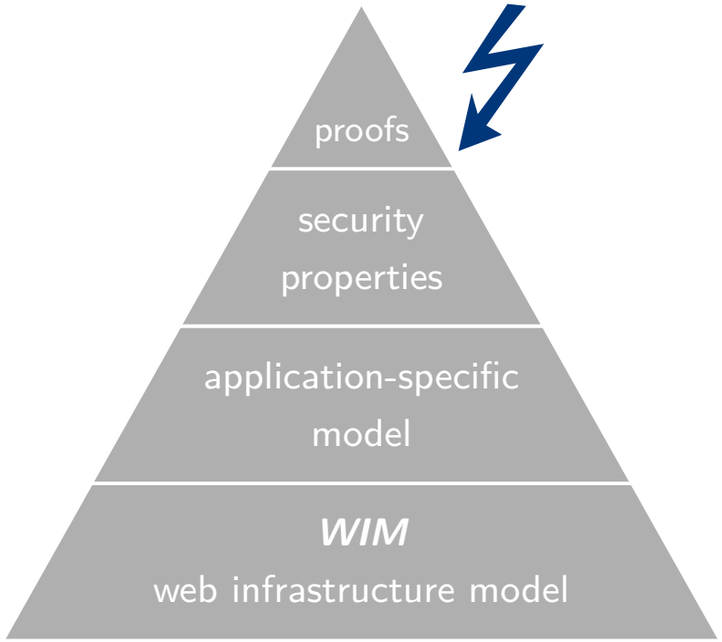
Similary for OpenID Connect.



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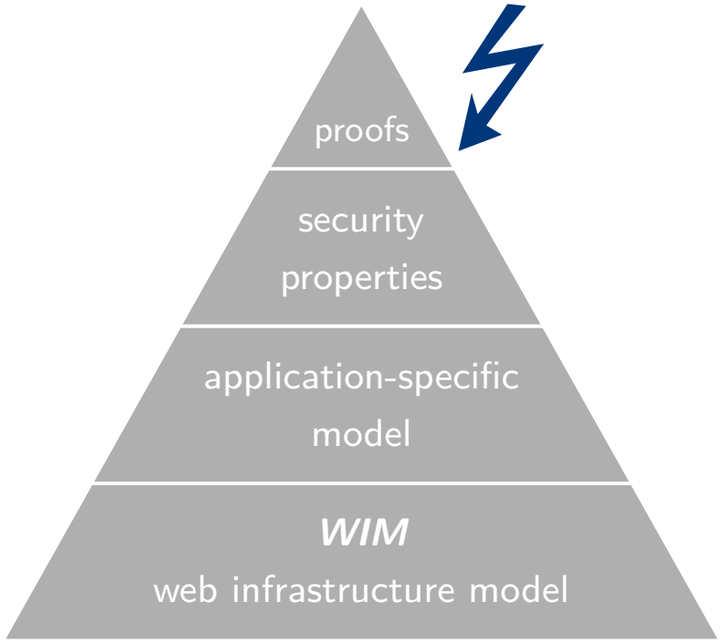
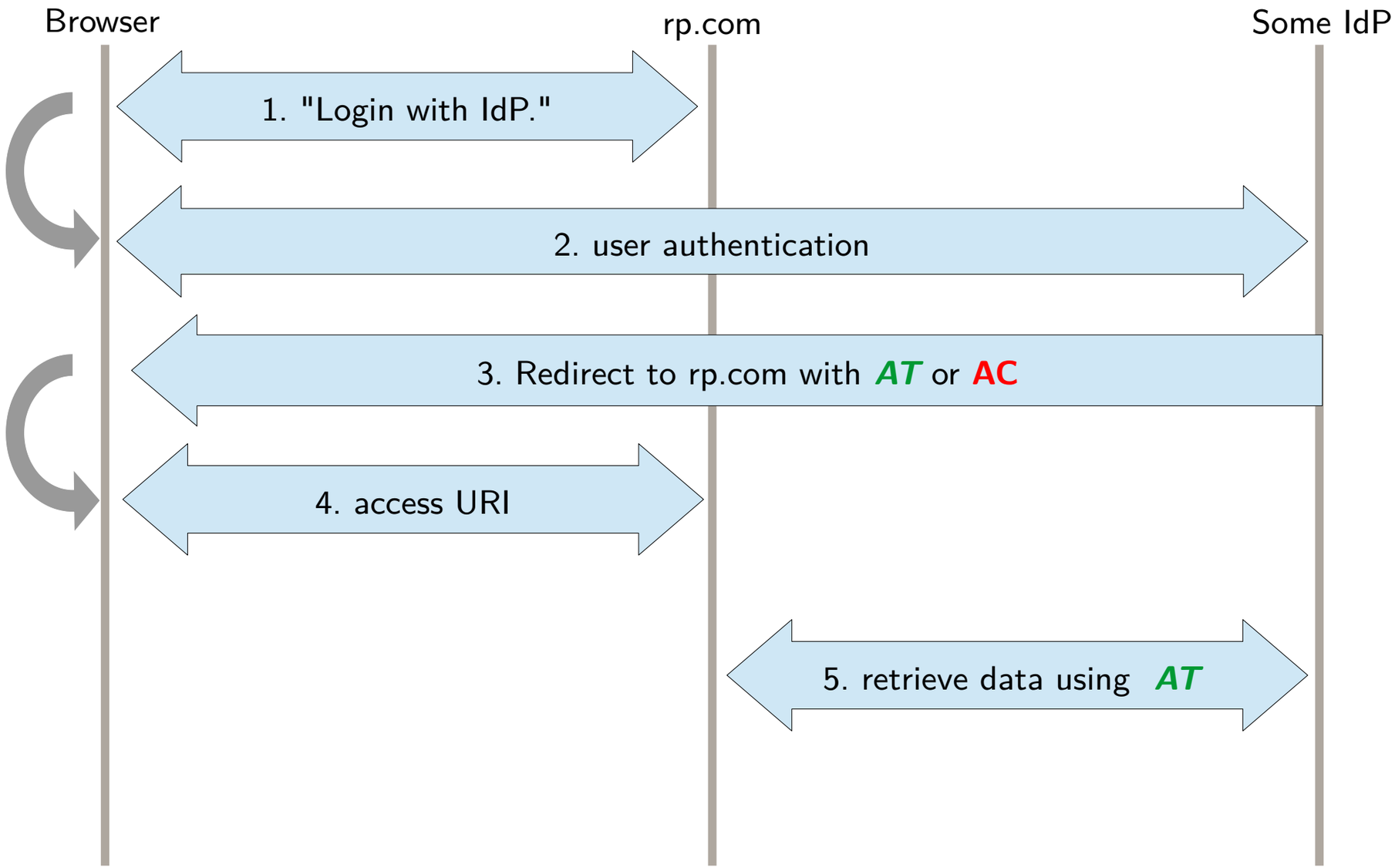
- Across Identity Provider State Reuse Attack

Similary for OpenID Connect.



307 Redirect Attack



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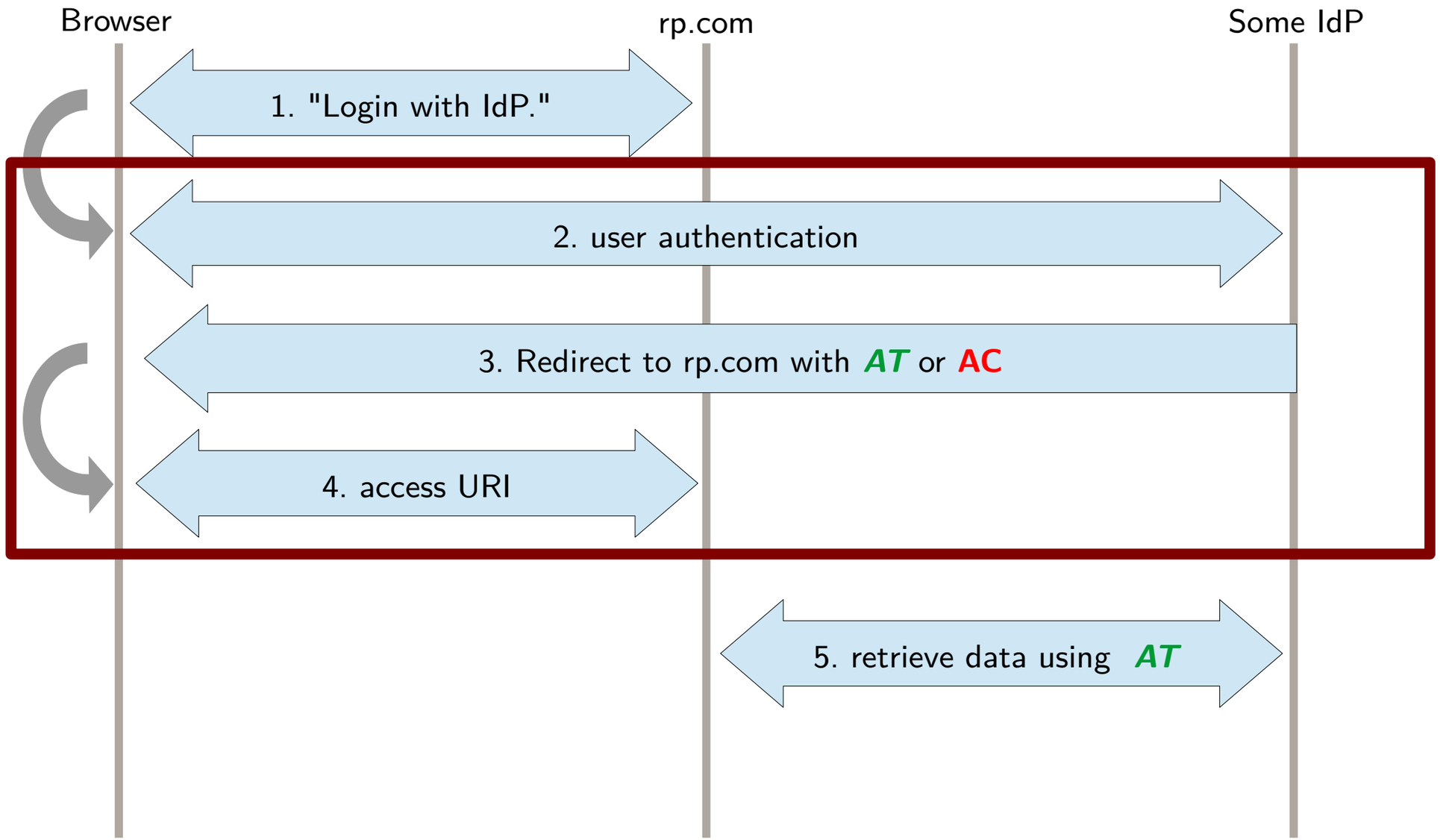

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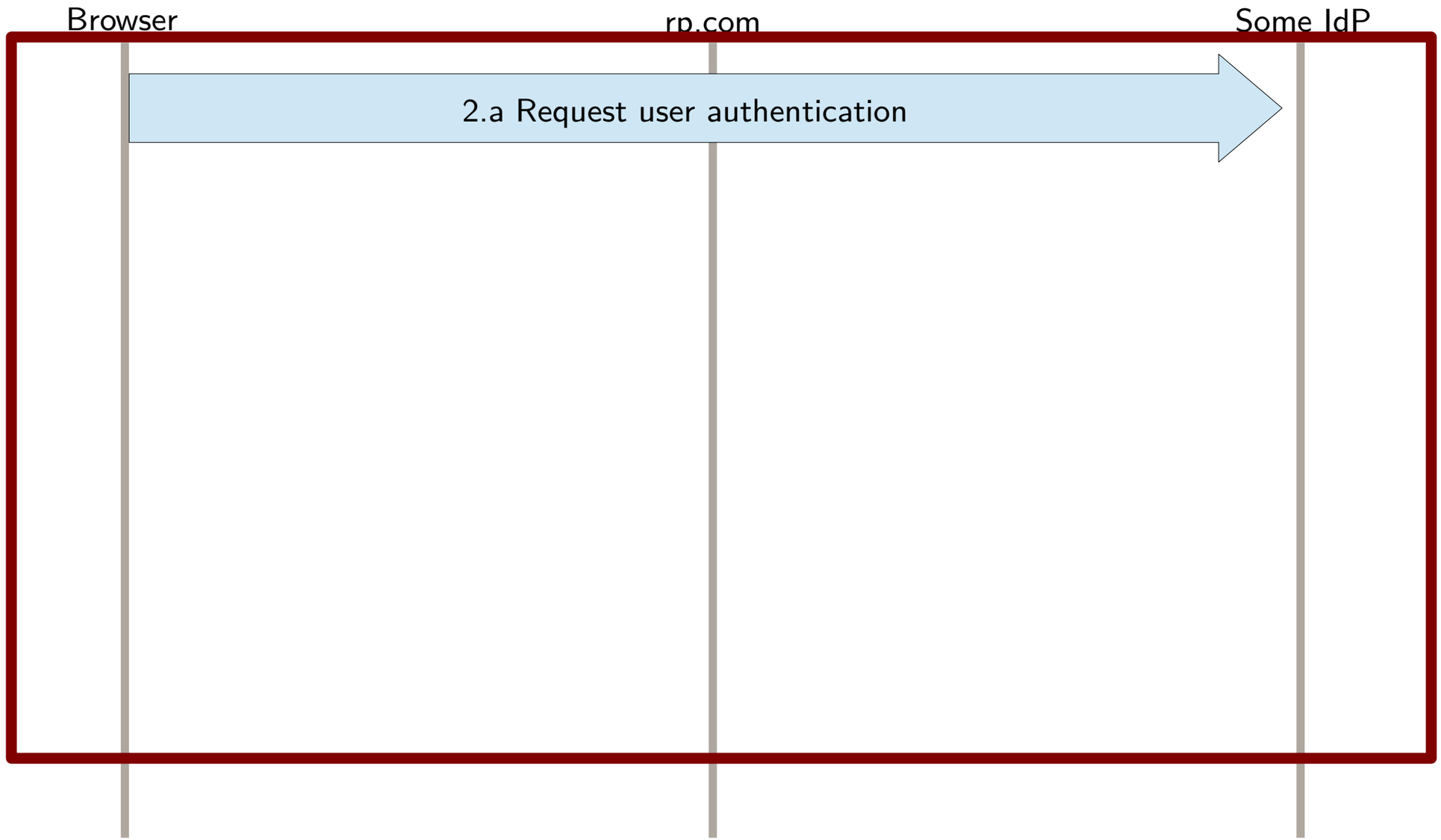

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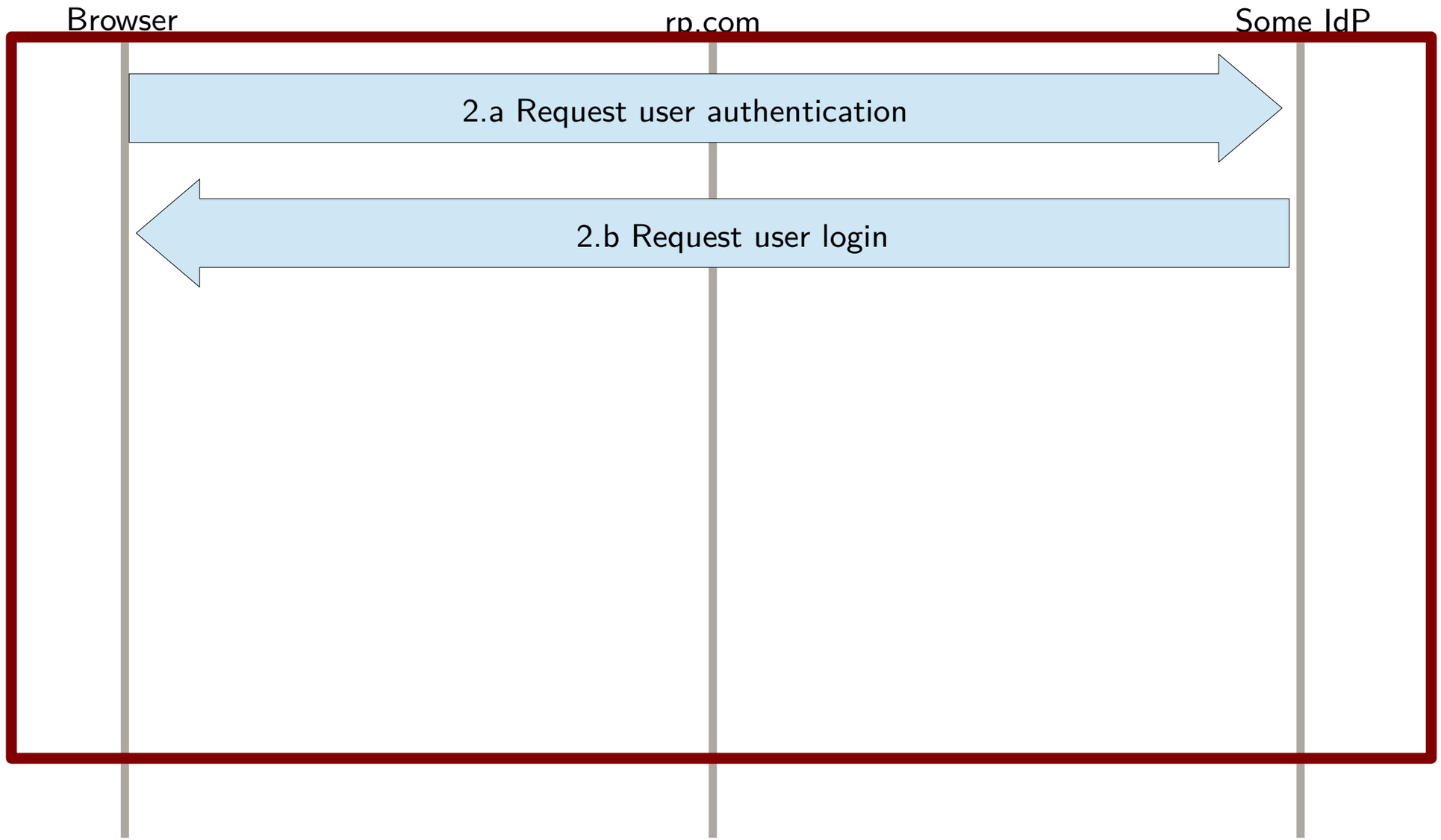

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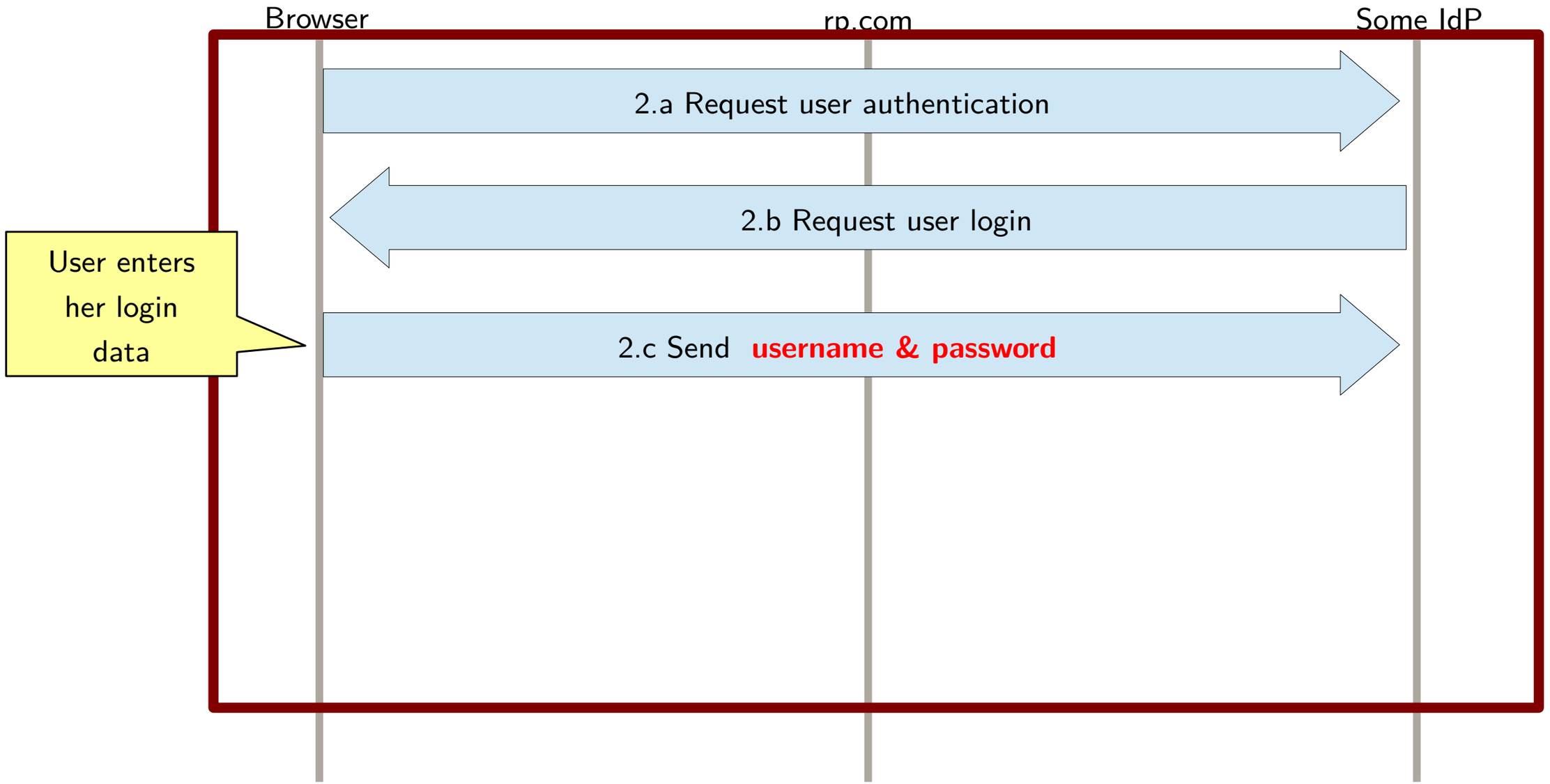

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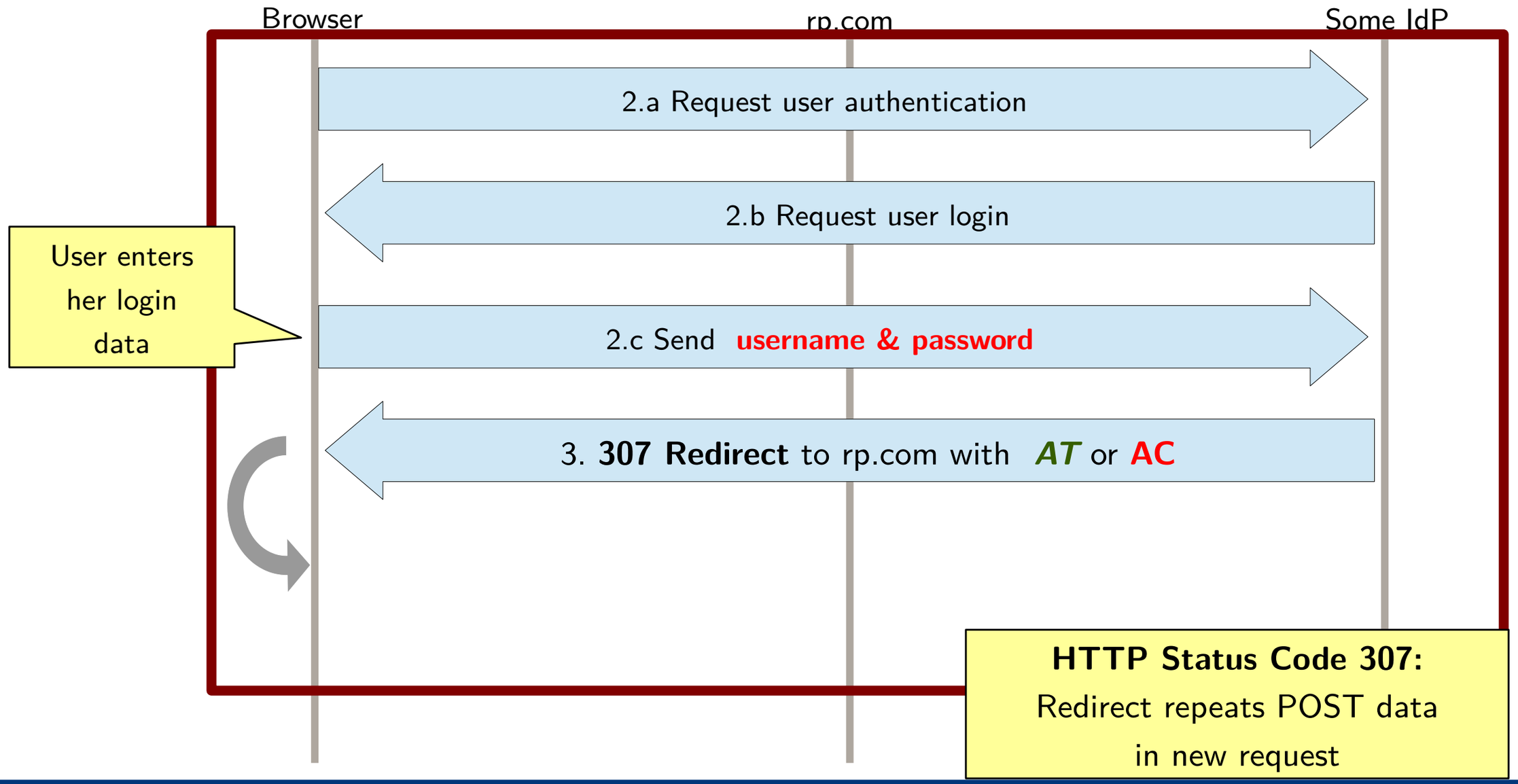

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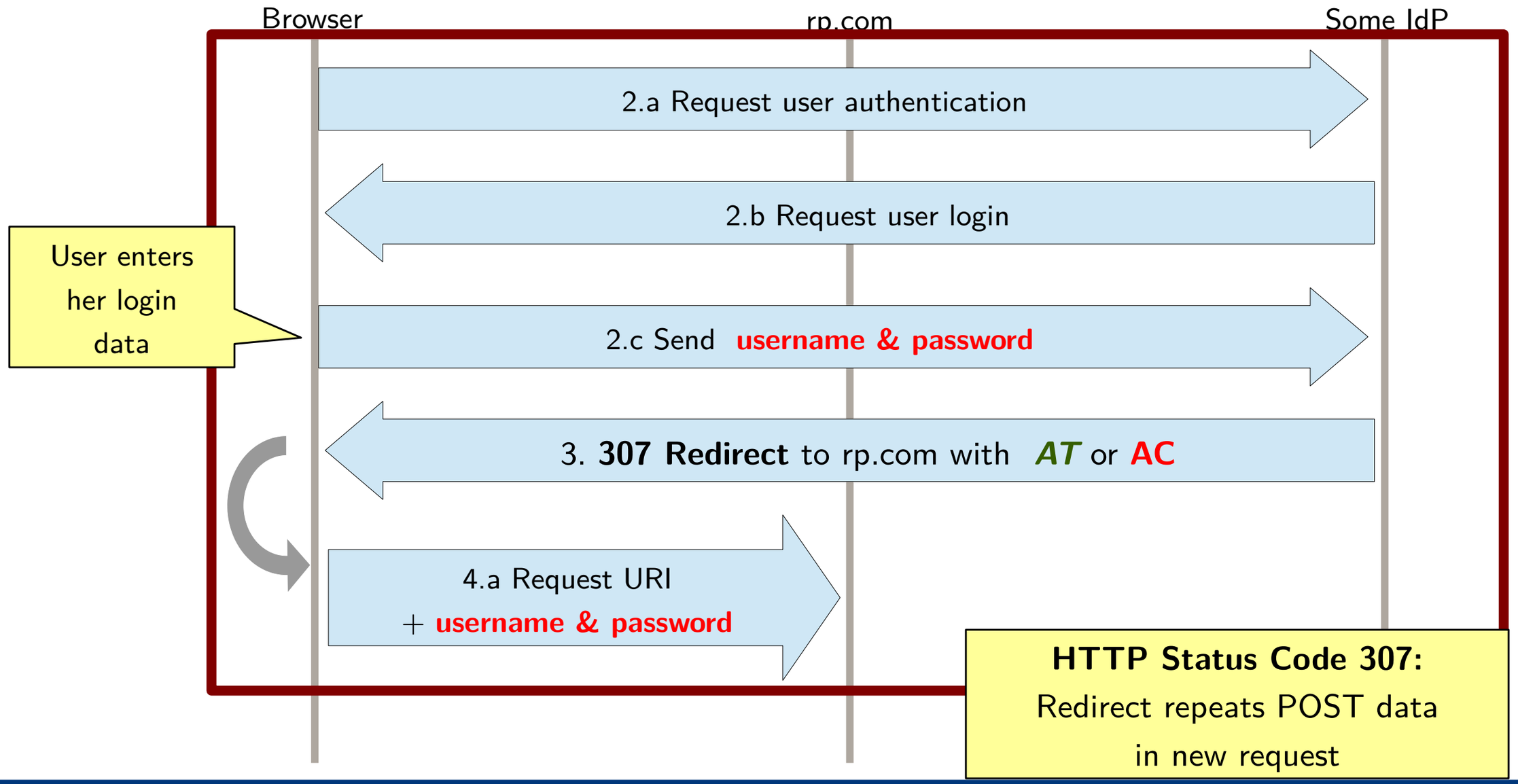

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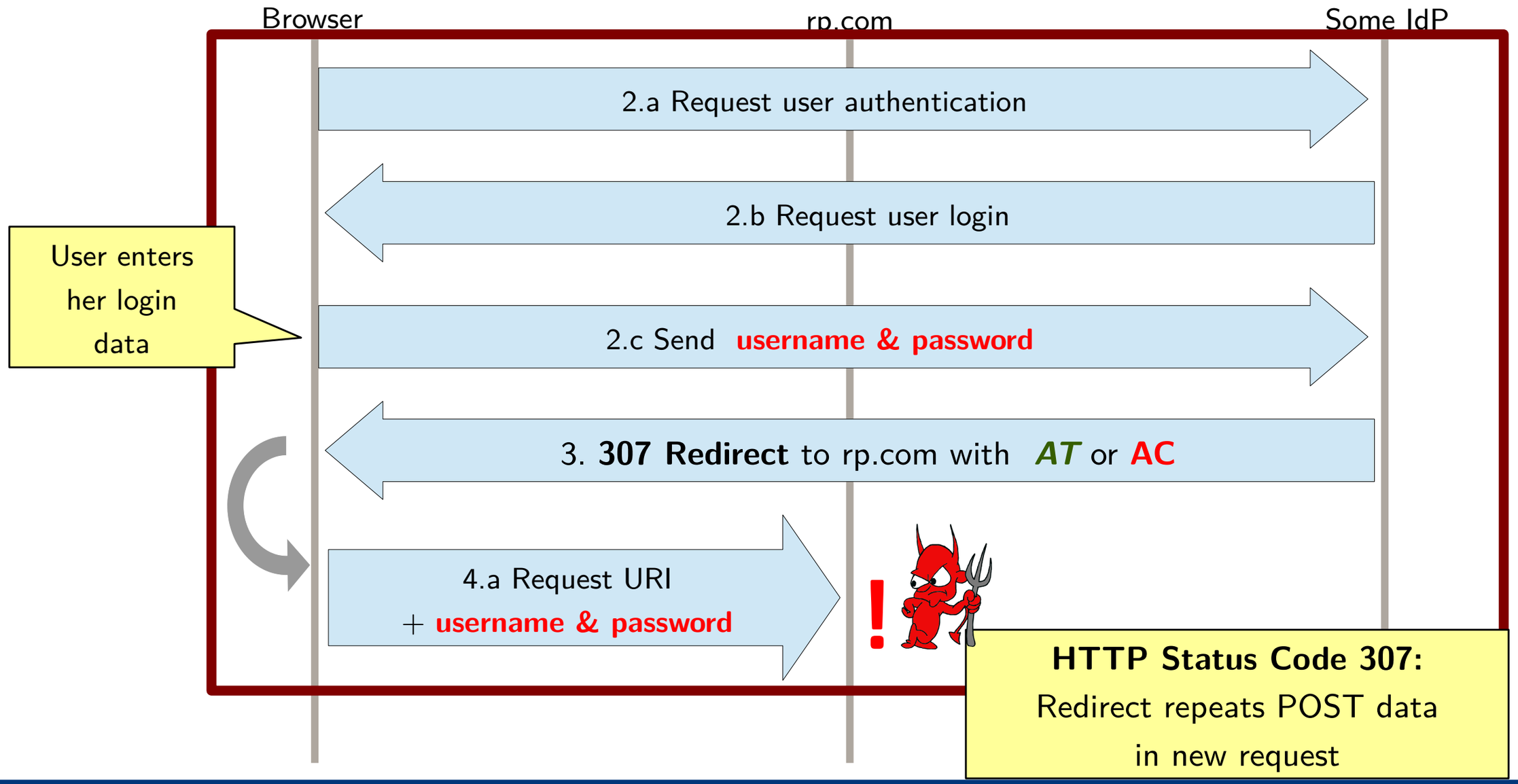


307 Redirect Attack

The attacker receives the username and password of the user.

OAuth standard says:

> **1.7. HTTP Redirections**
>
> This specification makes extensive use of HTTP redirections, in which the client or the authorization server directs the resource owner's user-agent to another destination. While the examples in this specification show the use of the HTTP 302 status code, any other method available via the user-agent to accomplish this redirection is allowed and is considered to be an implementation detail.

Mitigation:

Use status code 303 or any other method that does not forward POST data.

Theorem

Theorem

We proposed fixed to the standards and proved them secure:

Theorem

OAuth 2.0 and OIDC with fixes fulfill security properties

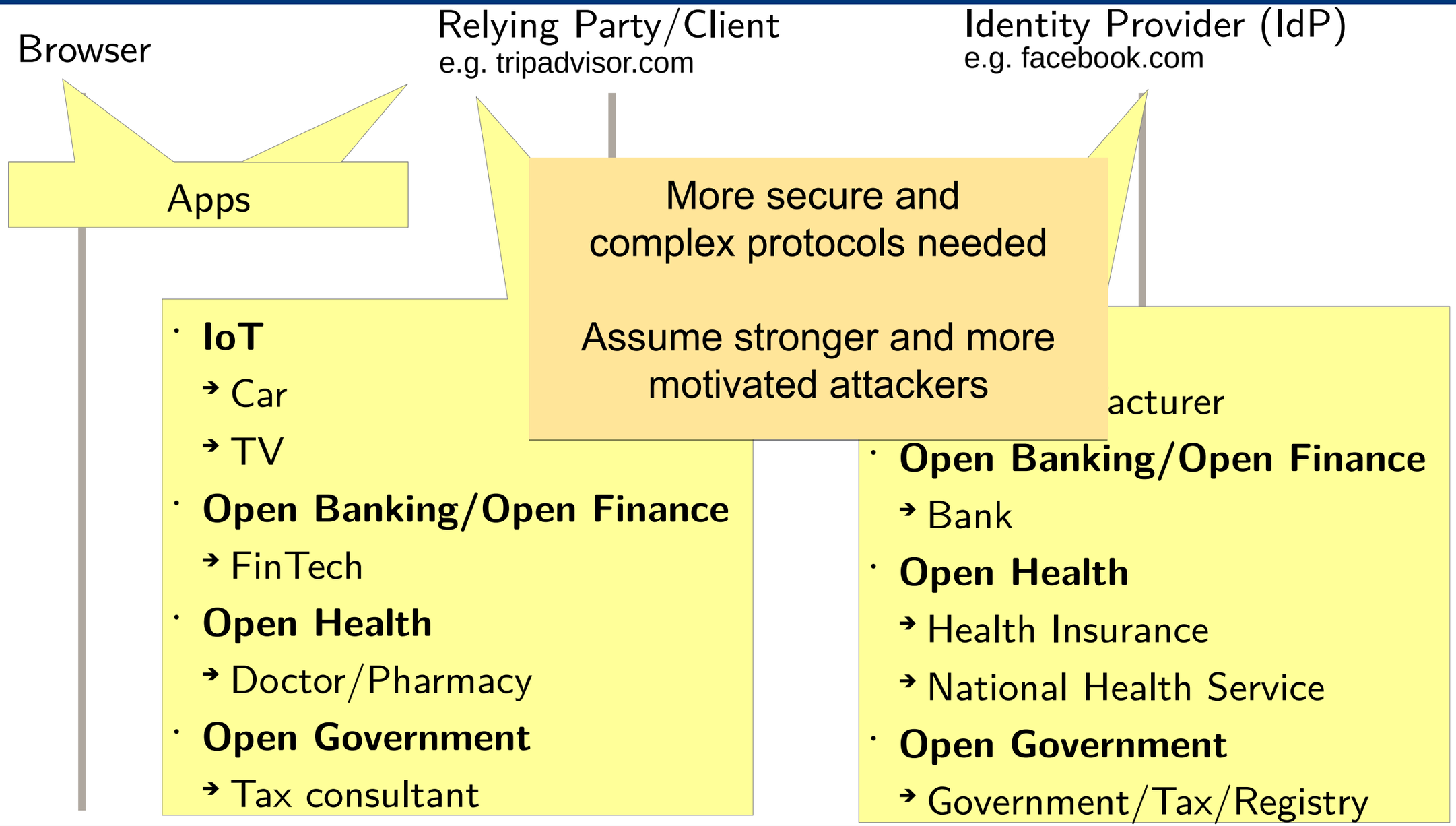
- Authentication
- Authorization
- Session Integrity

Impact

- Disclosed OAuth 2.0 attacks to the IETF Web Authorization Working Group in late 2015 (and had emergency meeting)
- Since then: In close contact with the IETF and OpenID Foundation to improve standards
- Initiated the OAuth Security Workshop (OSW) to foster the exchange between researchers, standardization groups, and industry.
  This year in its 10th edition (OSW 2025).

More Recent Case Studies:
New (High-Risk) Environments and
More Functionality/Flexibility

SSO: Today


Browser
Relying Party/Client
e.g. tripadvisor.com
Identity Provider (IdP)
e.g. facebook.com
Apps
More secure and
complex protocols needed
Assume stronger and more
motivated attackers
• IoT
➔ Car
➔ TV
• Open Banking/Open Finance
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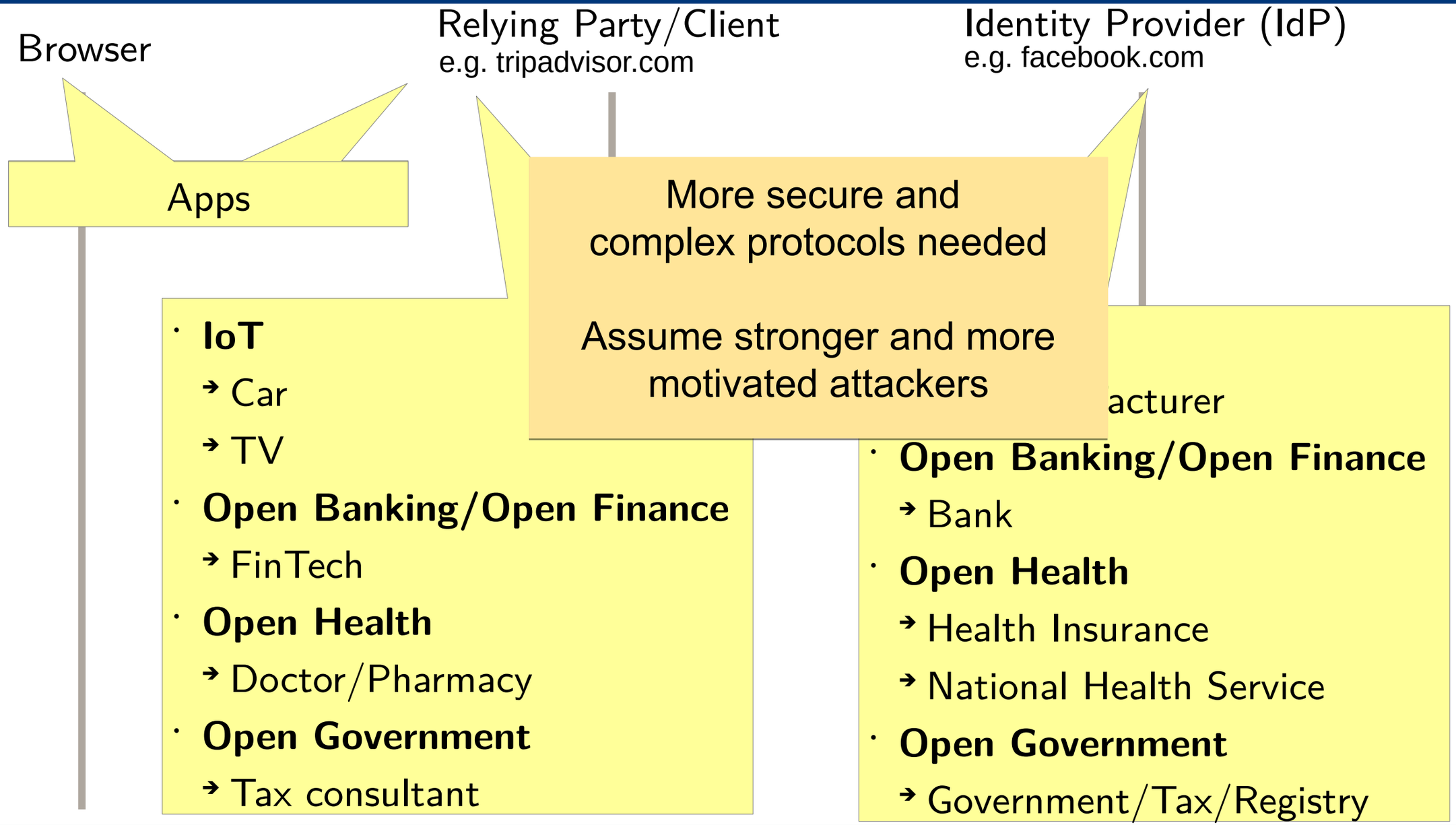
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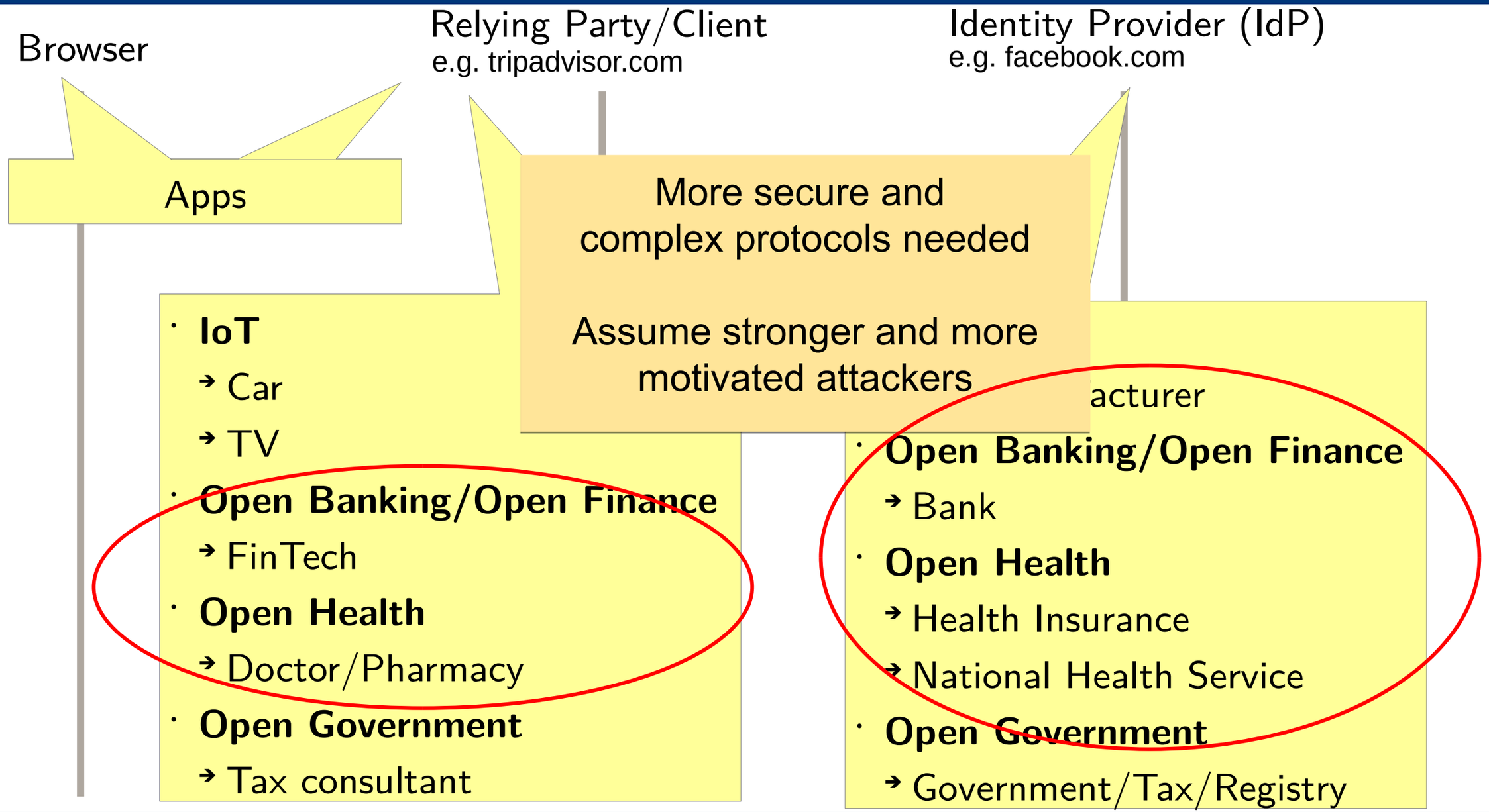
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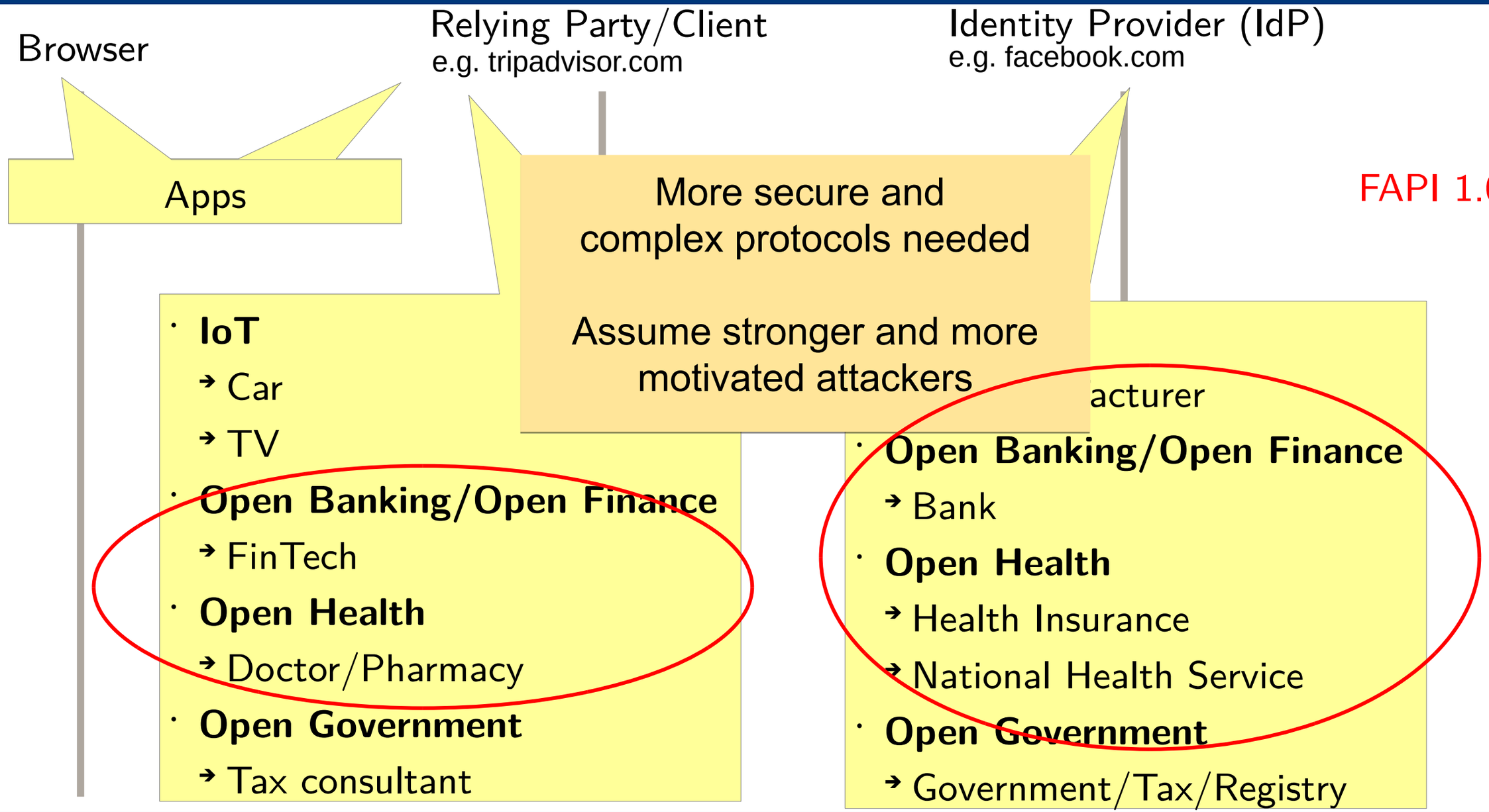
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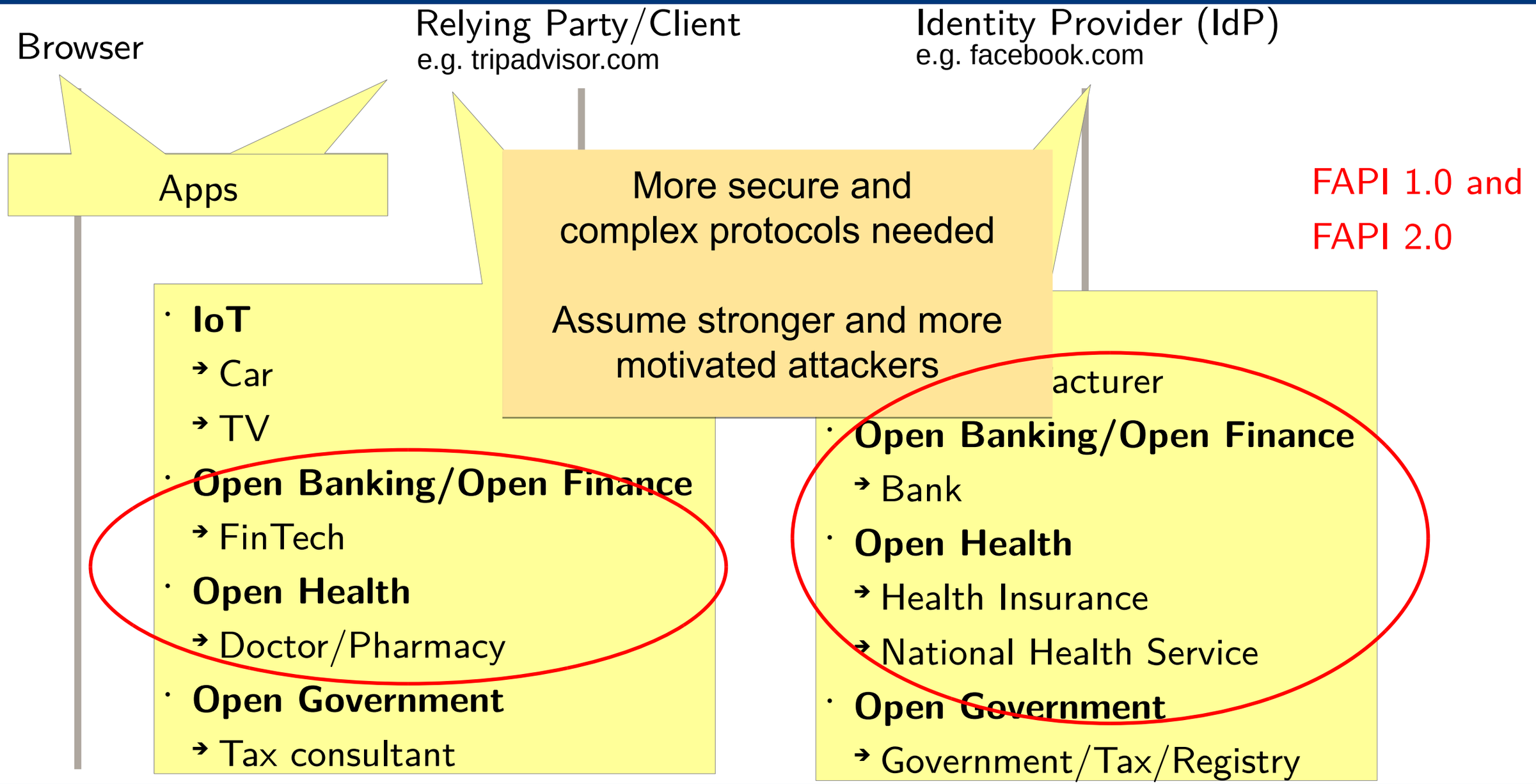
SSO: Today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Background: FAPI

- Open Banking UK
- Open Banking Brazil
- Open Insurance Brazil
- Open Finance Brazil
- Australia's Consumer Data Standards
- Open Banking Saudi Arabia
- Financial Data Exchange
- New Zealand's core payment clearing house payments.nz
- Norway's national health data sharing

OPEN BANKING

OpenInsurance

openfinance

البنك المركزي السعودي
SAMA
Saudi Central Bank

CONSUMER DATA STANDARDS

FINANCIAL™ DATA EXCHANGE

paymentsnz®

Norsk helsenett

⇒ **Many millions of users in high-risk environments**

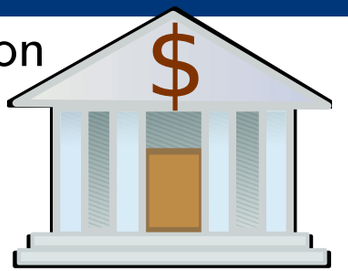
OAuth 2.0: Authorization Code Mode

OAuth 2.0: Authorization Code Mode

Browser

Client

Authorization Server



Resource Server

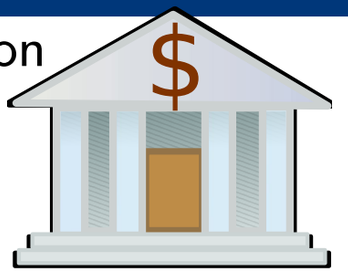
OAuth 2.0: Authorization Code Mode

Browser

Client

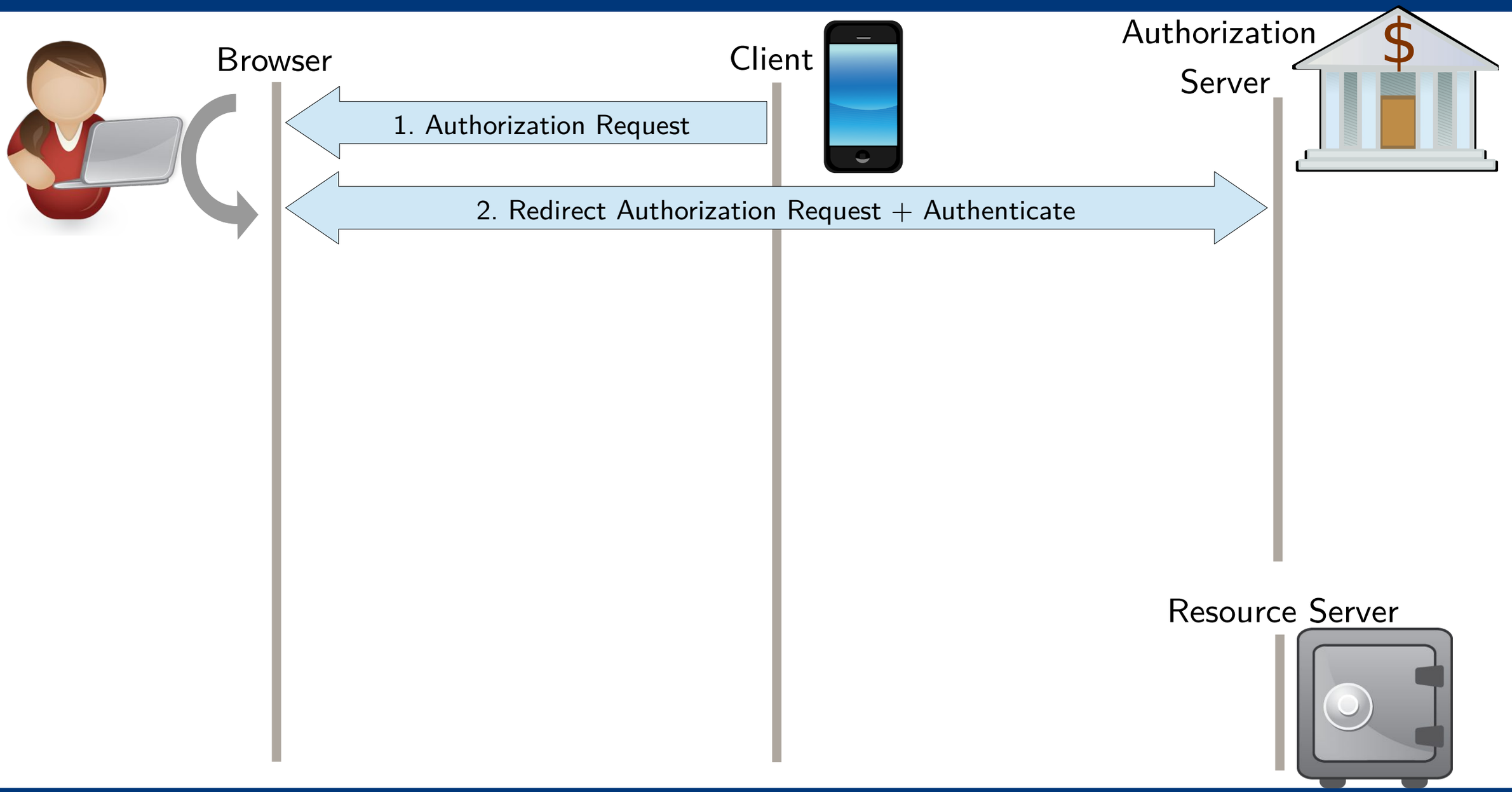
1. Authorization Request

Authorization Server

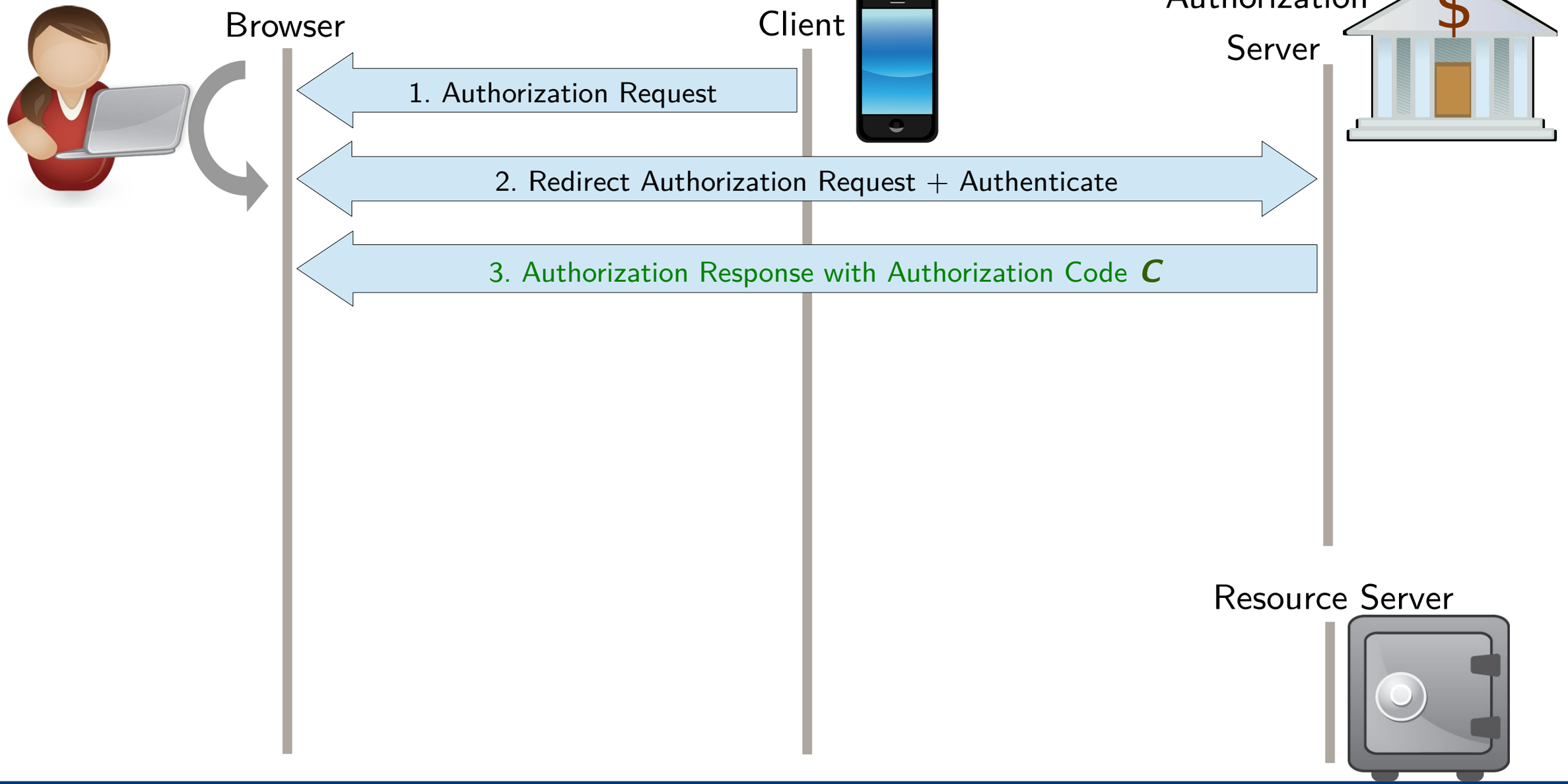


Resource Server

OAuth 2.0: Authorization Code Mode



OAuth 2.0: Authorization Code Mode



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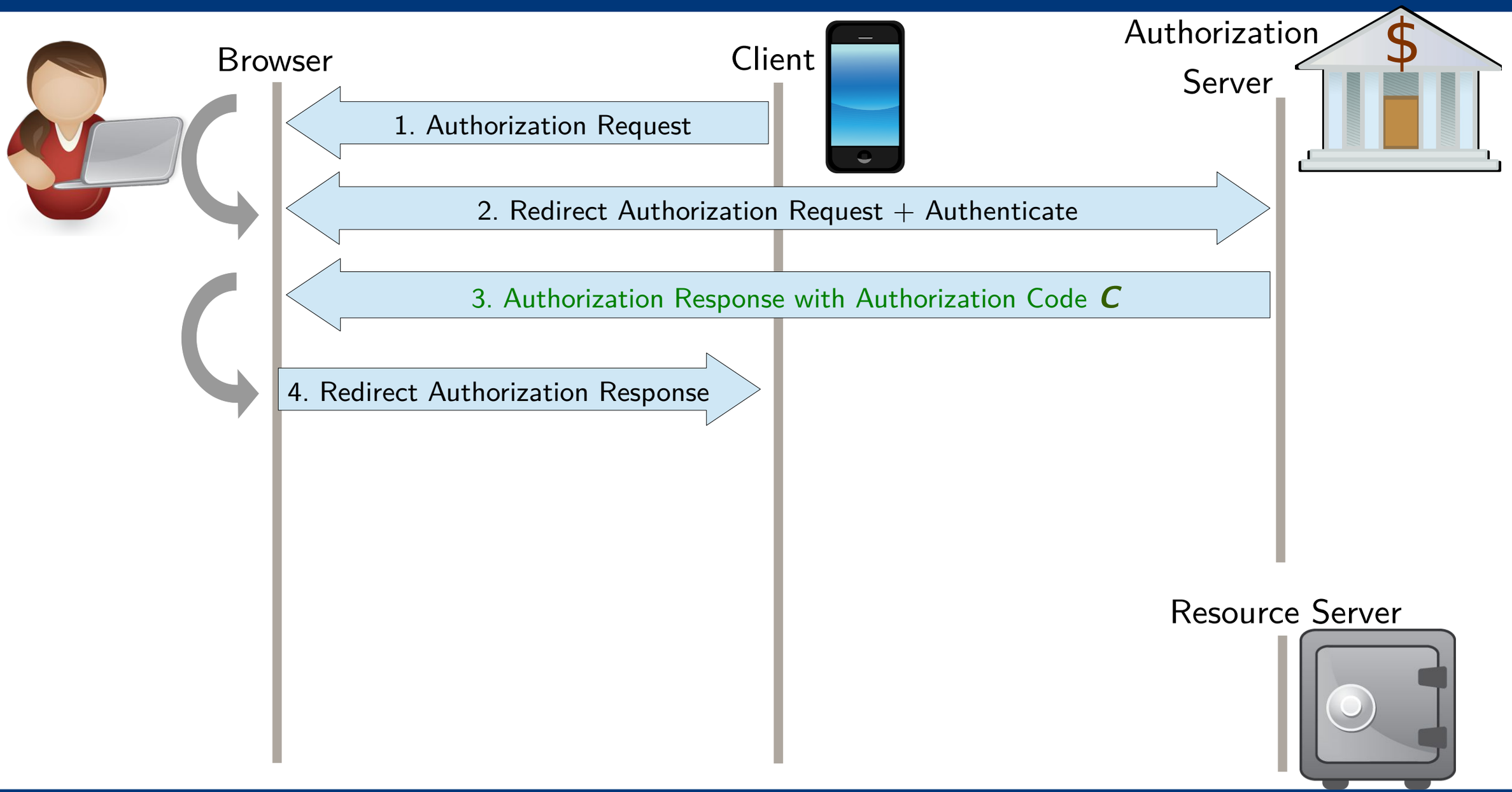

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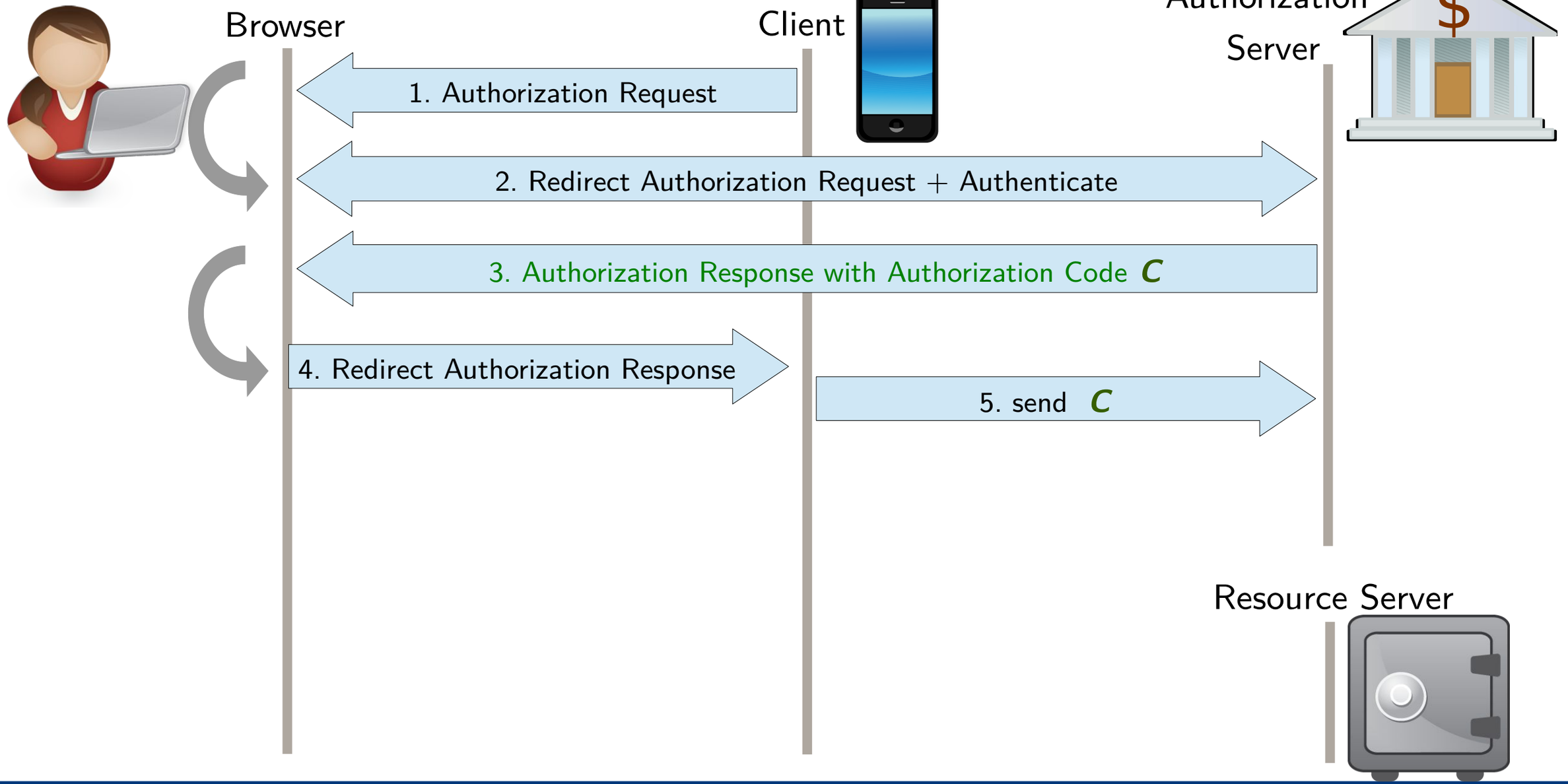
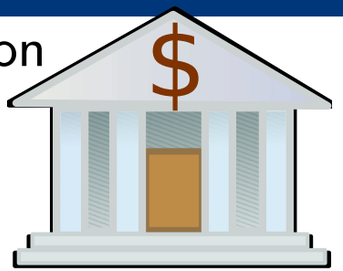


OAuth 2.0: Authorization Code Mode

Browser

Client

Authorization Server



1. Authorization Request

2. Redirect Authorization Request + Authenticate

3. Authorization Response with Authorization Code ***C***

4. Redirect Authorization Response

5. send ***C***

6. send Access Token ***AT***

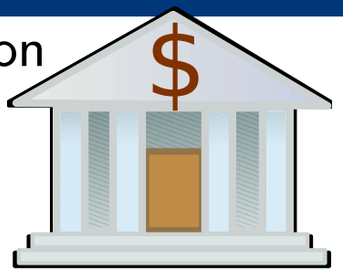
Resource Server

OAuth 2.0: Authorization Code Mode

Browser

Client

Authorization Server



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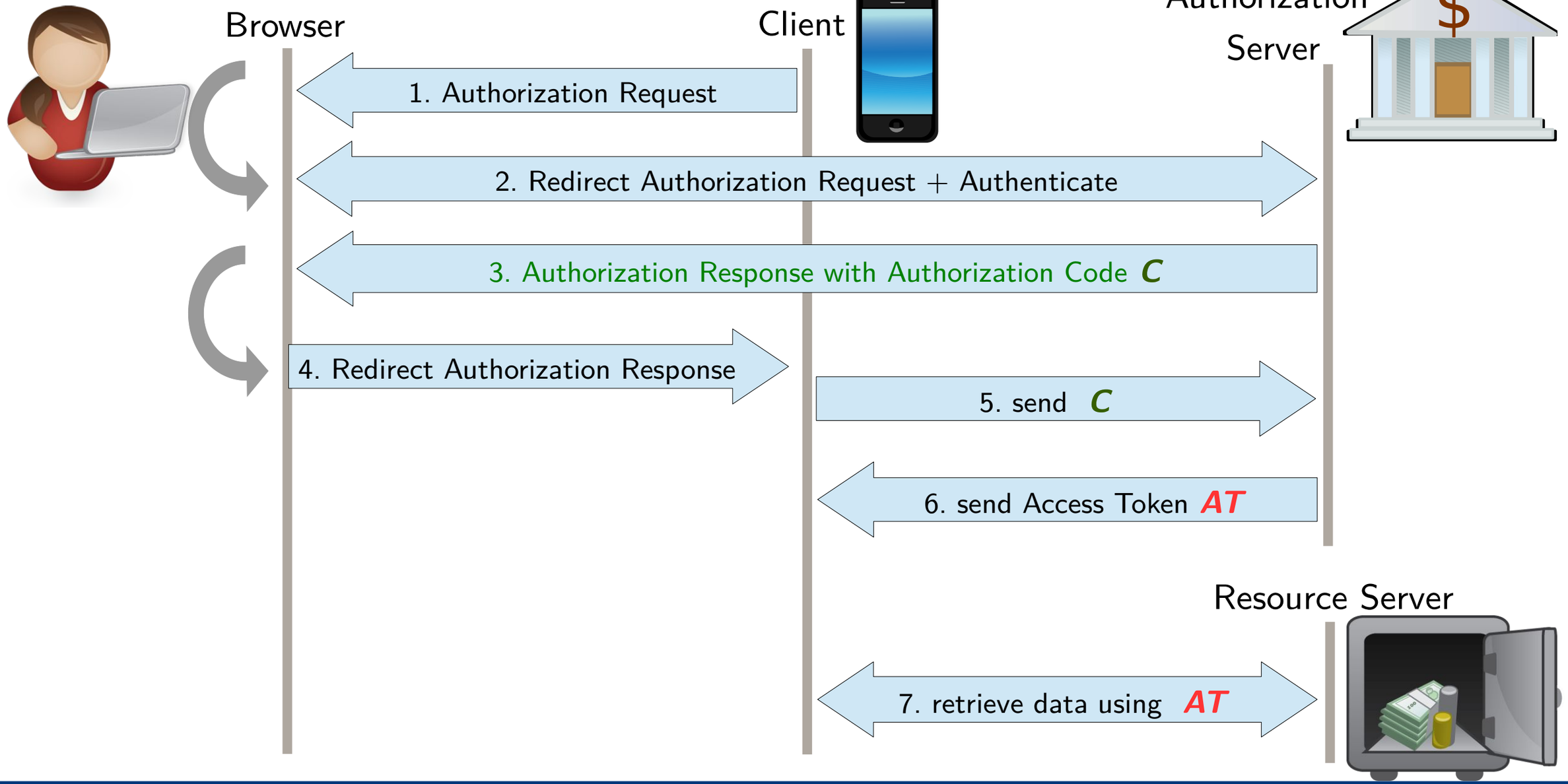
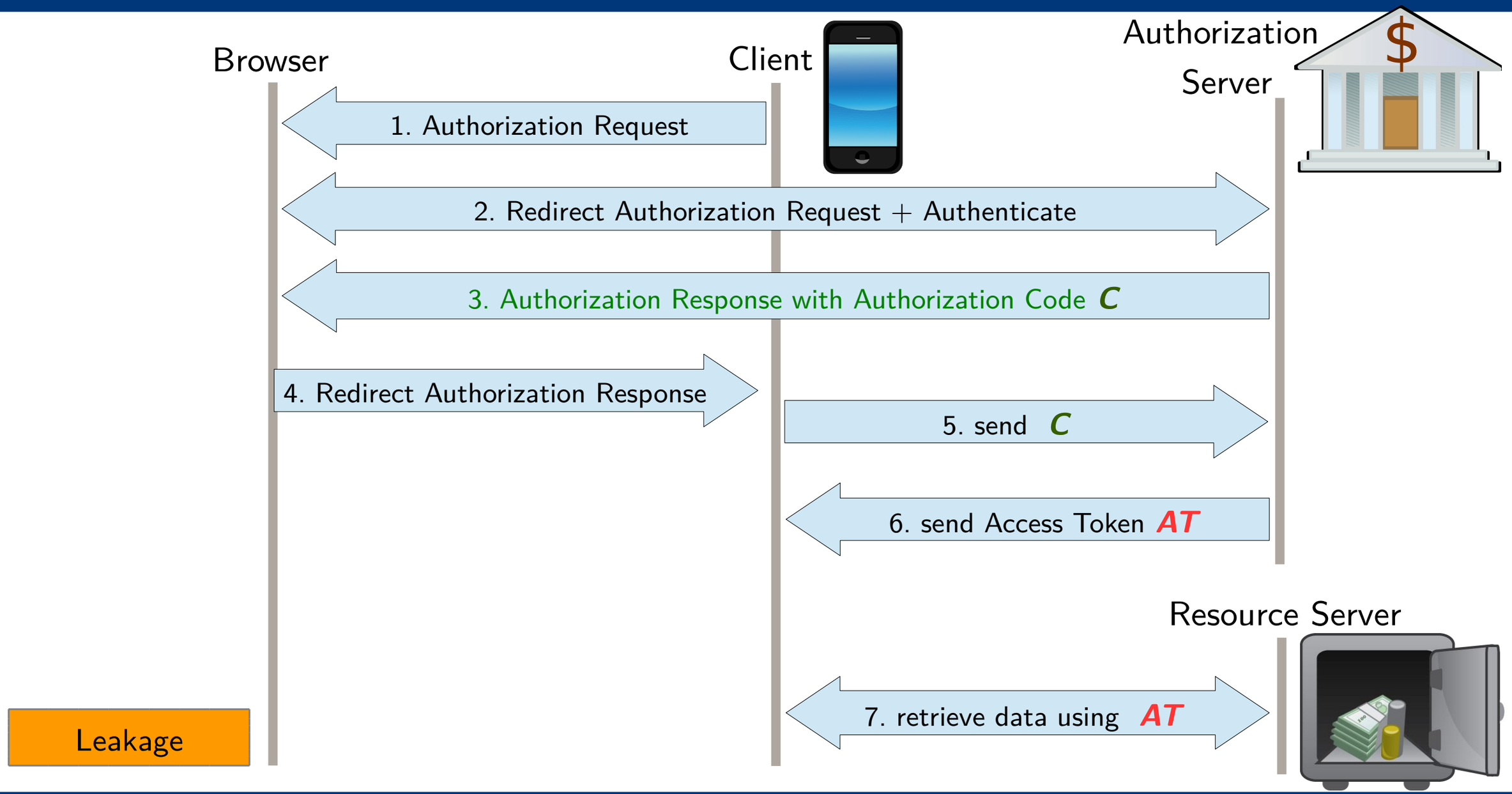
Resource Server

7. retrieve data using ***AT***

OAuth 2.0: Authorization Code Mode



FAPI: Secure, even if ...



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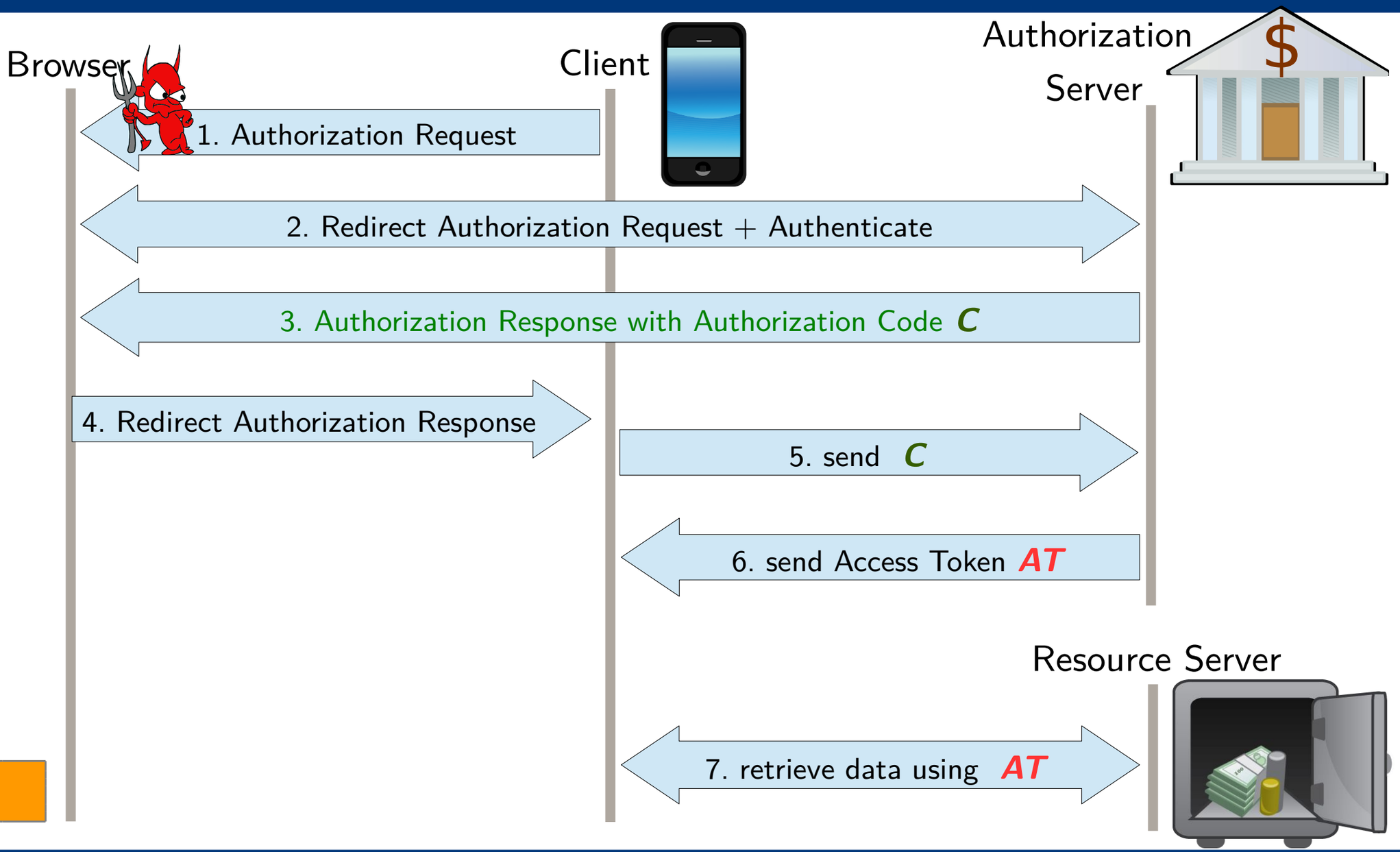

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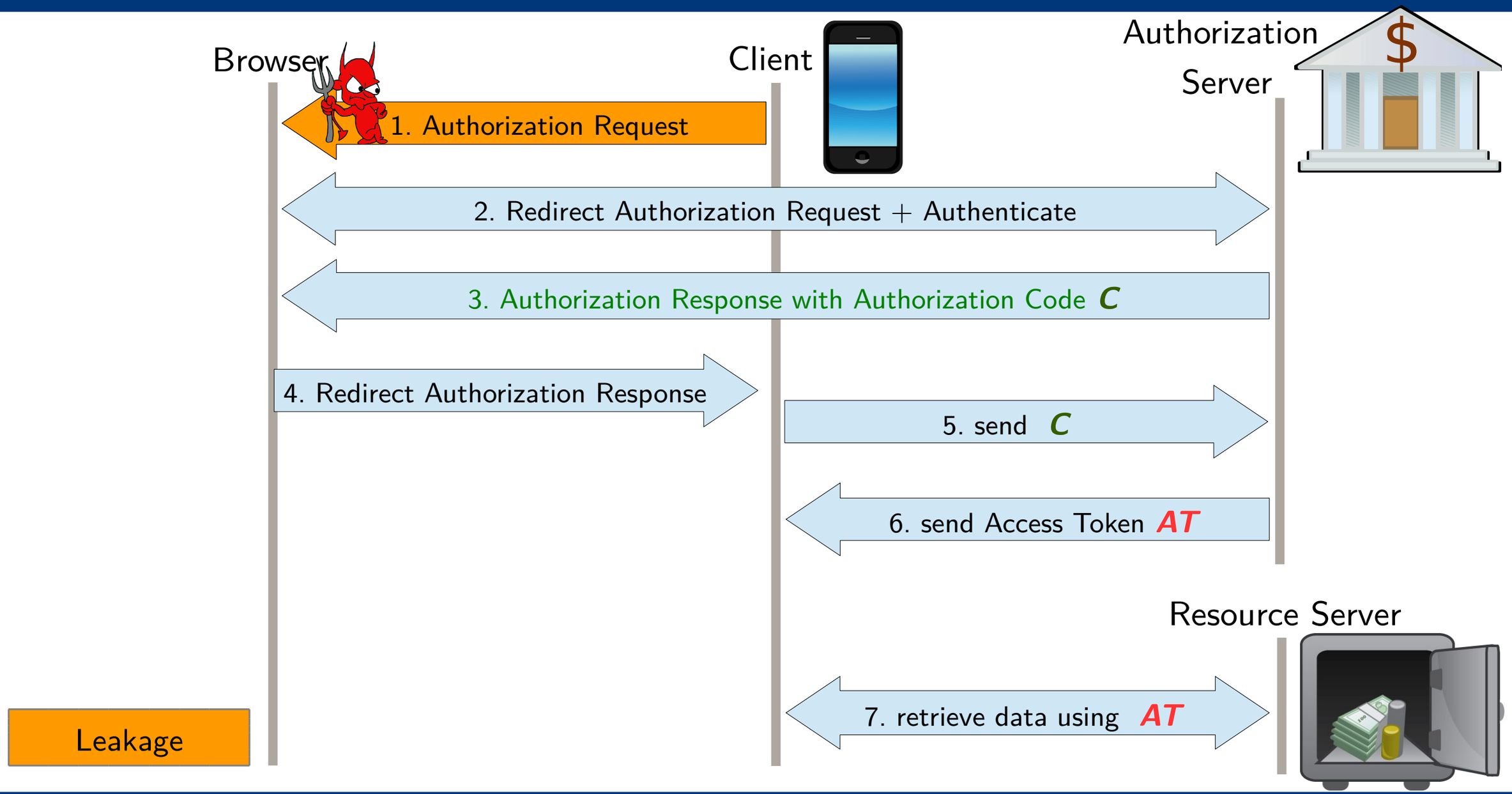

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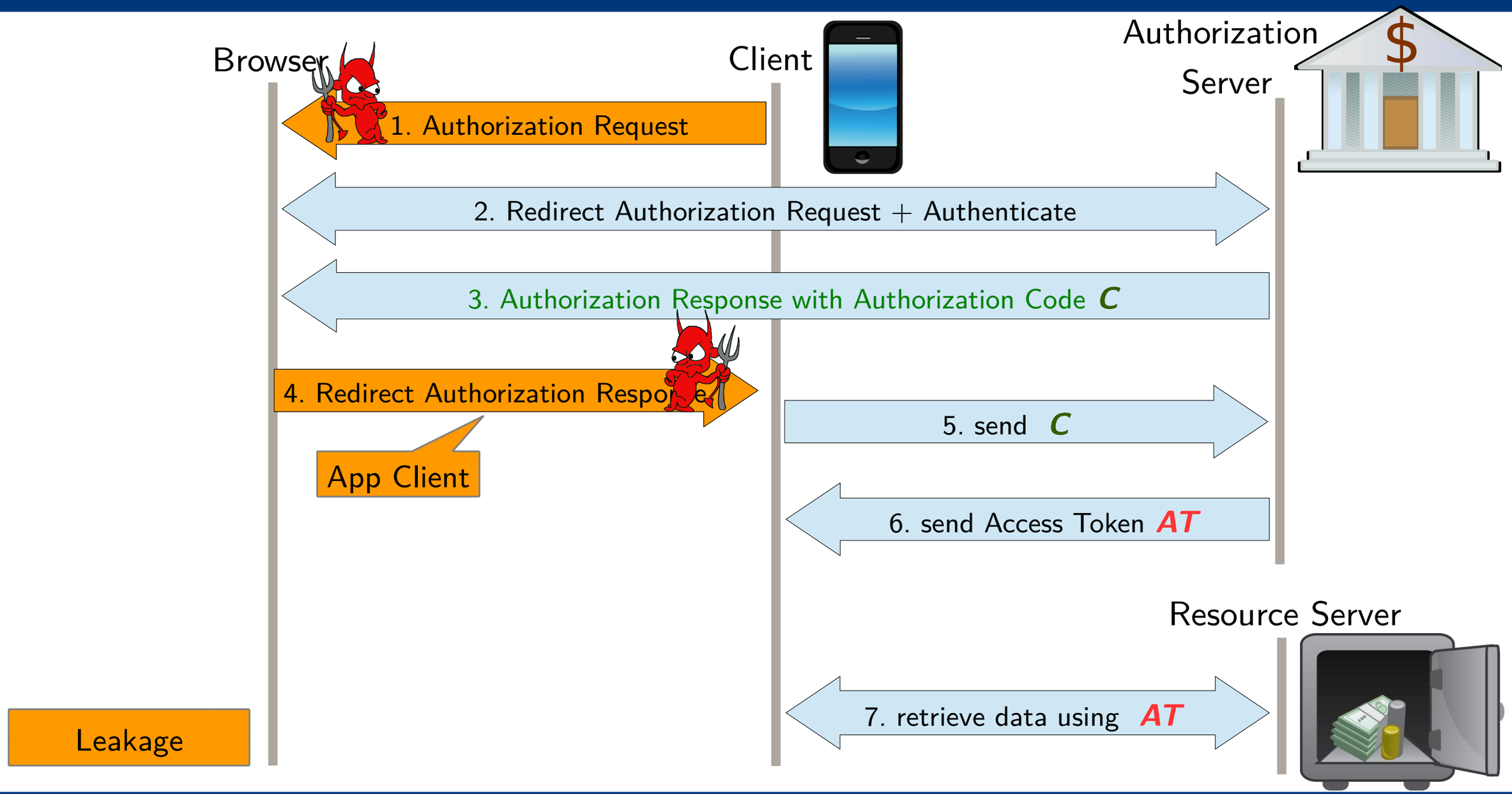

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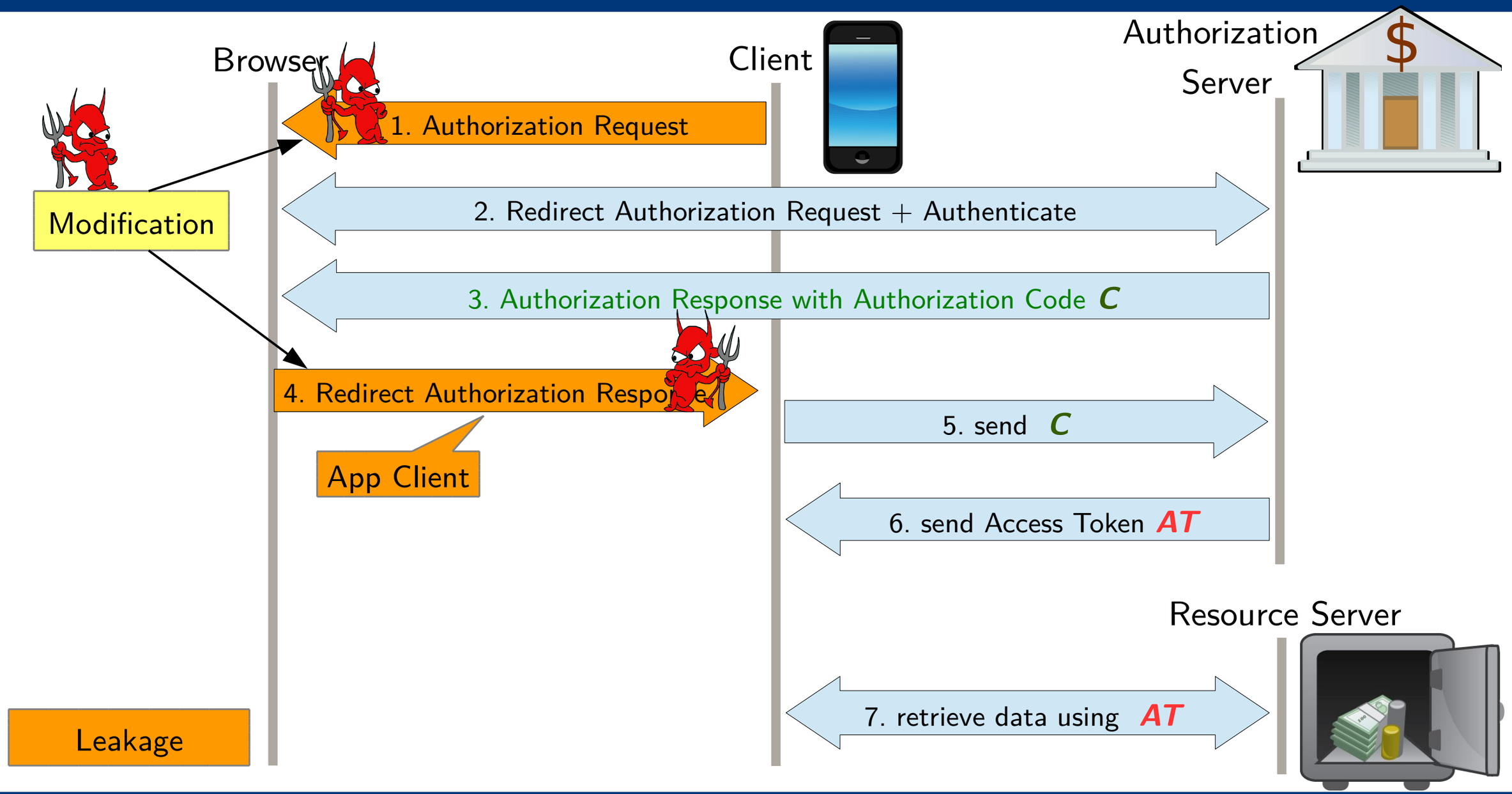

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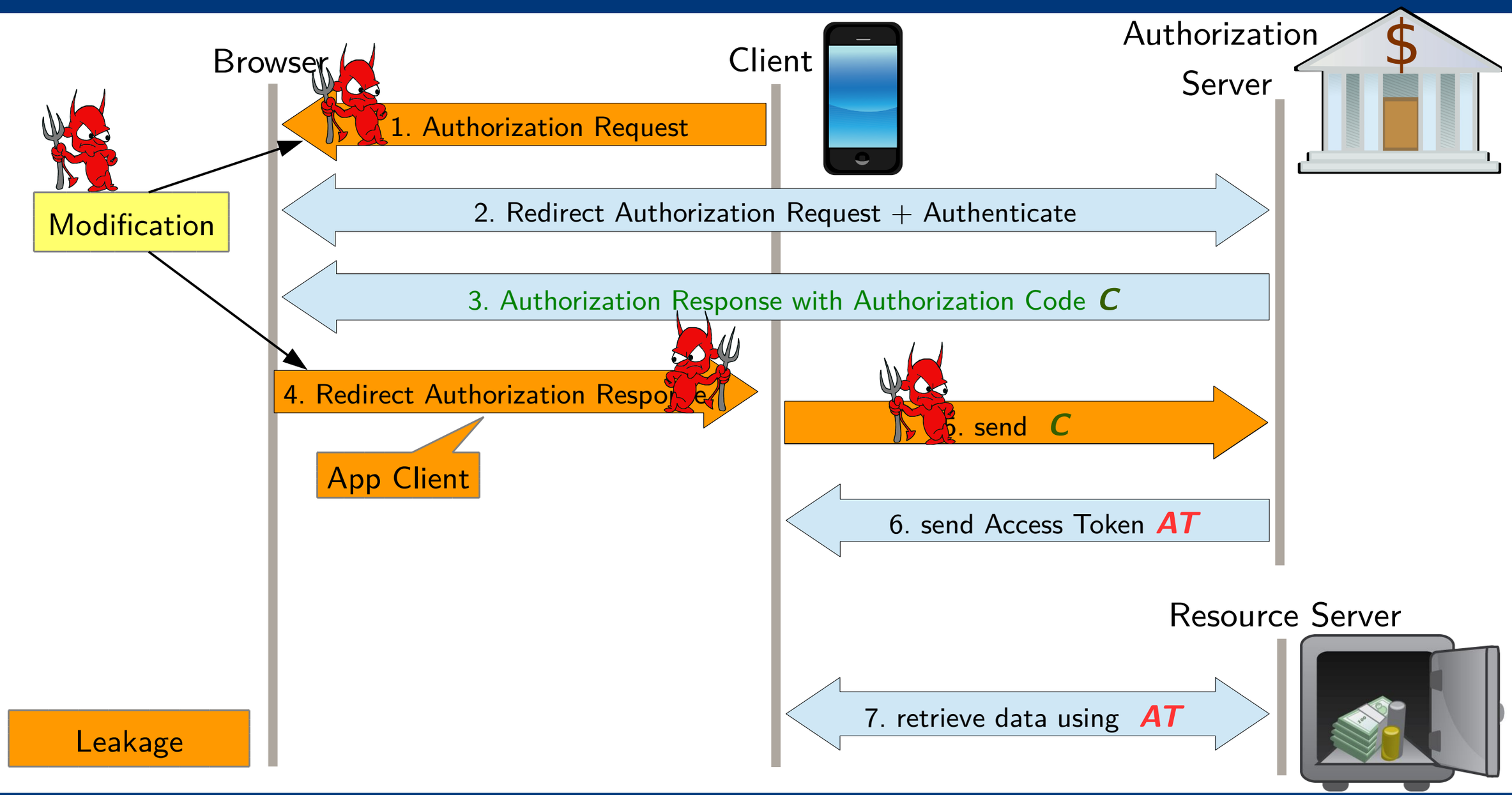

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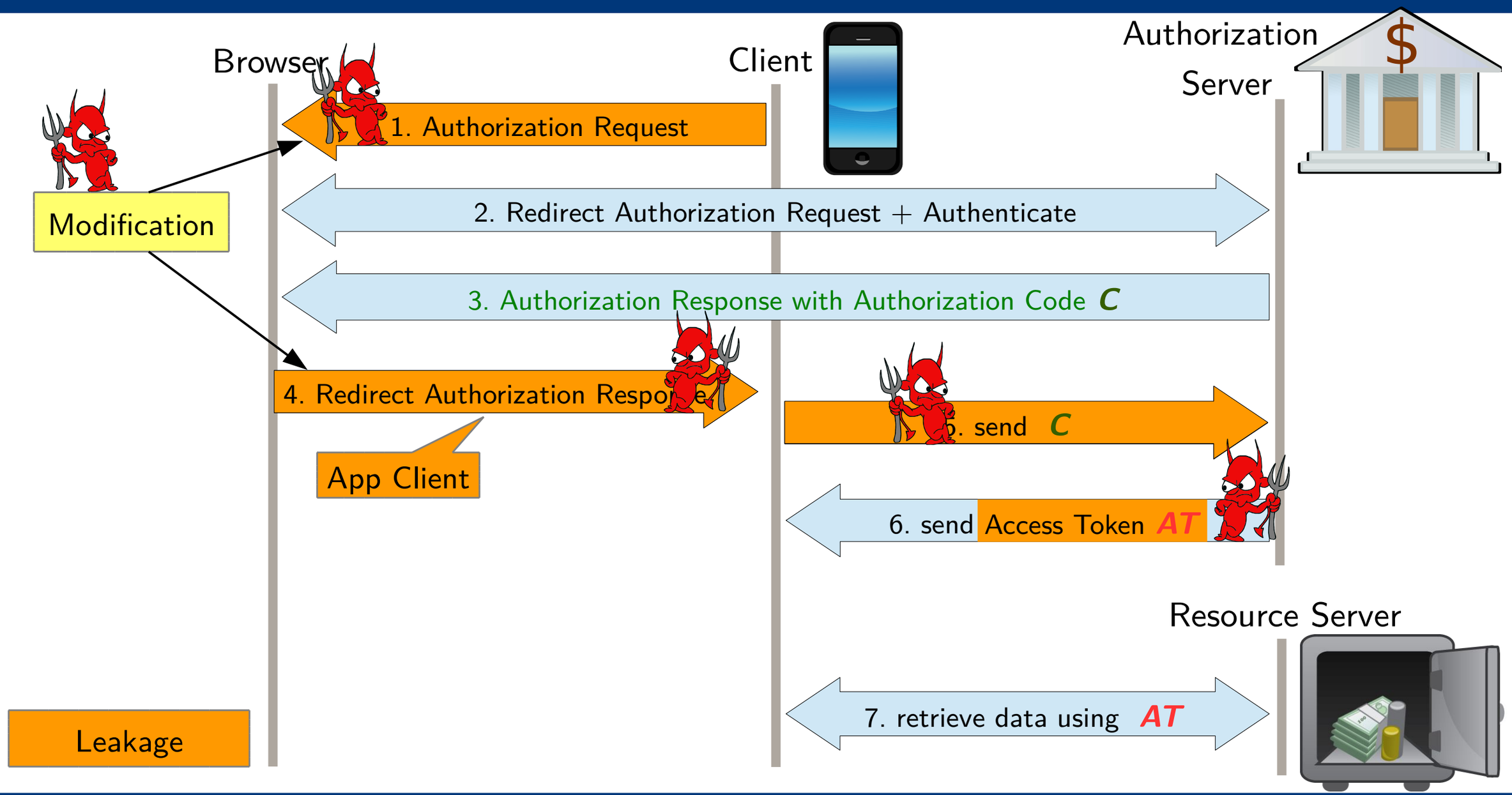
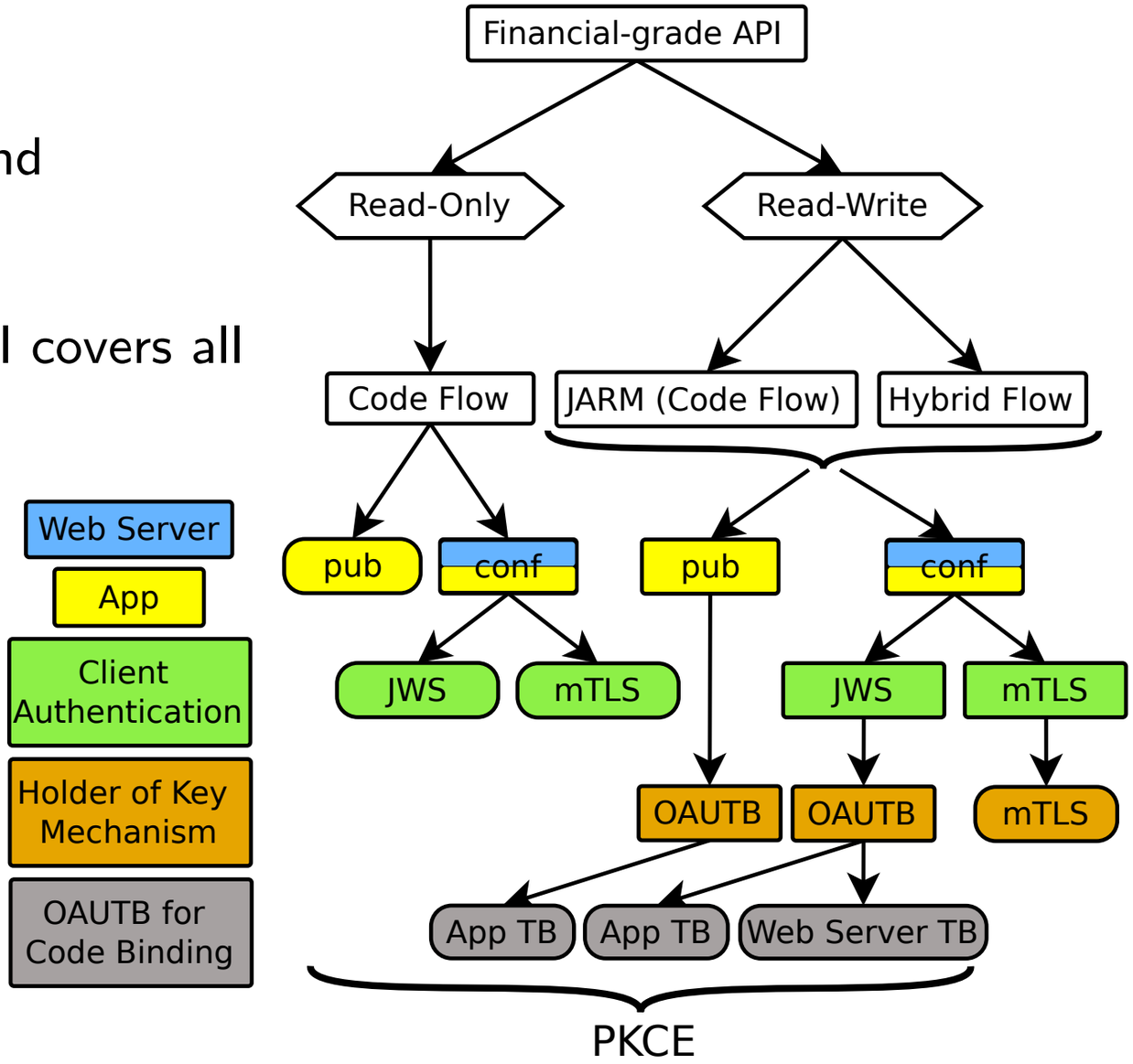
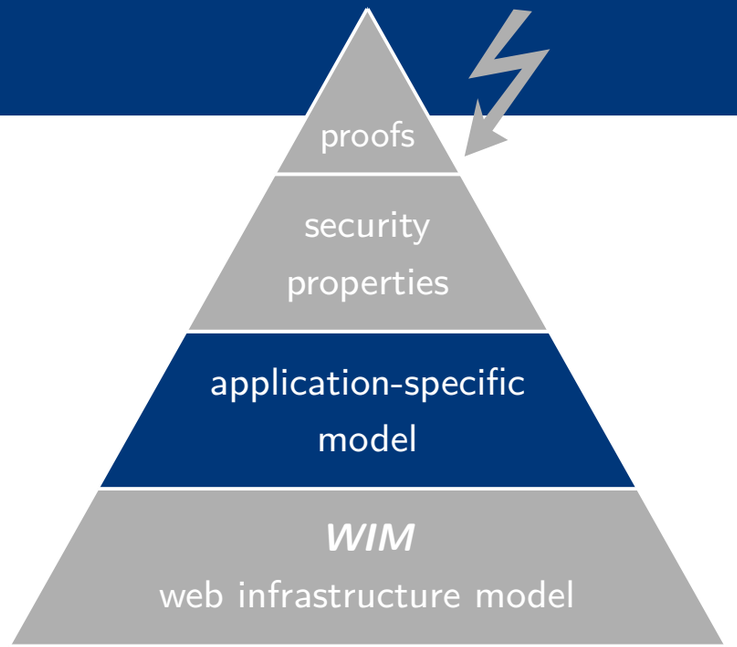


FAPI: Model

- FAPIa has
  many options and
  configurations
- Our WIM model covers all
  of them





New Defense Mechanisms

- Token Binding
- Proof Key for Code Exchange (PKCE)
- Improved Client Authentication
- Signed Authorization Request
- Signed Authorization Response (JARM)

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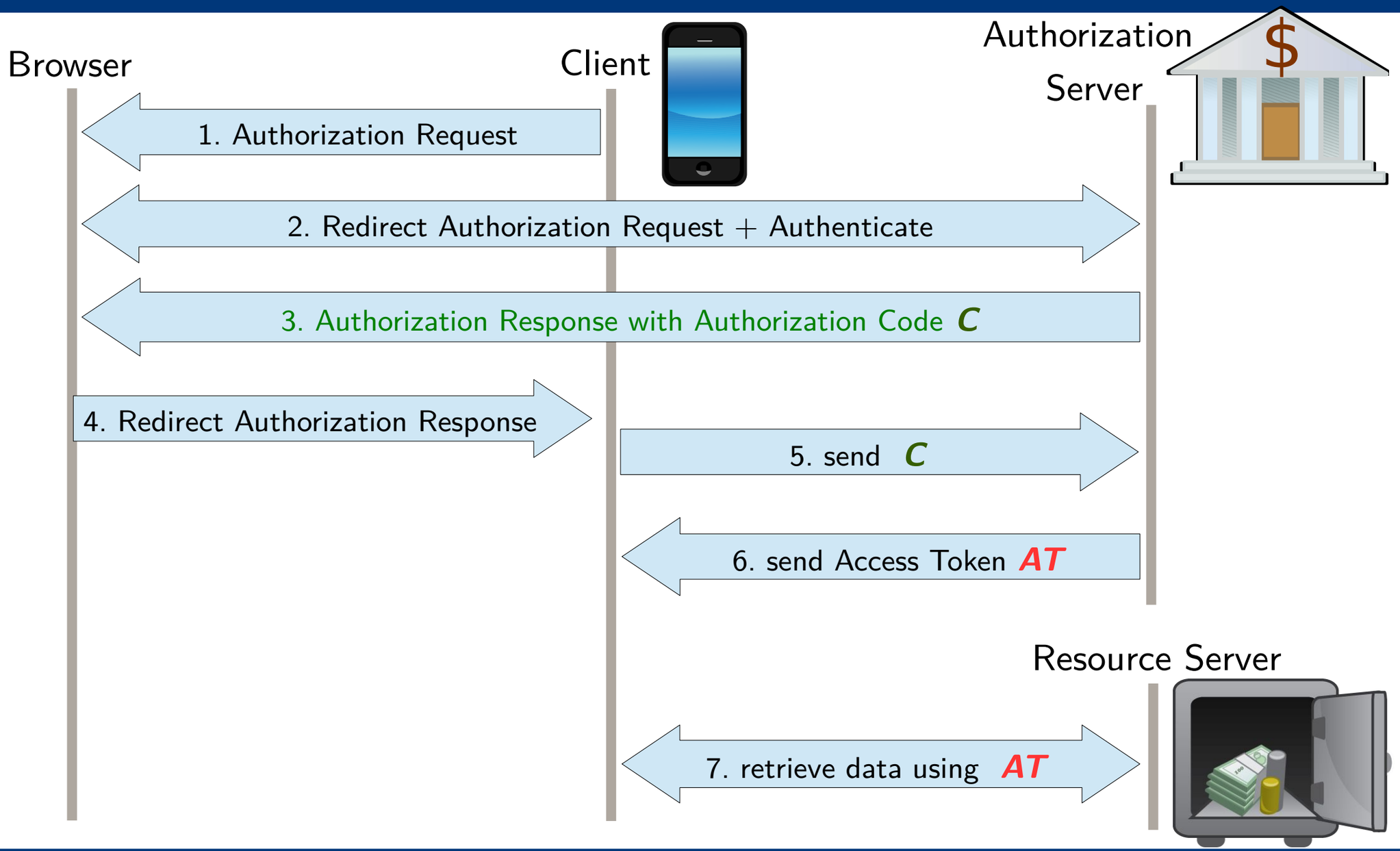
Example: Binding Access Tokens



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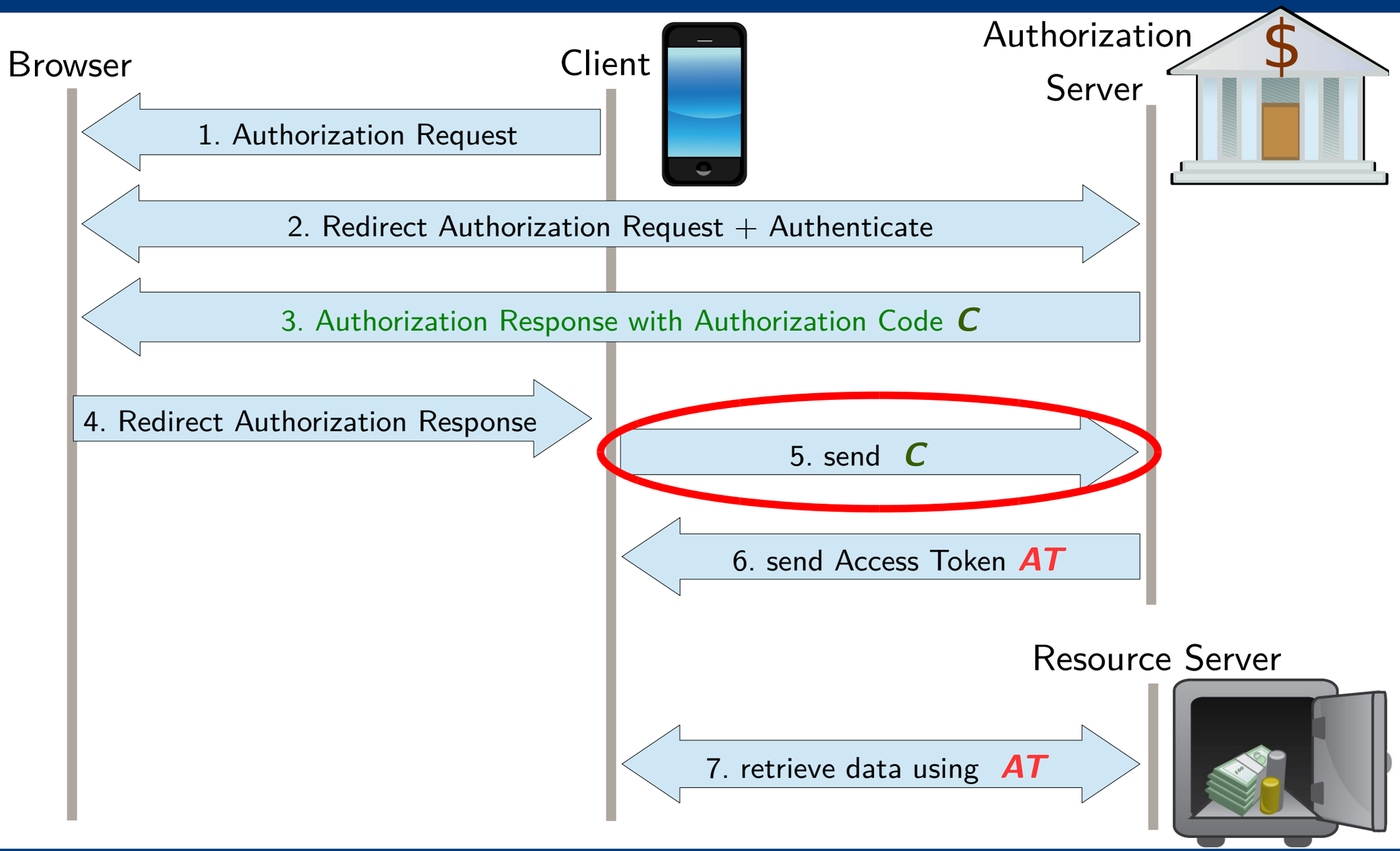

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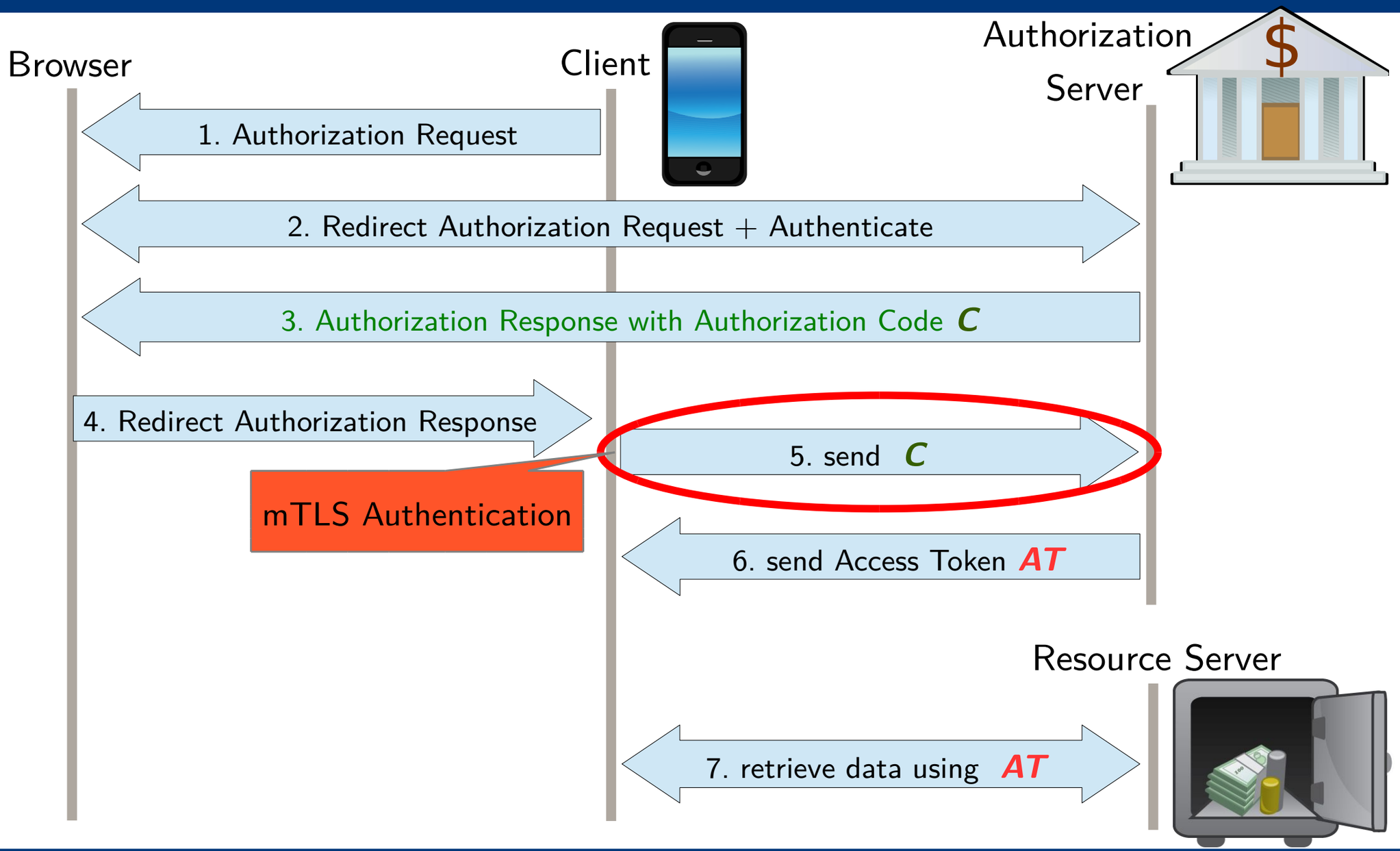

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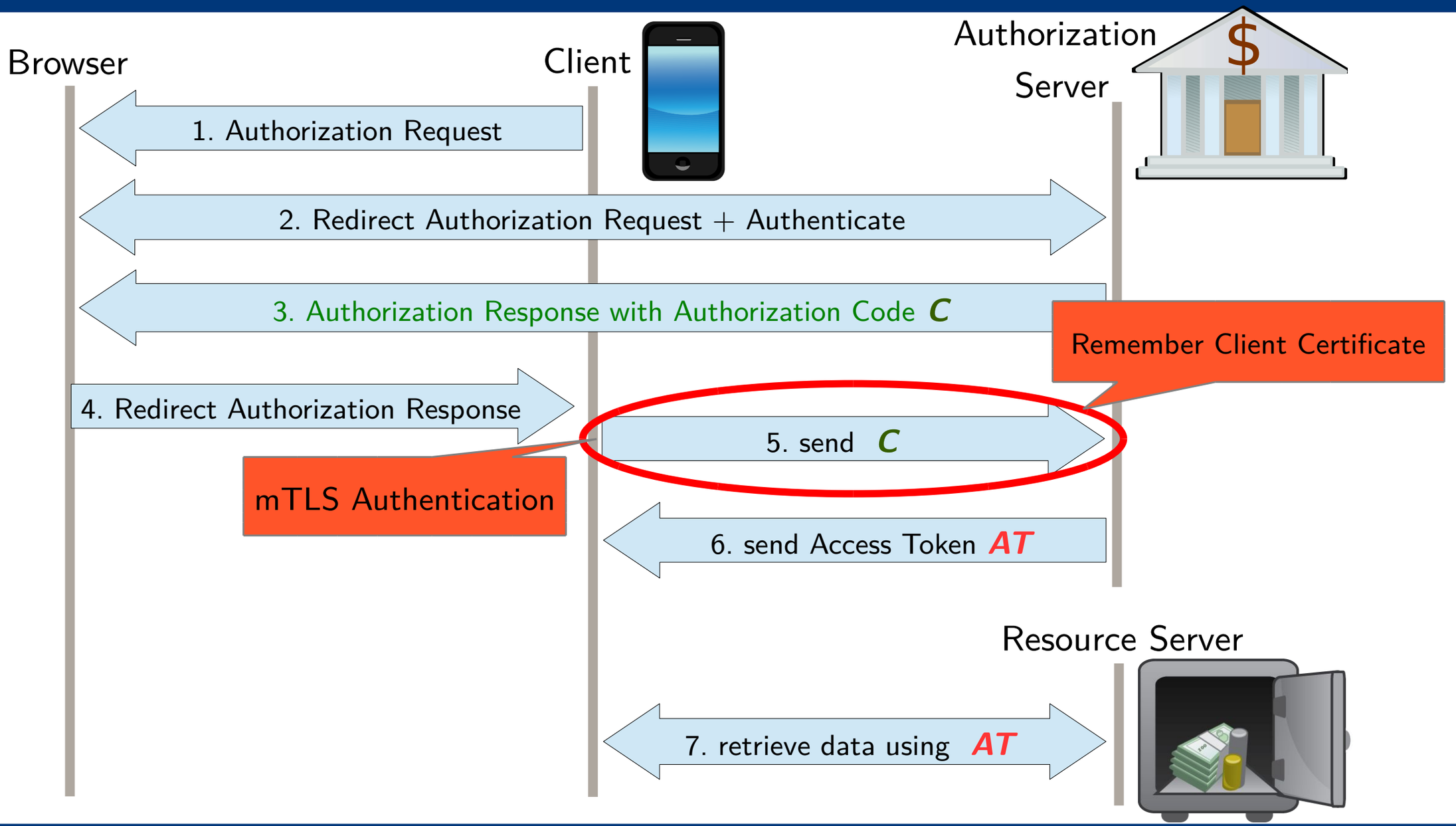

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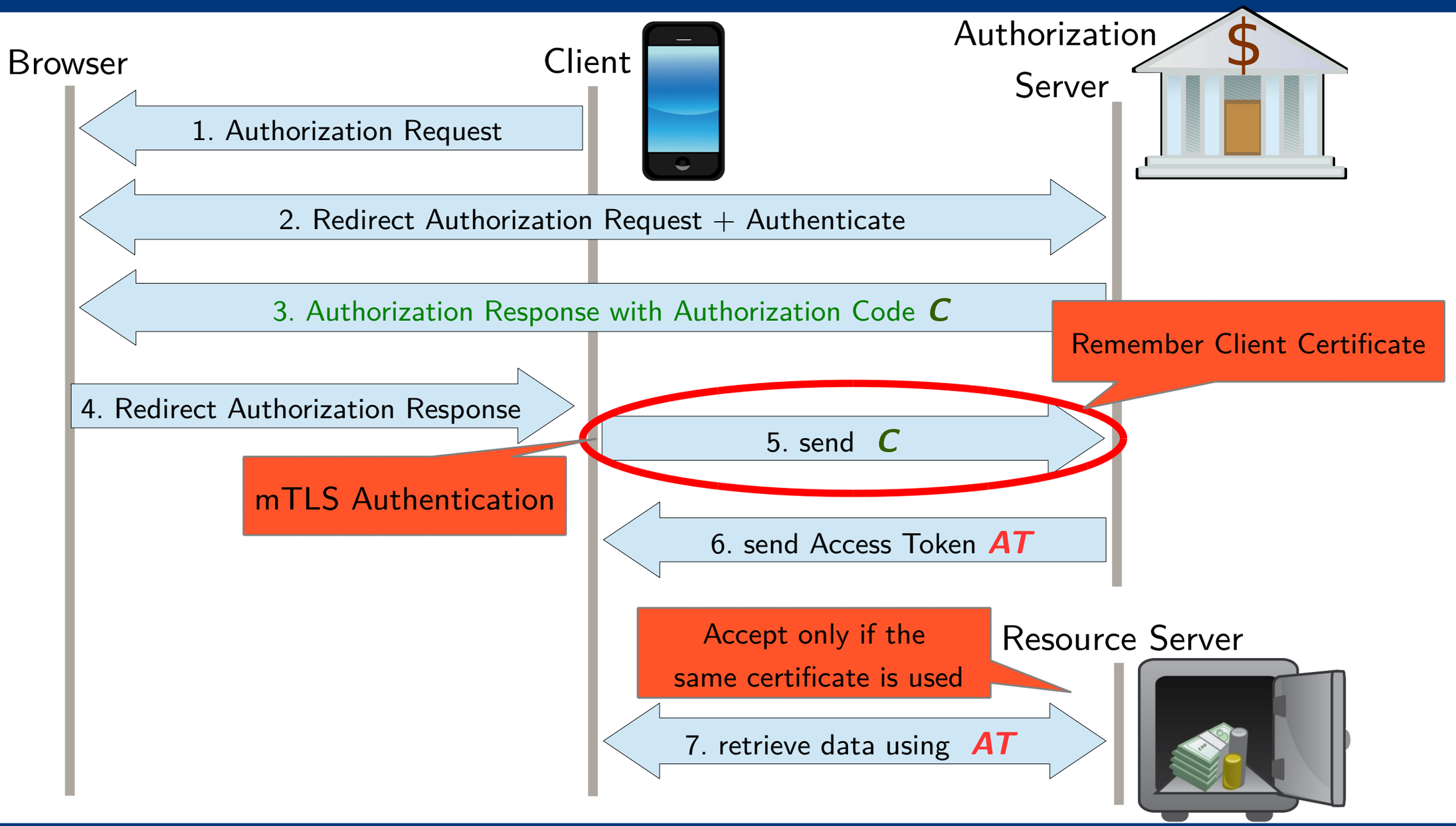

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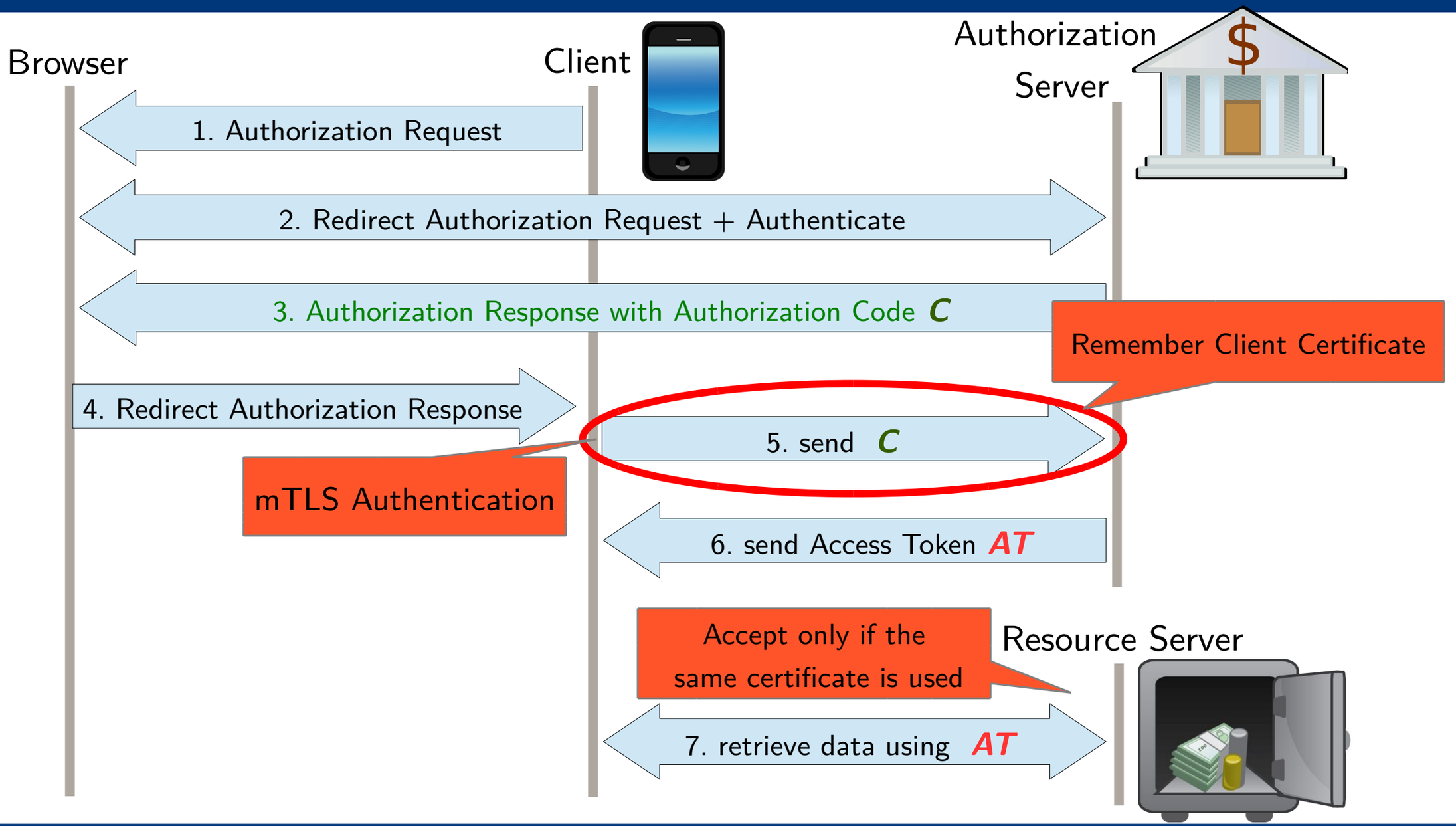

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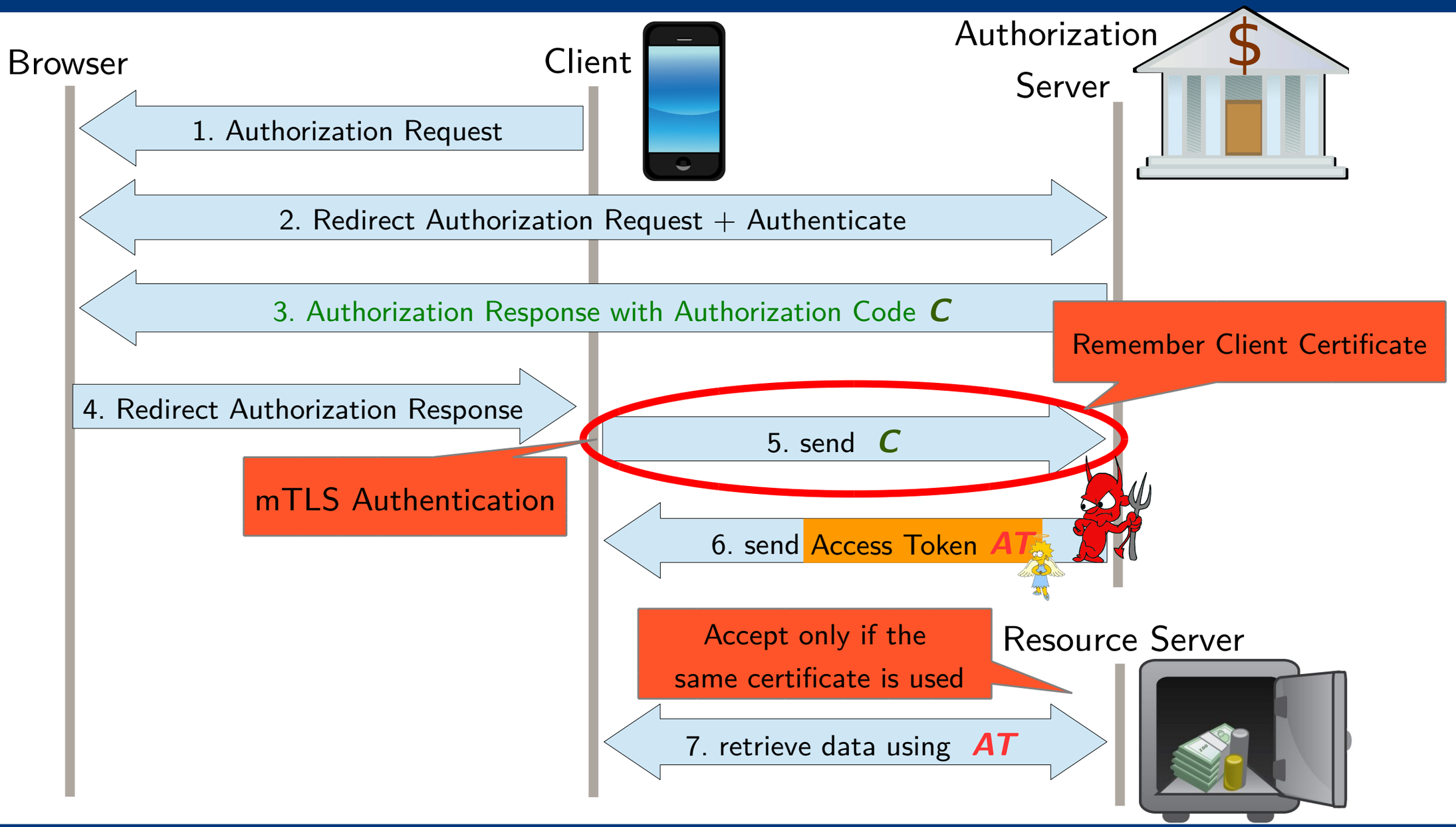


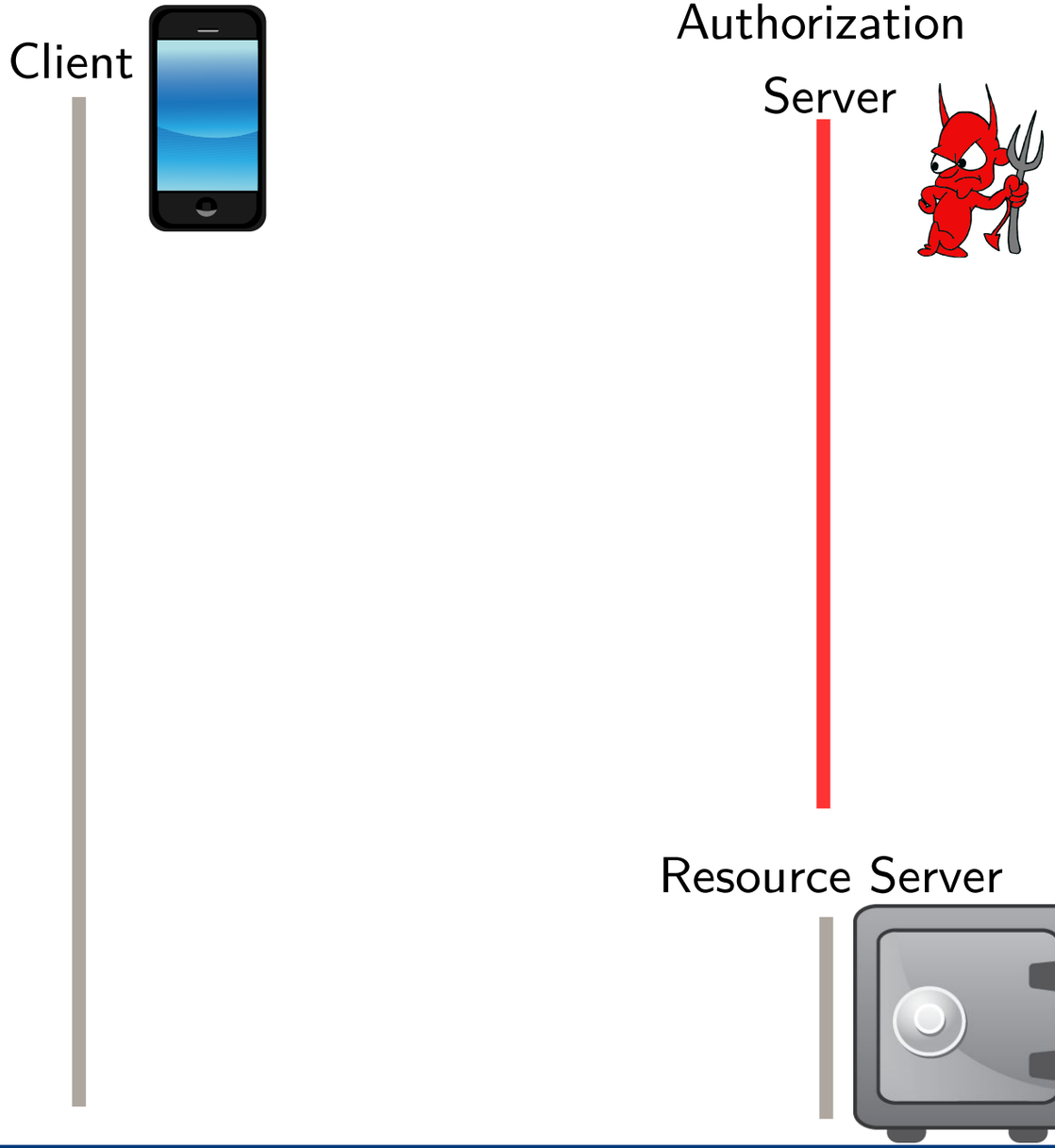
Cuckoo's Token Attack

Cuckoo's Token Attack

Client

Resource Server

Cuckoo's Token Attack



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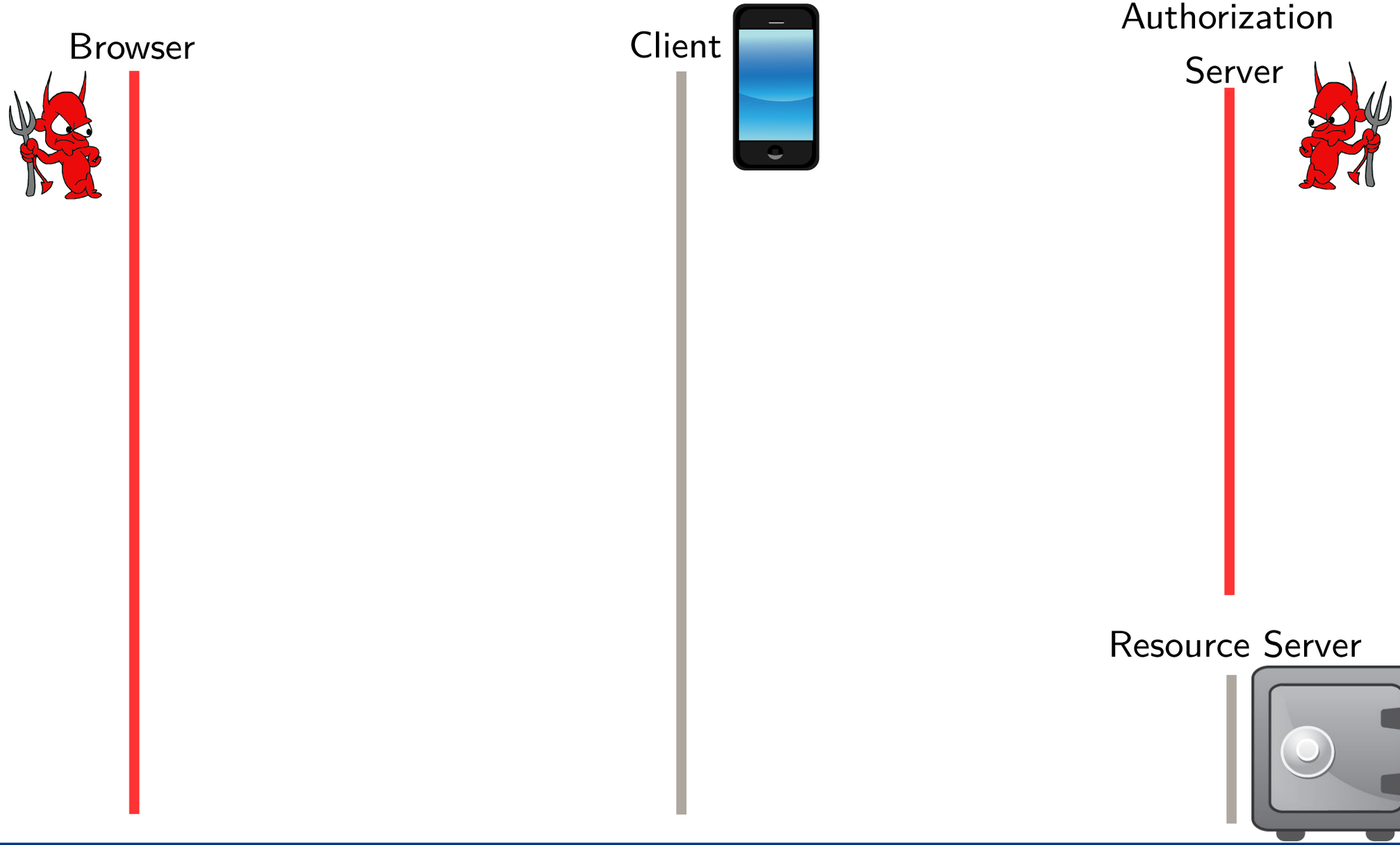

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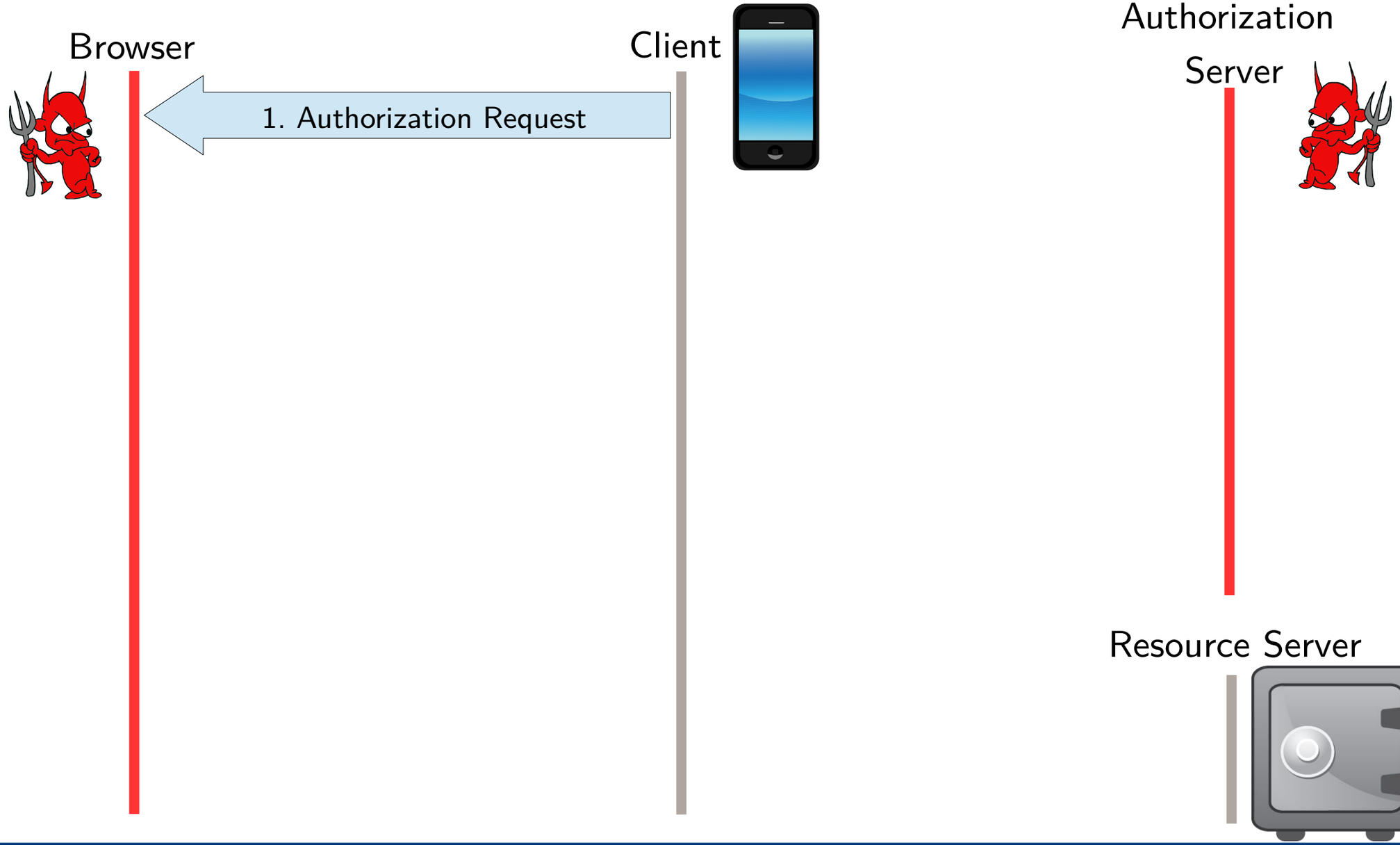

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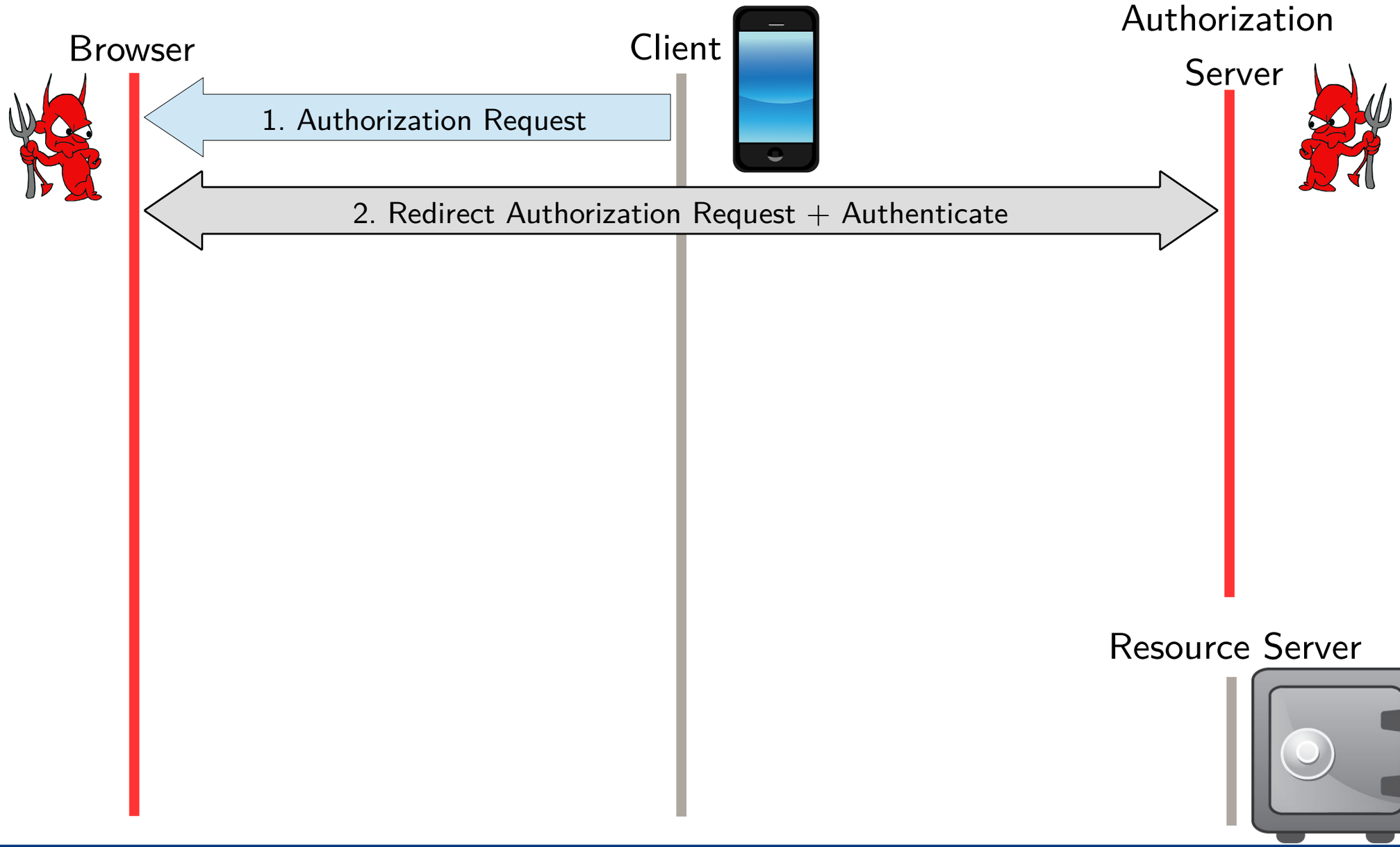

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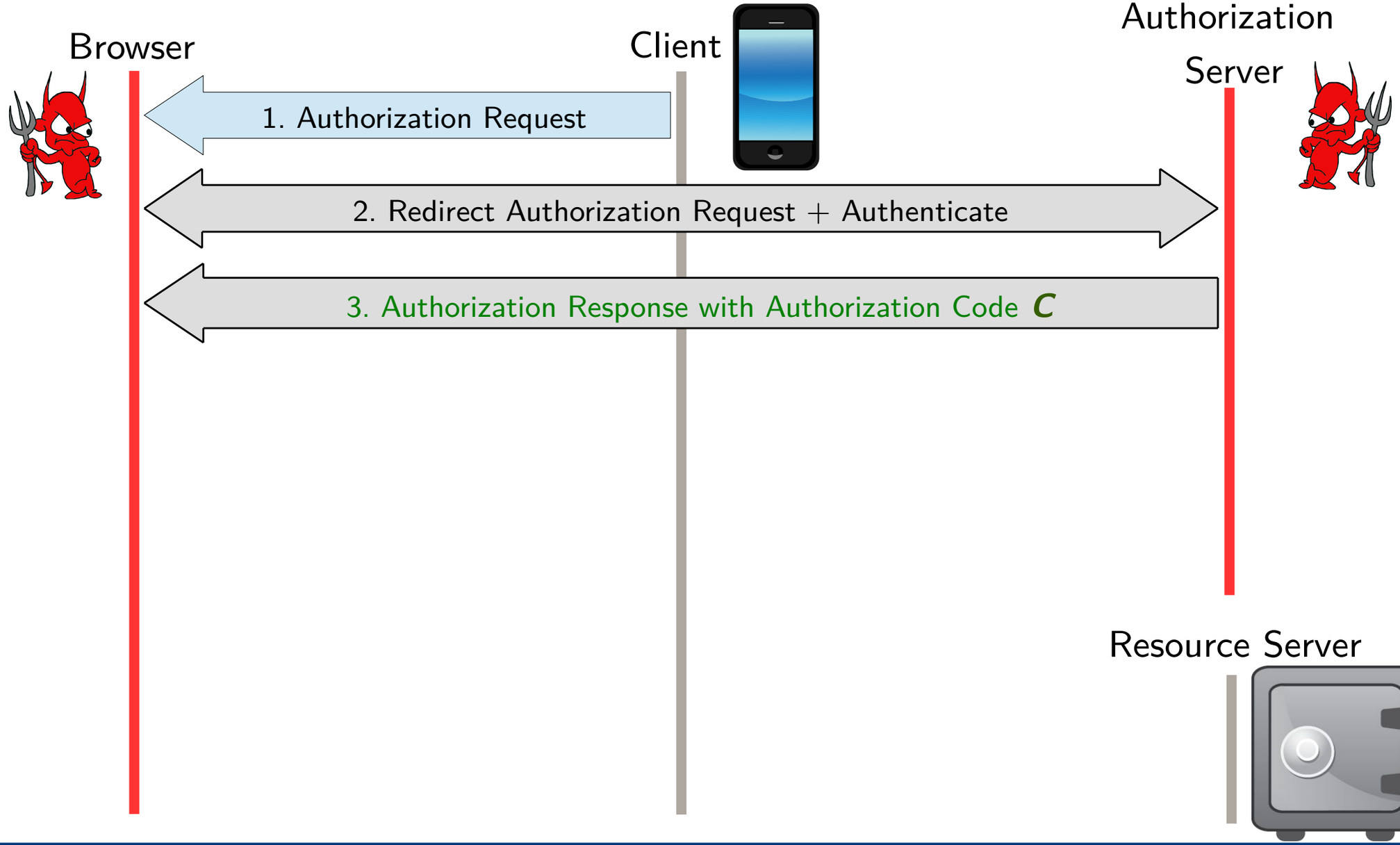

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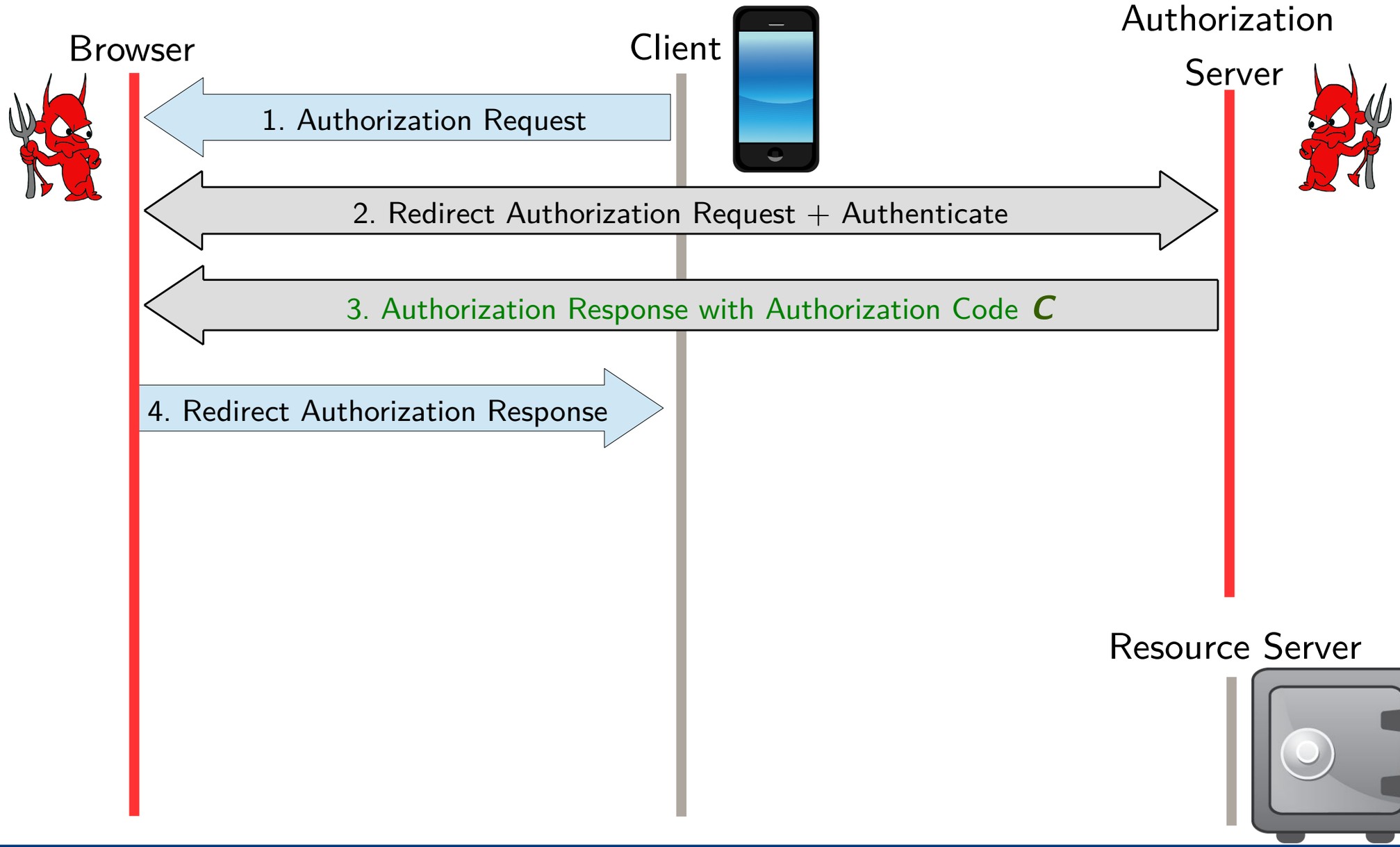

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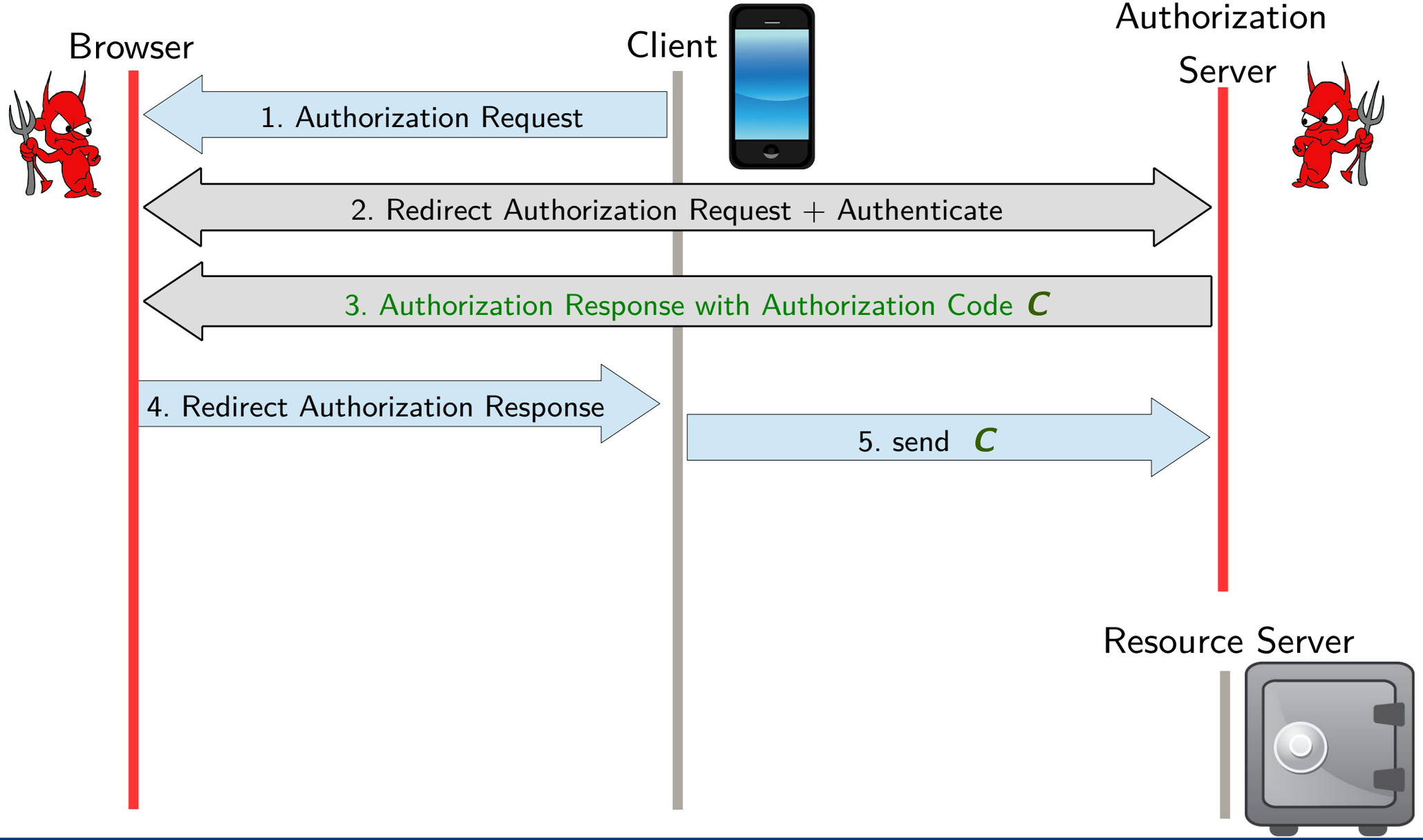

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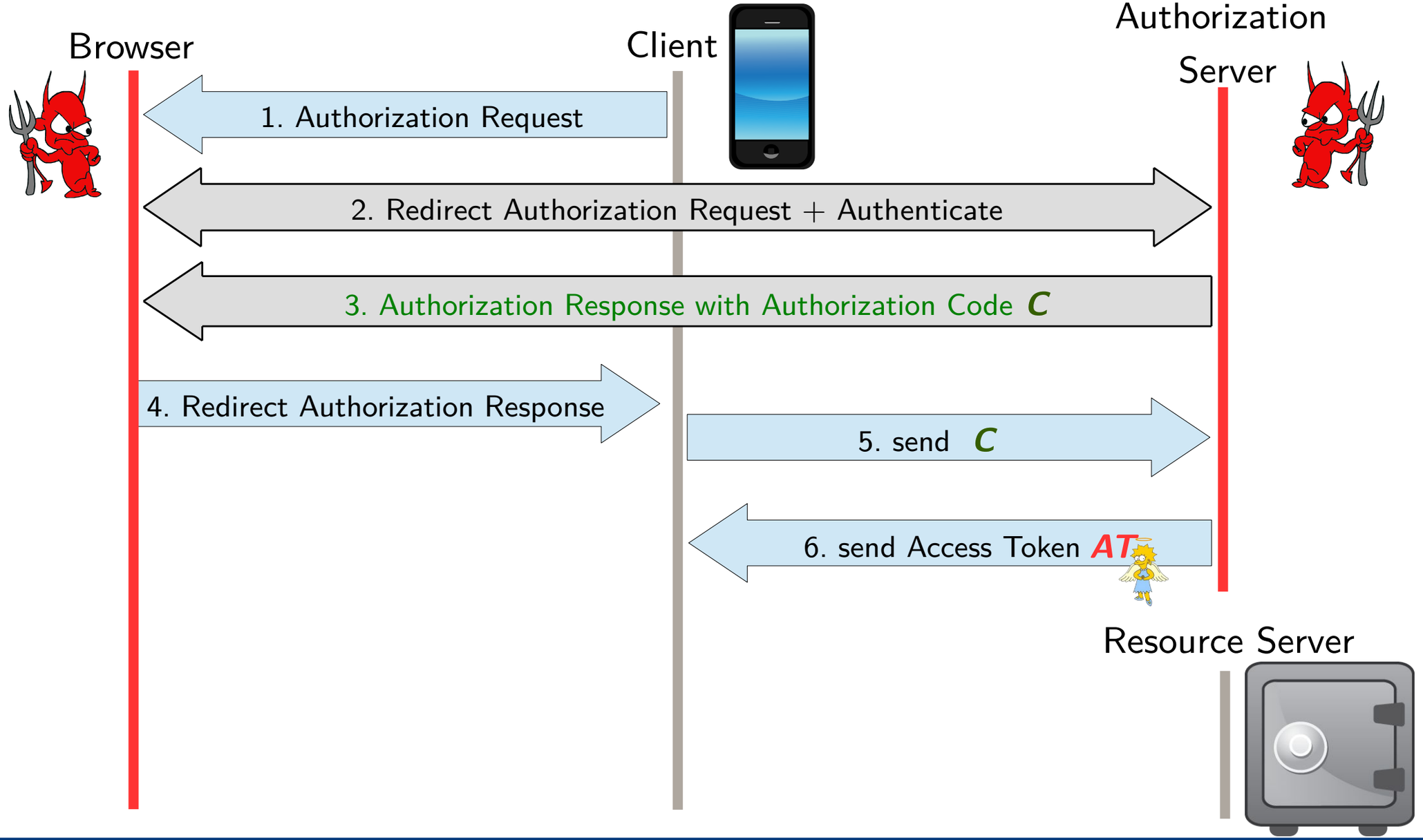

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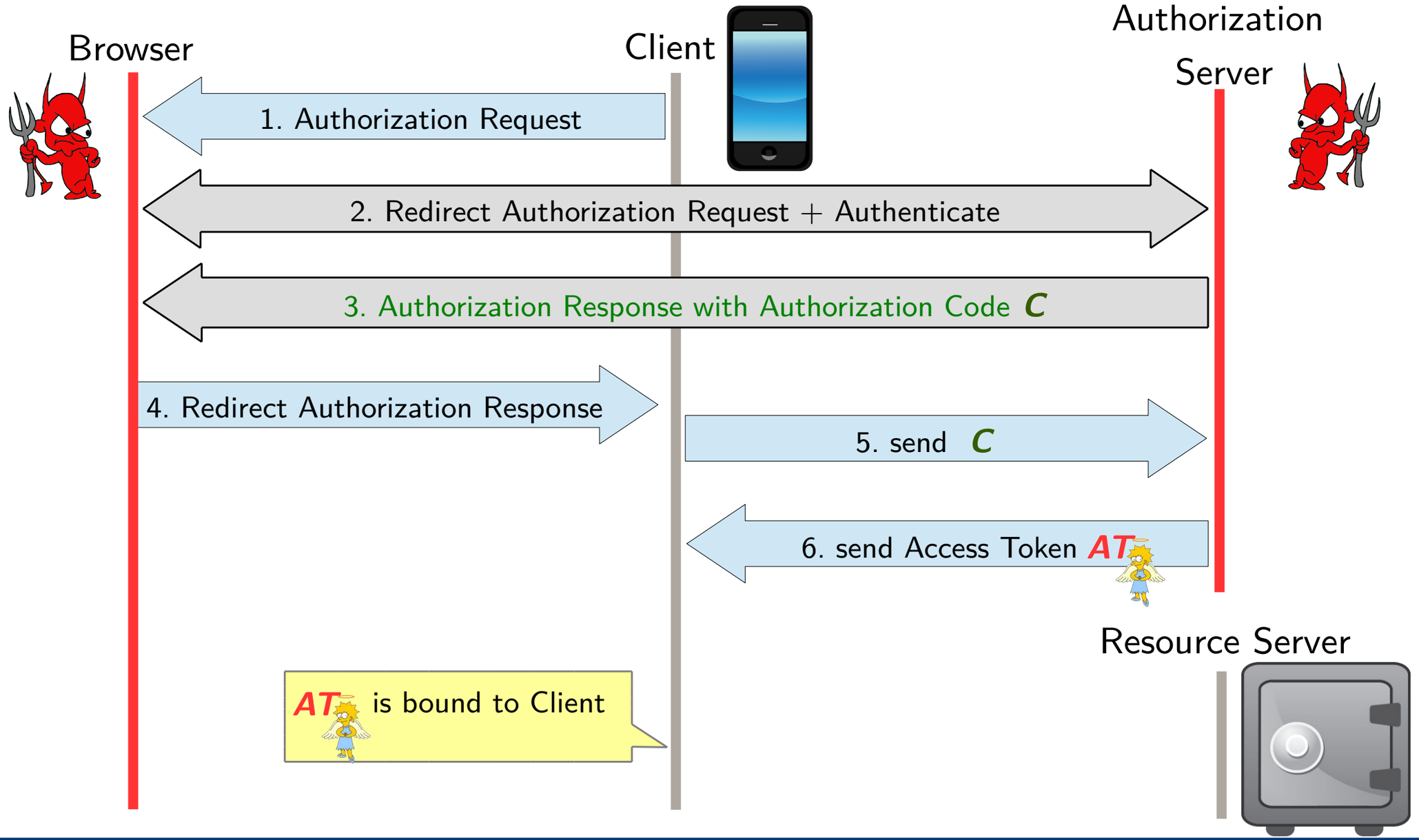

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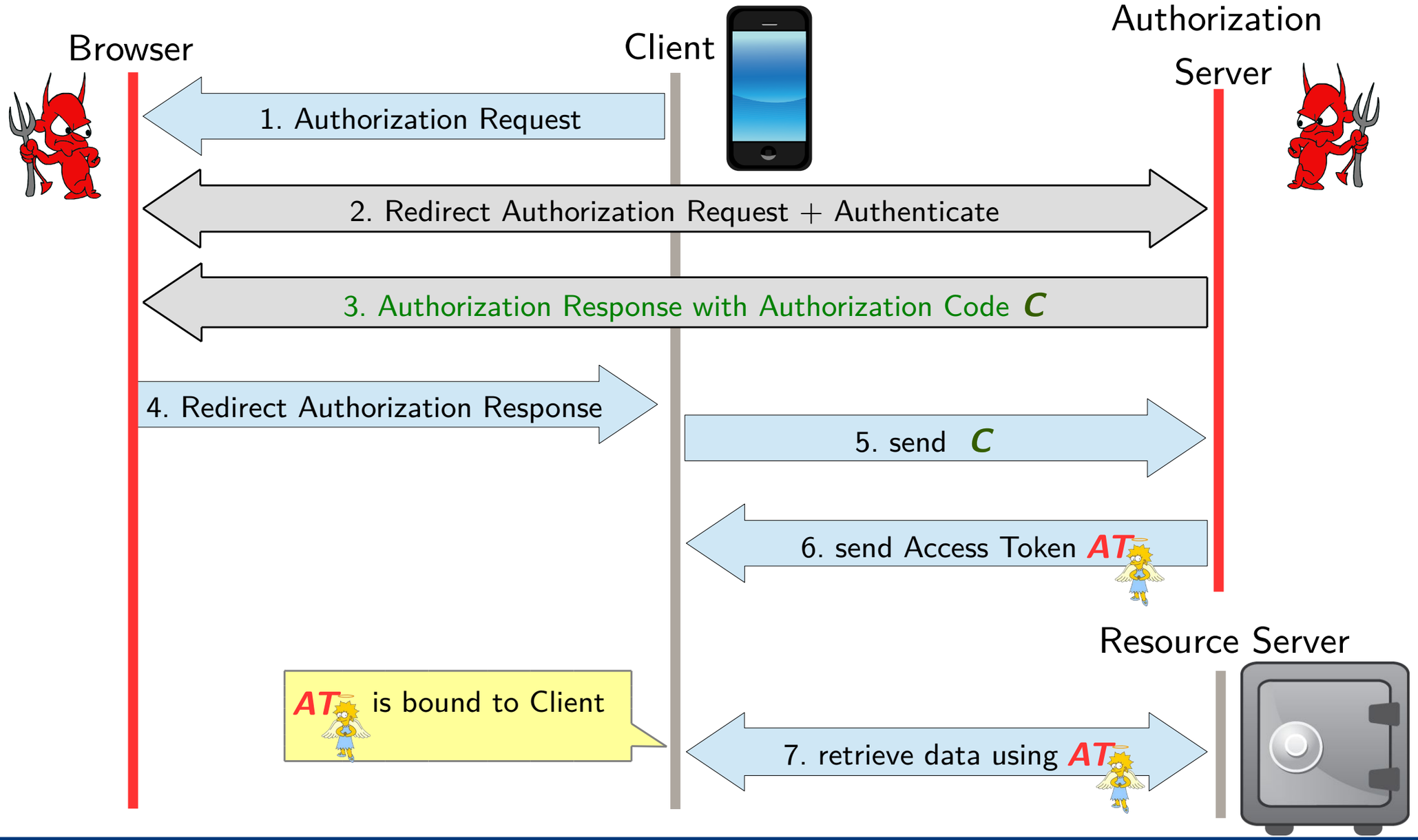

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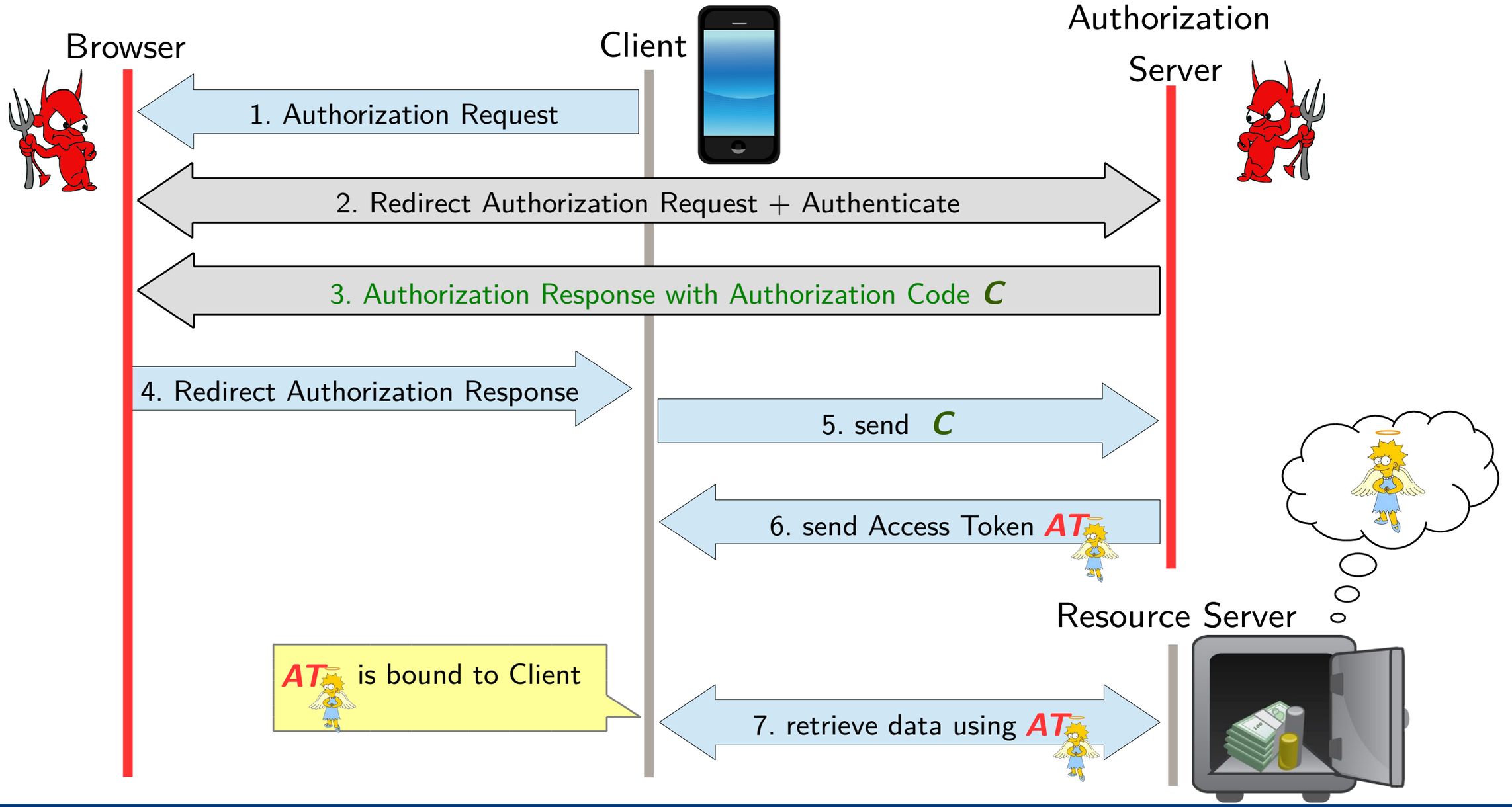
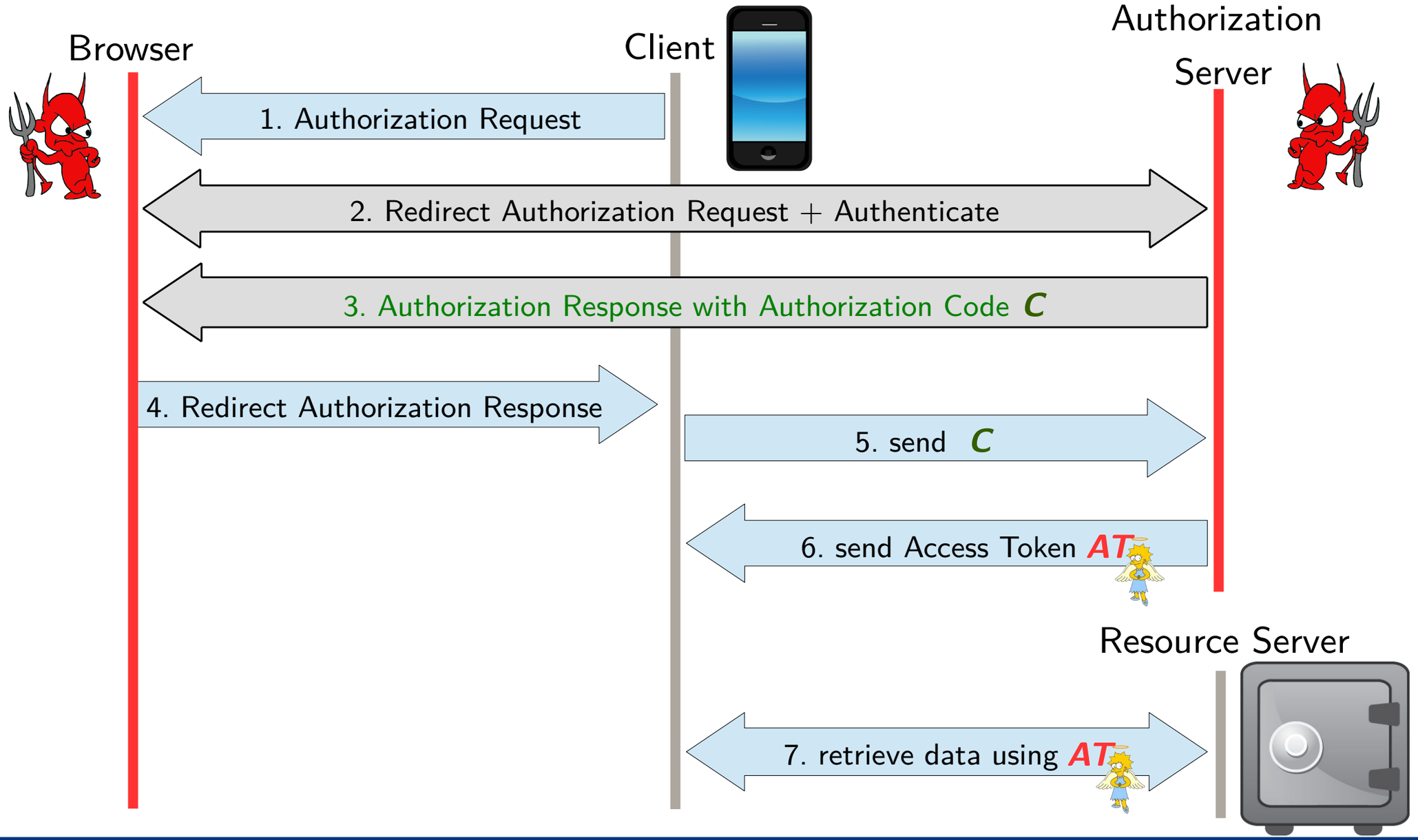


Mitigation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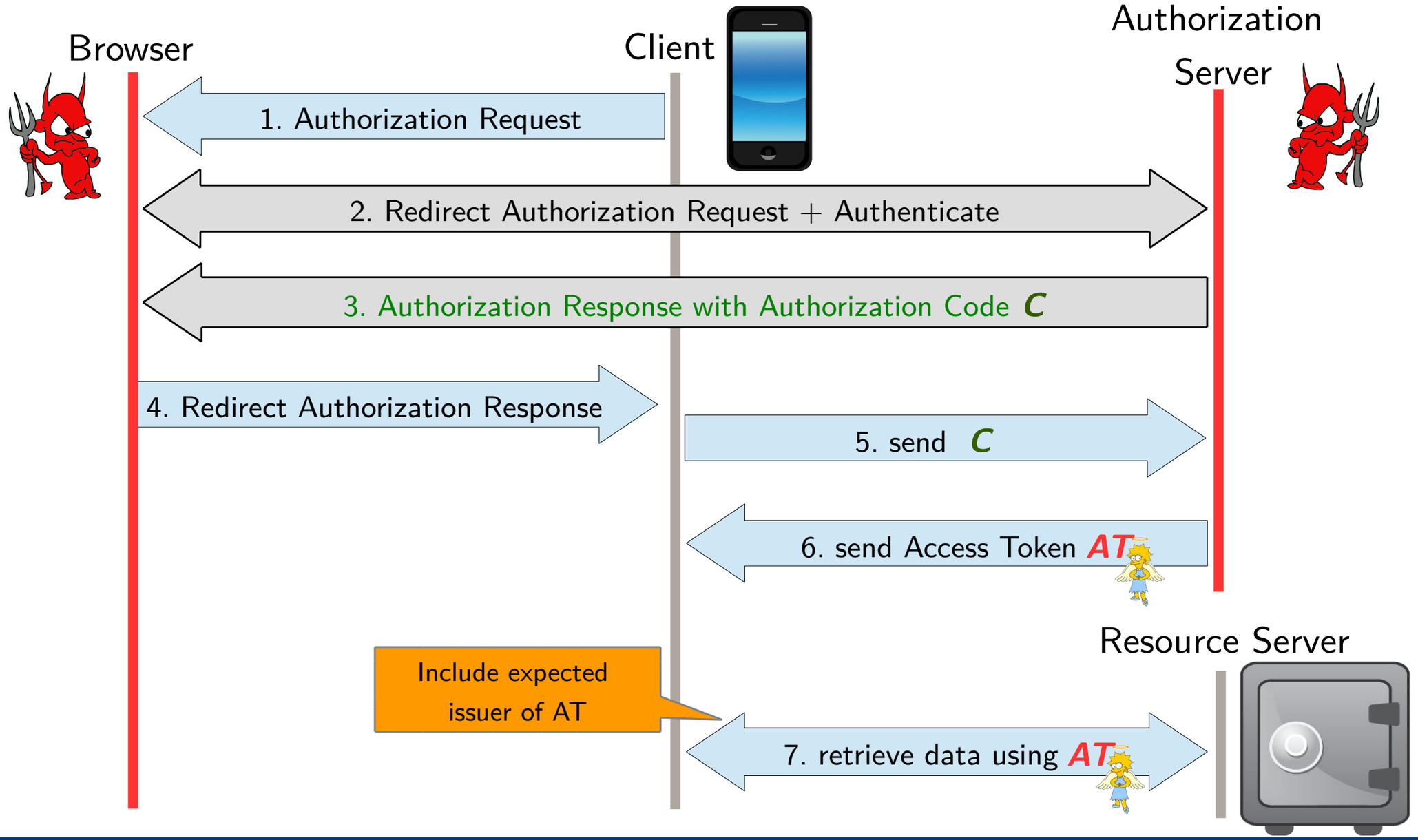
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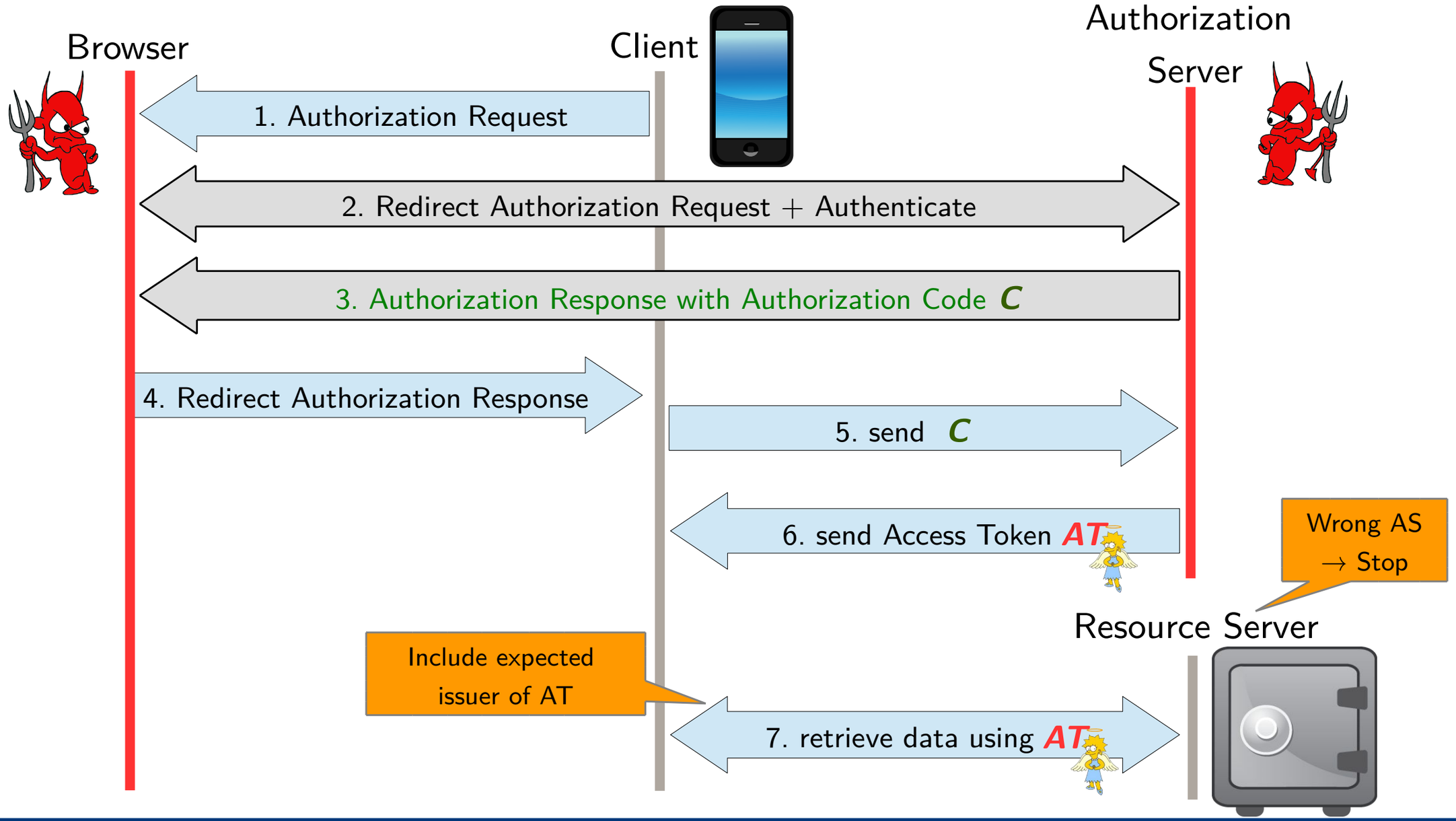
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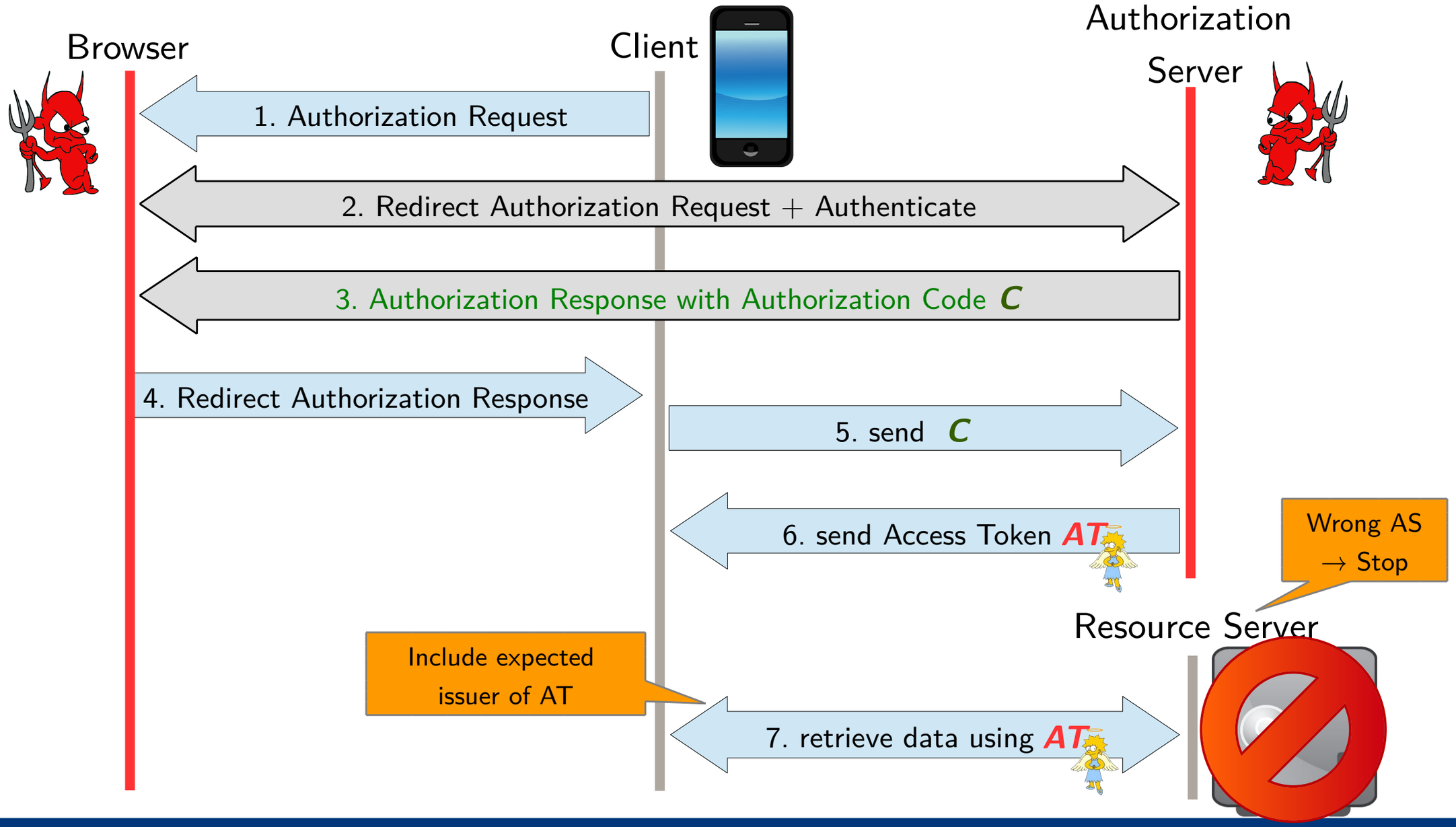
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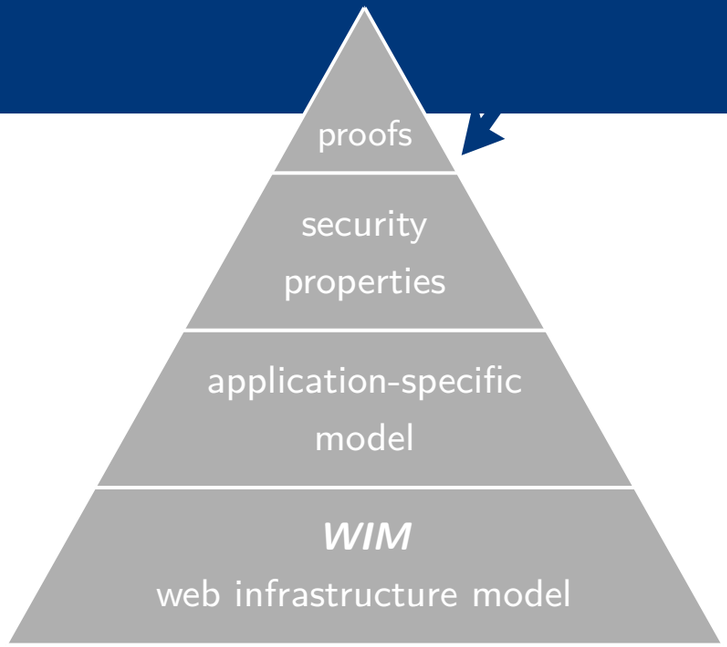
Attacks Found Through Our Formal Analysis

- Cuckoo's Token Attack
- Access Token Injection
- PKCE Chosen Challenge Attack
- Authorization Request Leak Attacks



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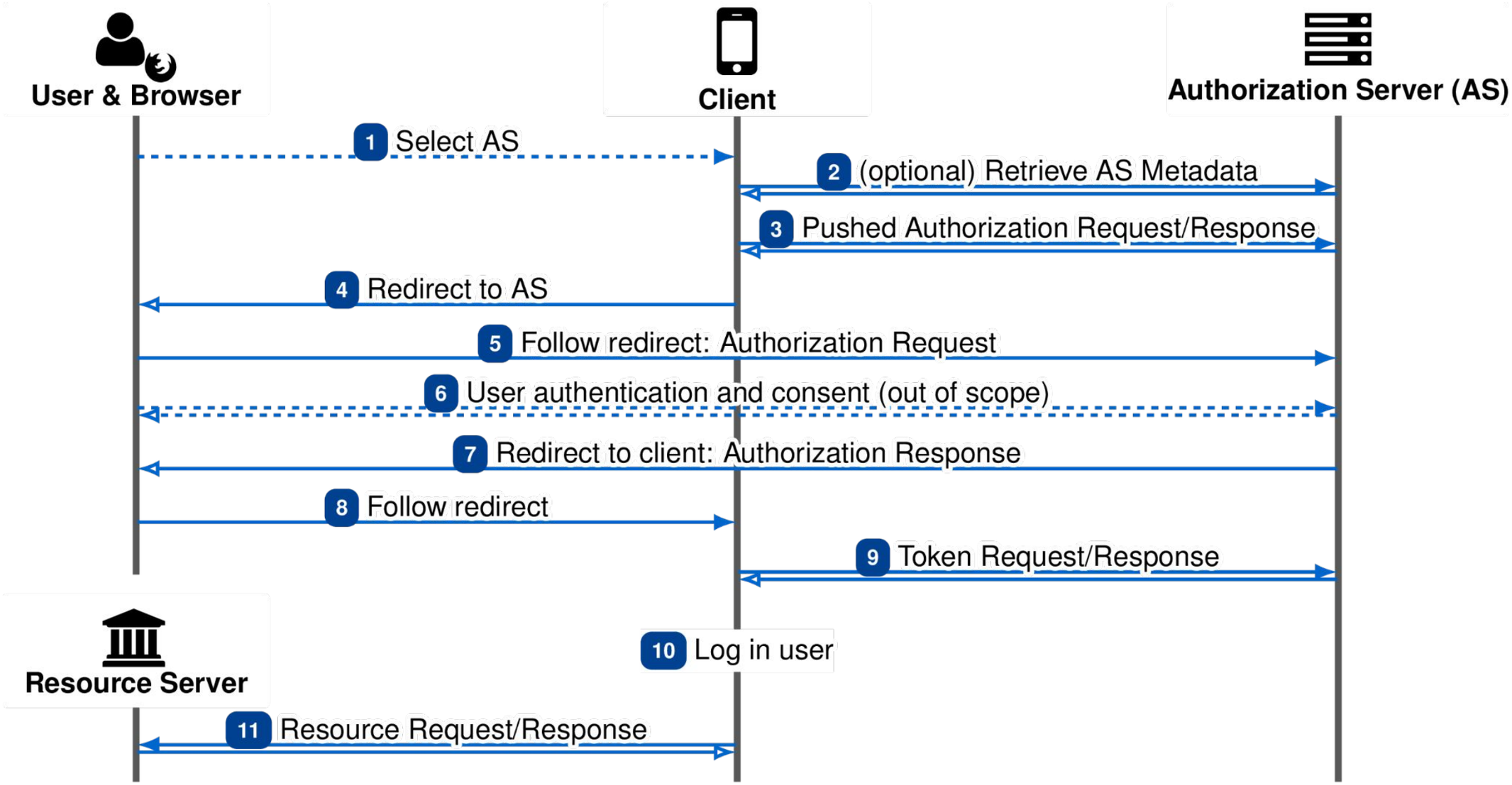

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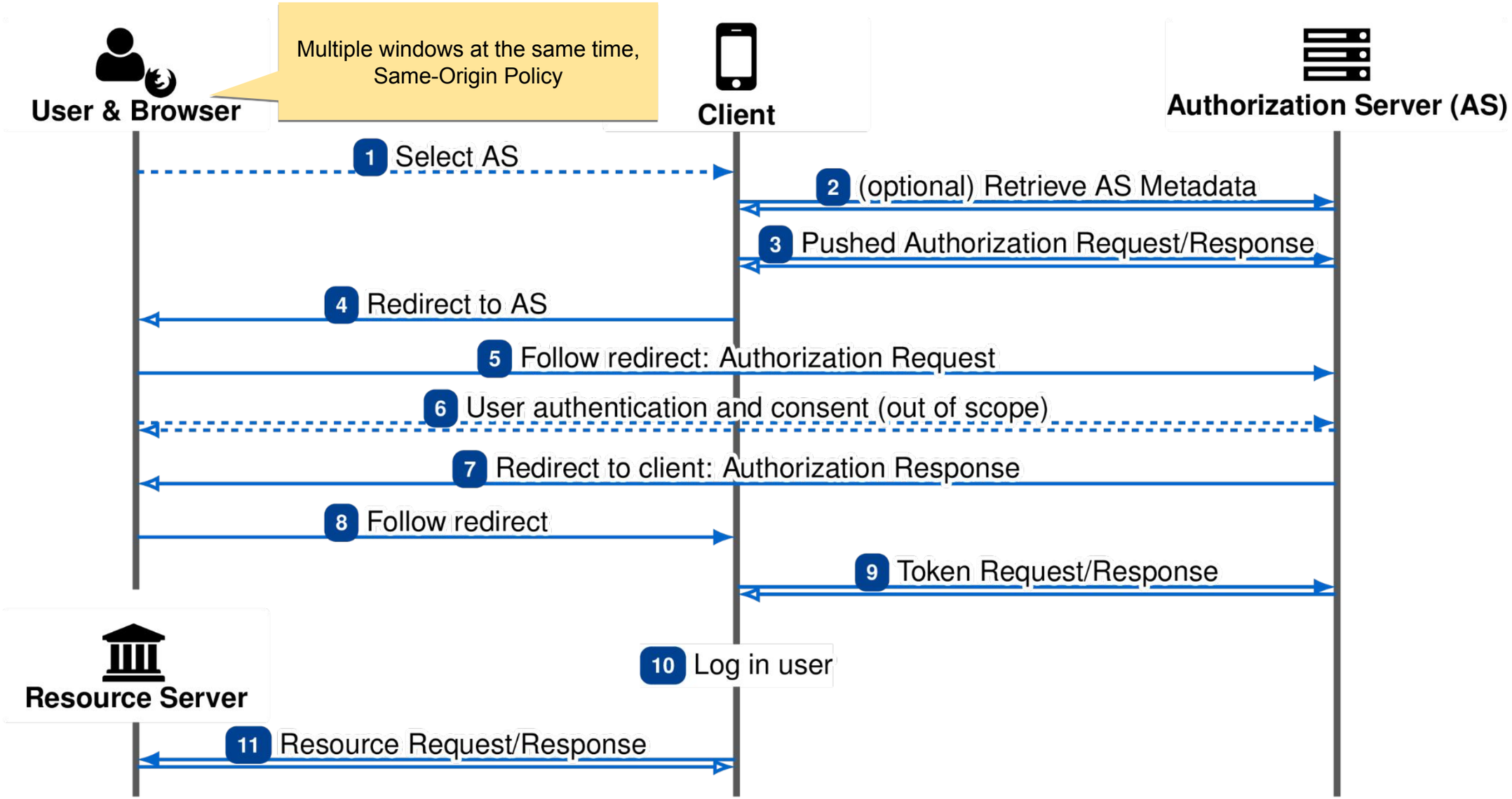

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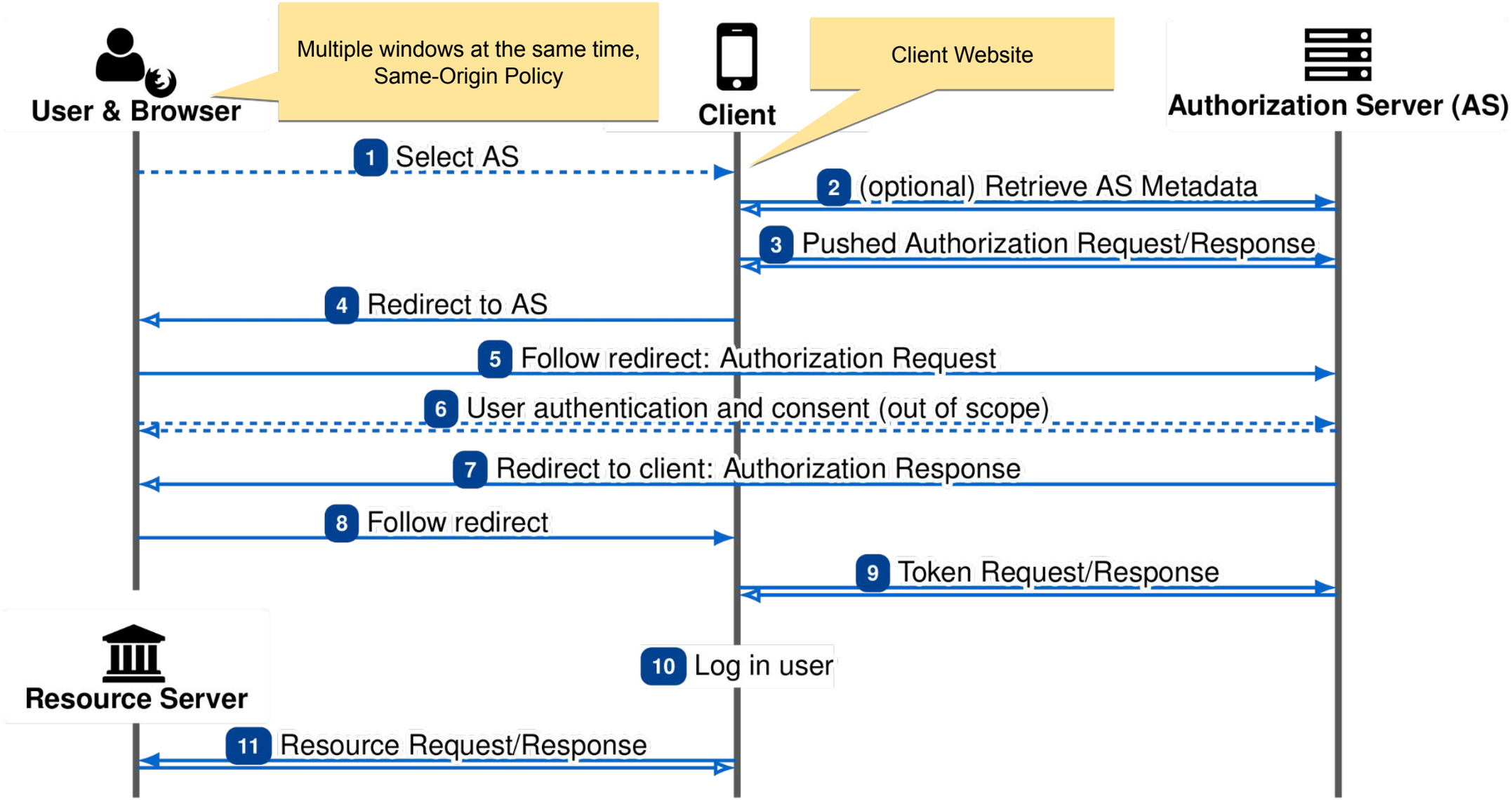

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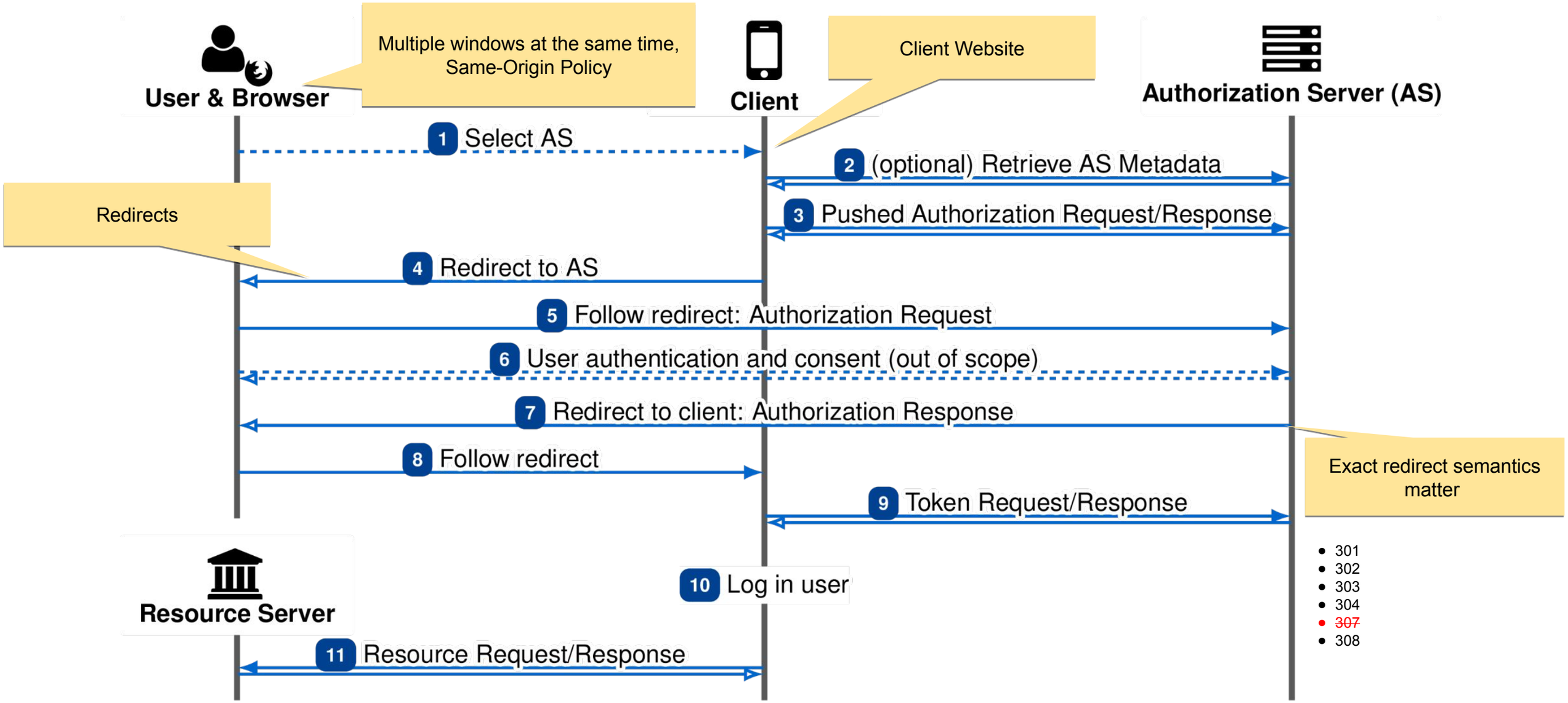

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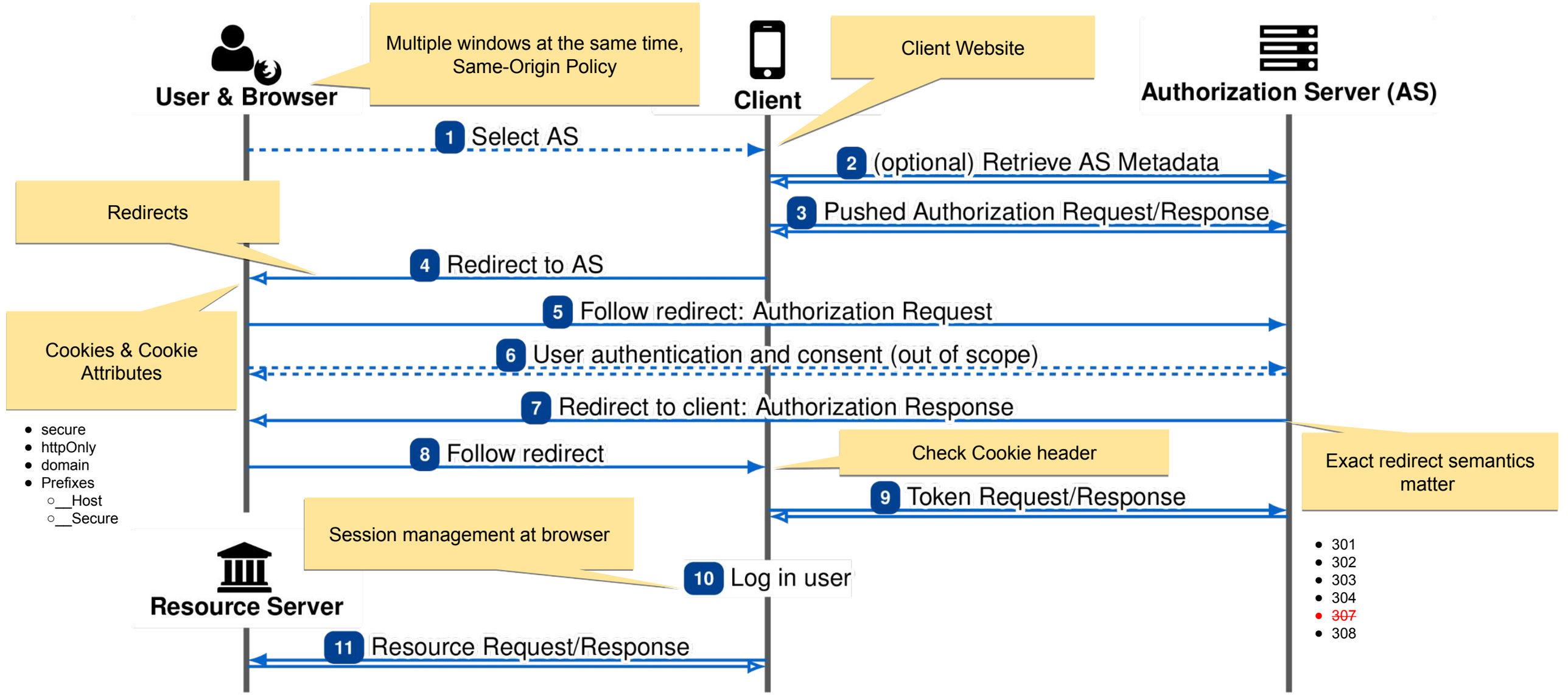

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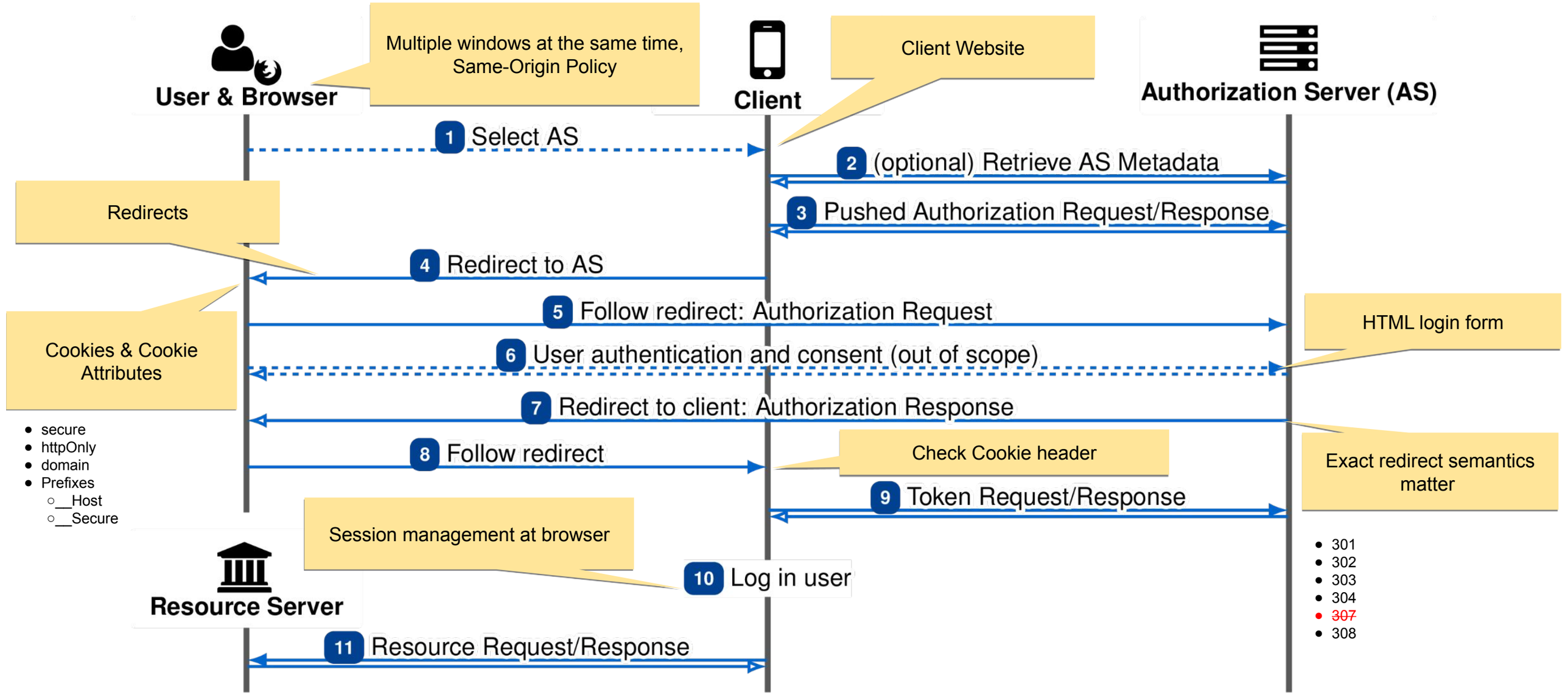

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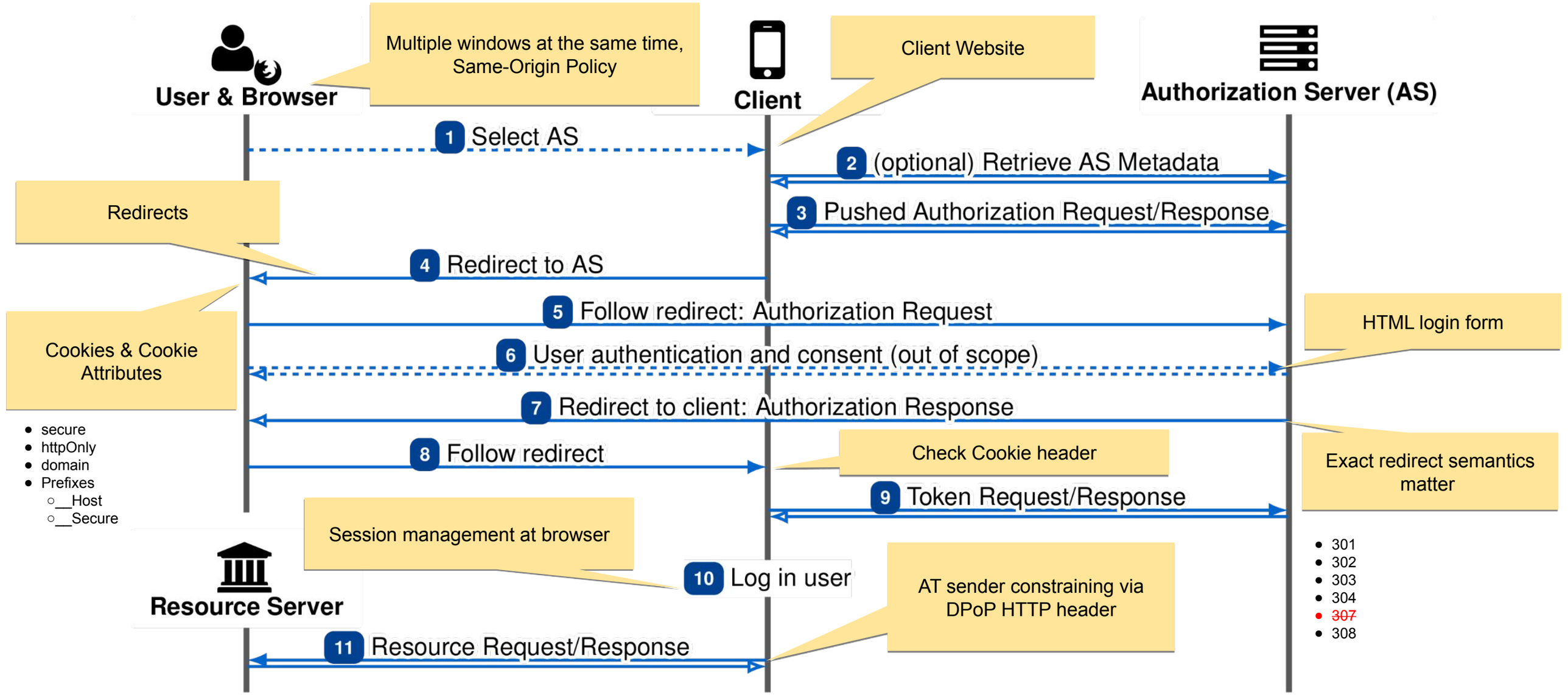

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Protocols/Standards We Have Analyzed So Far

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- OpenID Connect
- OpenID FAPI 1.0 and FAPI 2.0
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- GNAP
- Mozilla BrowserID
- OID4VP/VCI (ongoing work)
- Web Payment APIs
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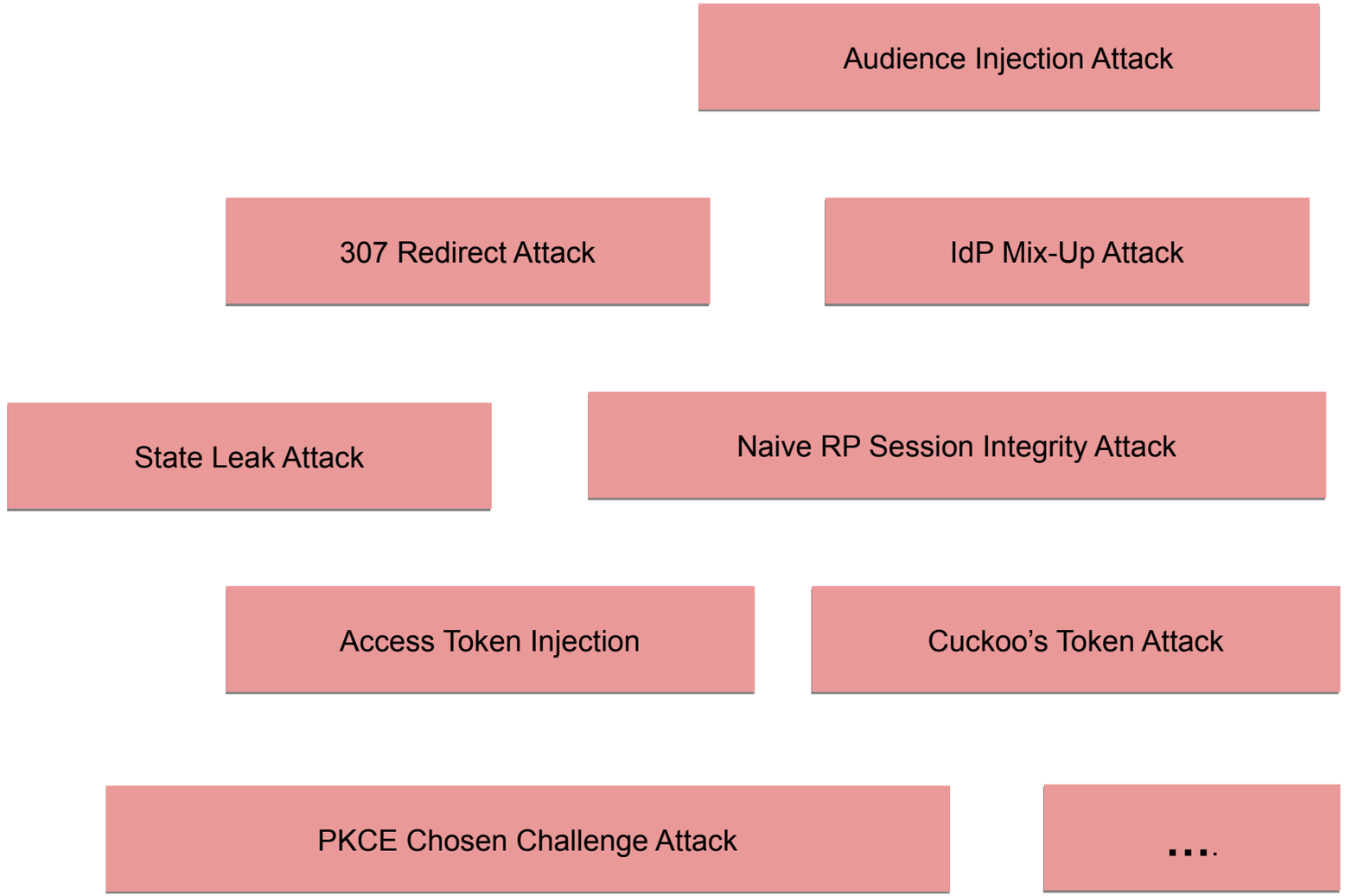

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under submission: affects several standards related to OAuth 2.0, OpenID Connect, FAPI, CIBA, ...

Audience Injection Attack

307 Redirect Attack

IdP Mix-Up Attack

State Leak Attack

Naive RP Session Integrity Attack

Access Token Injection

Cuckoo's Token Attack

PKCE Chosen Challenge Attack

....

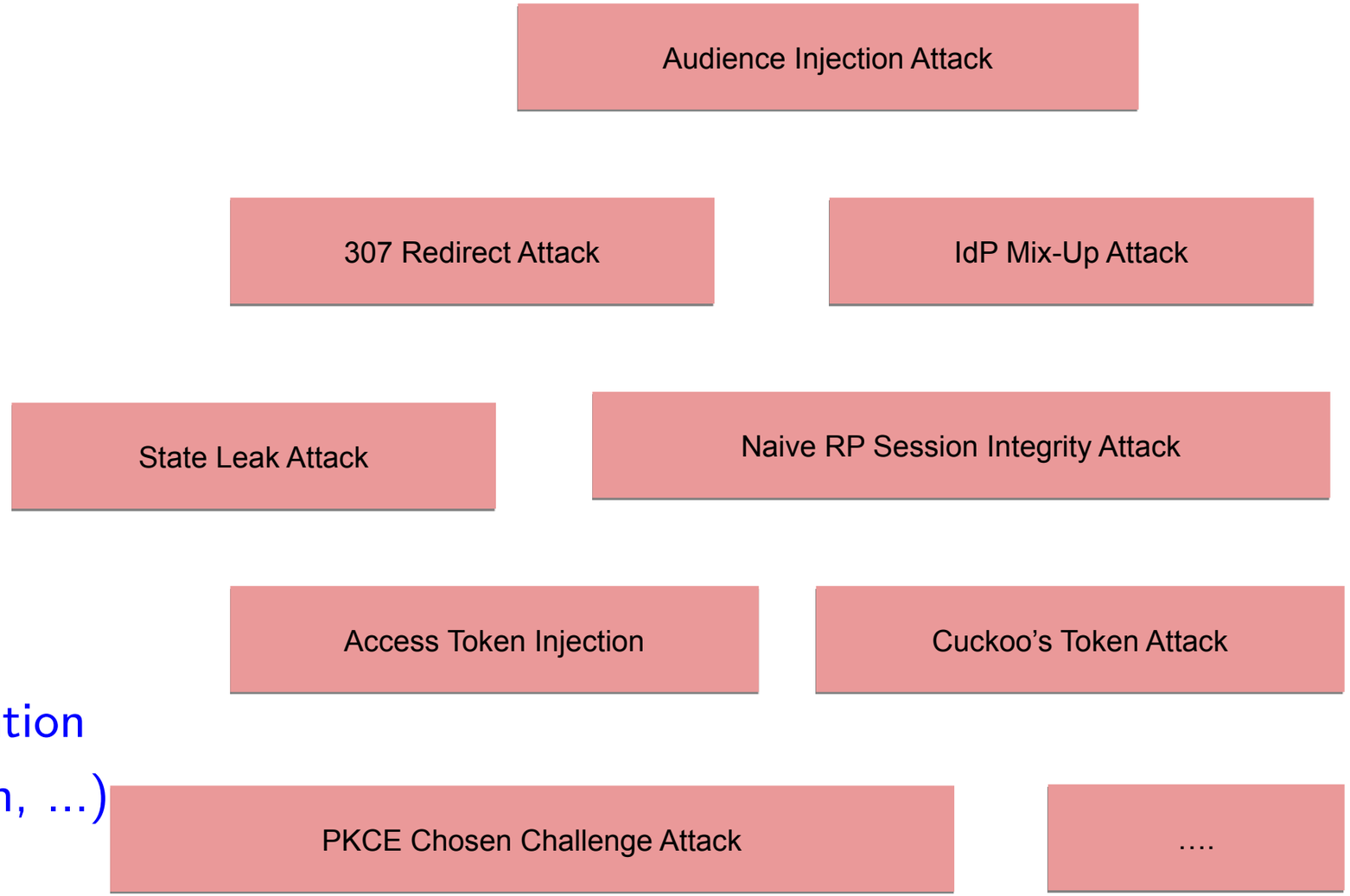
But this was not about breaking things ...

We always started out with

1) Modeling
2) Formalizing security properties
3) Trying to prove properties

Our findings resulted in fixed/improved and formally analyzed standards.

Close interaction with standardization bodies (IETF, OpenID Foundation, ...)



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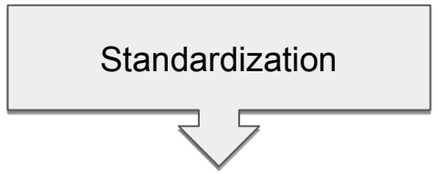

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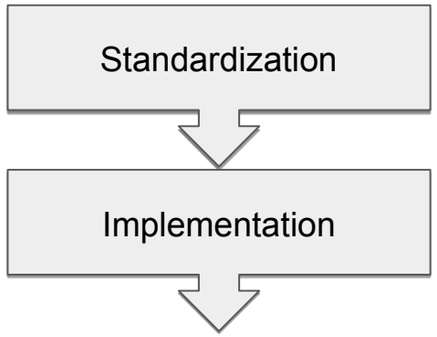

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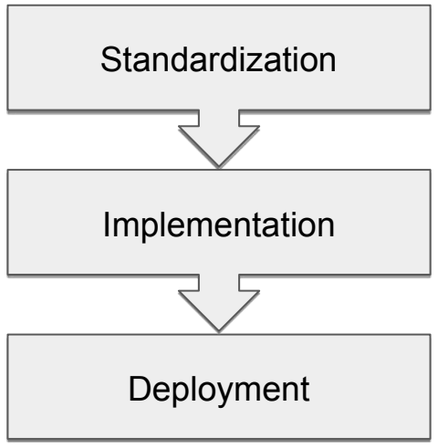

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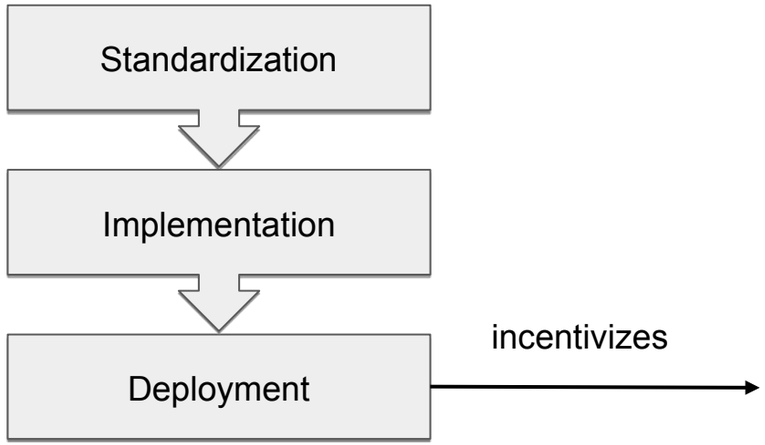

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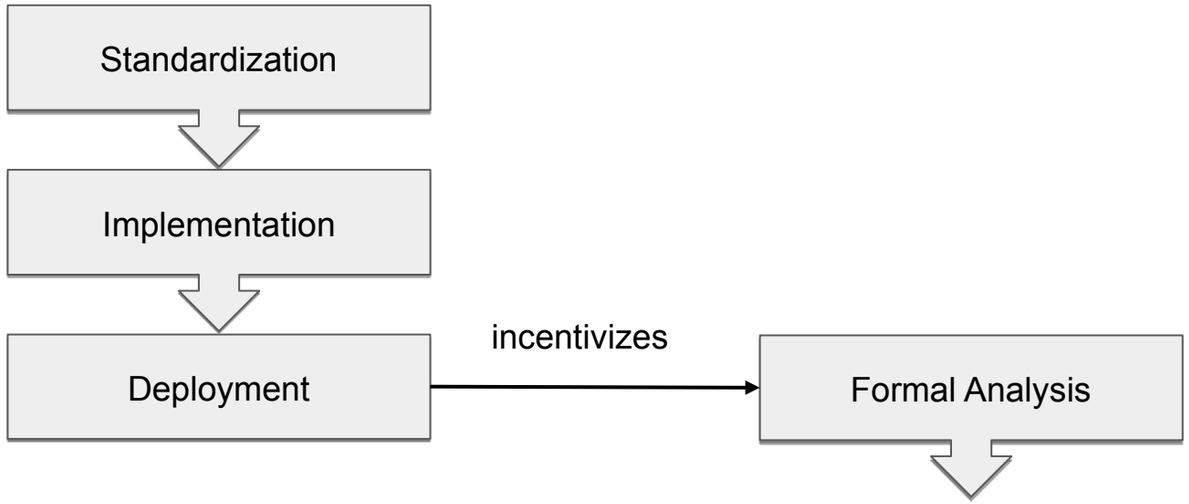

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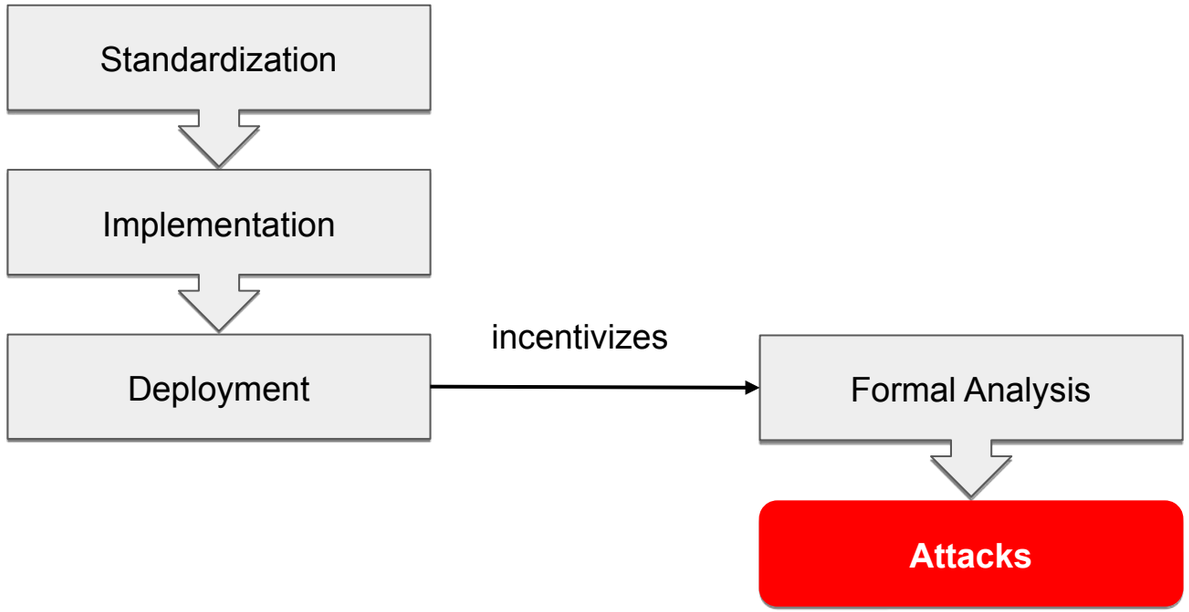

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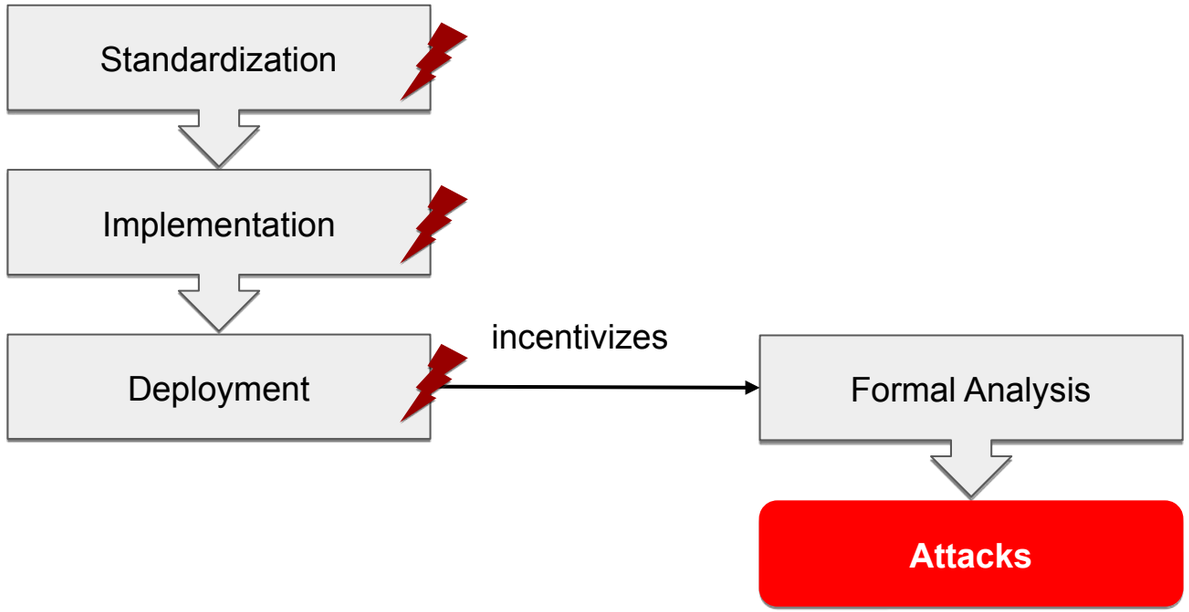

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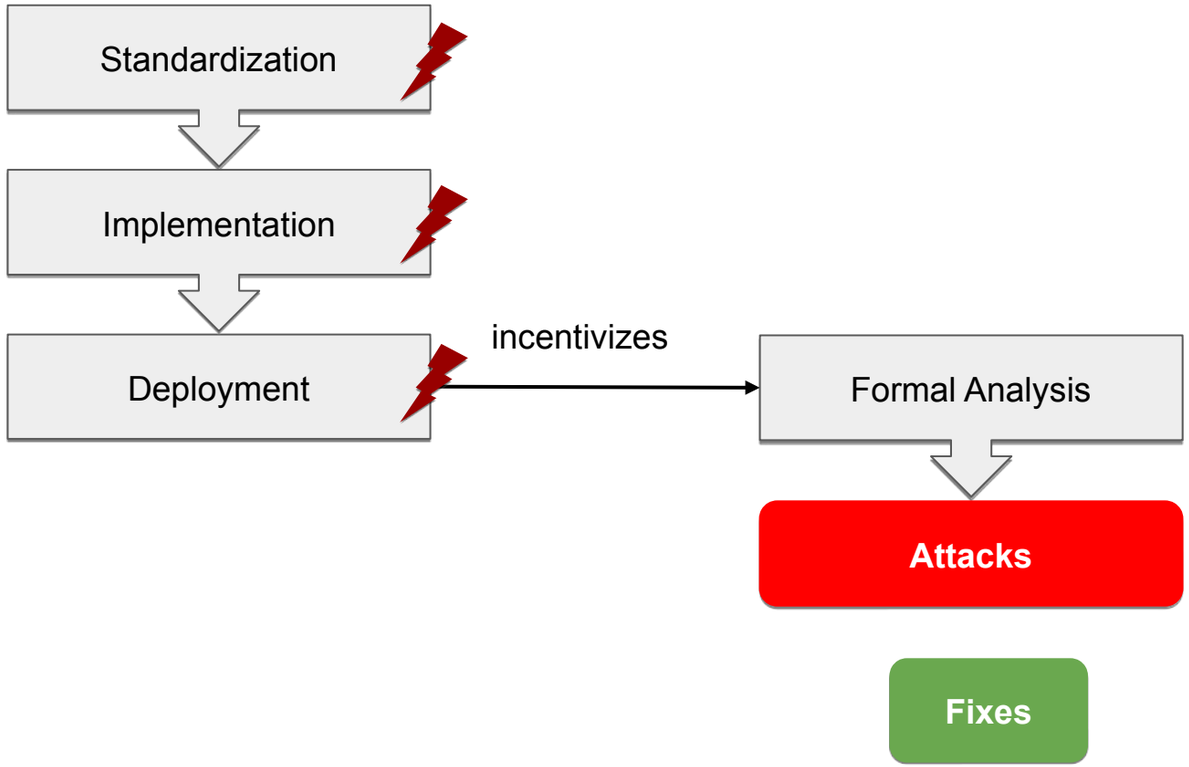

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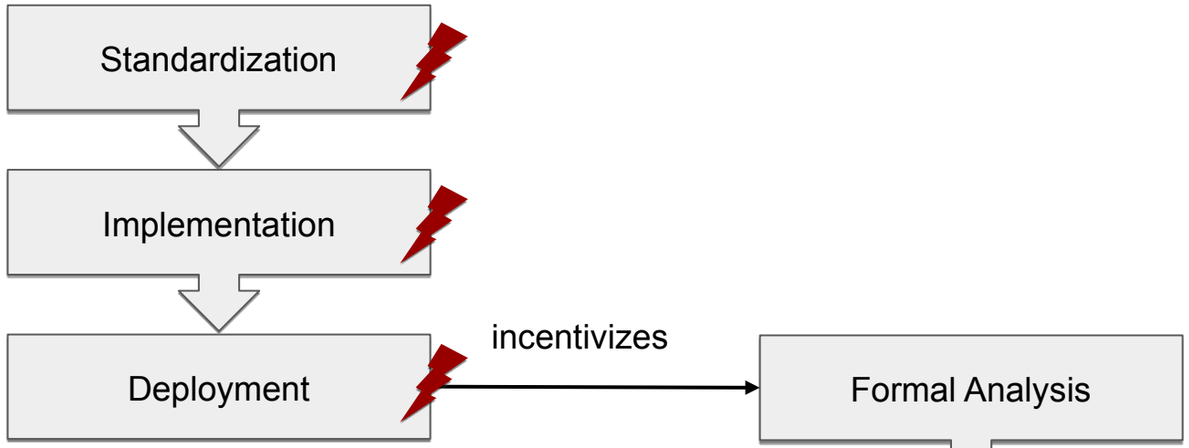

Fixes

Working closely with Standardization Bodies

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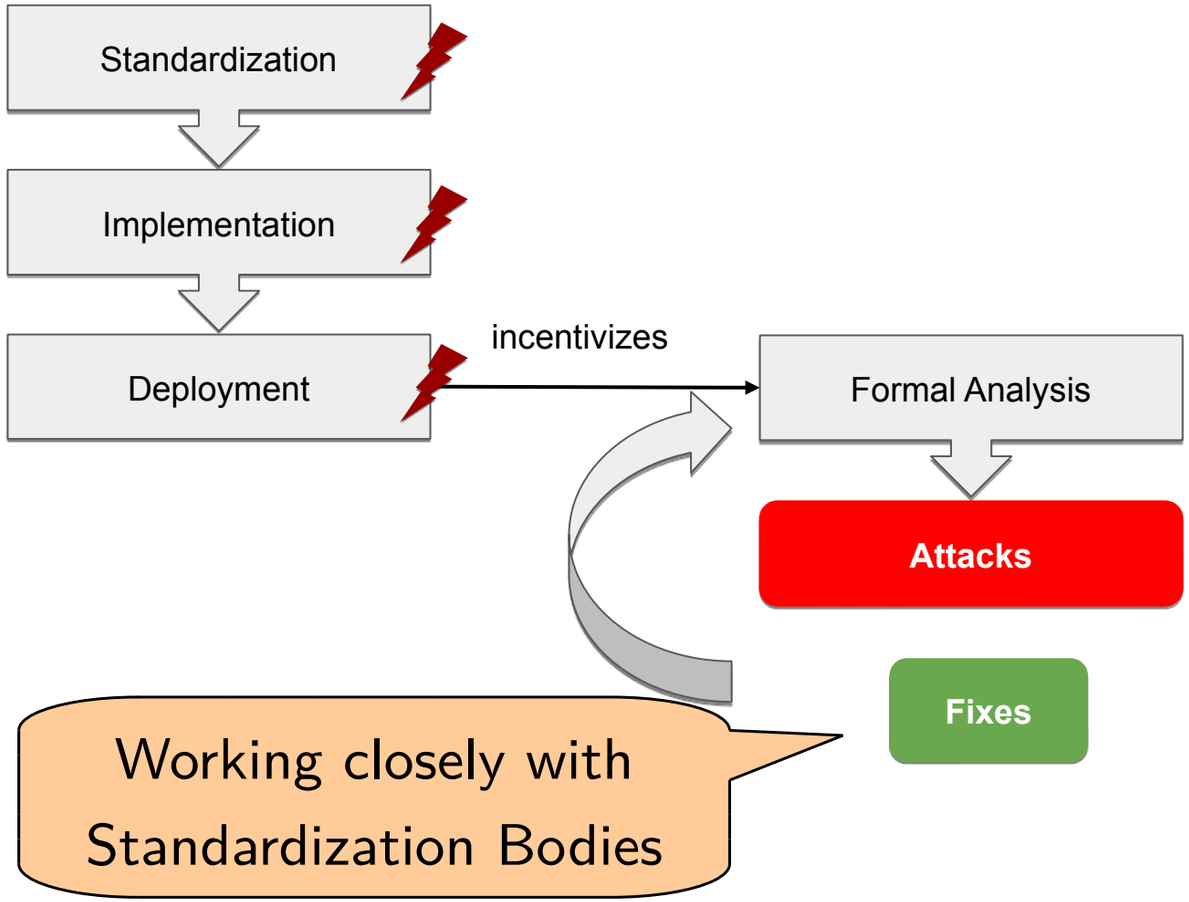

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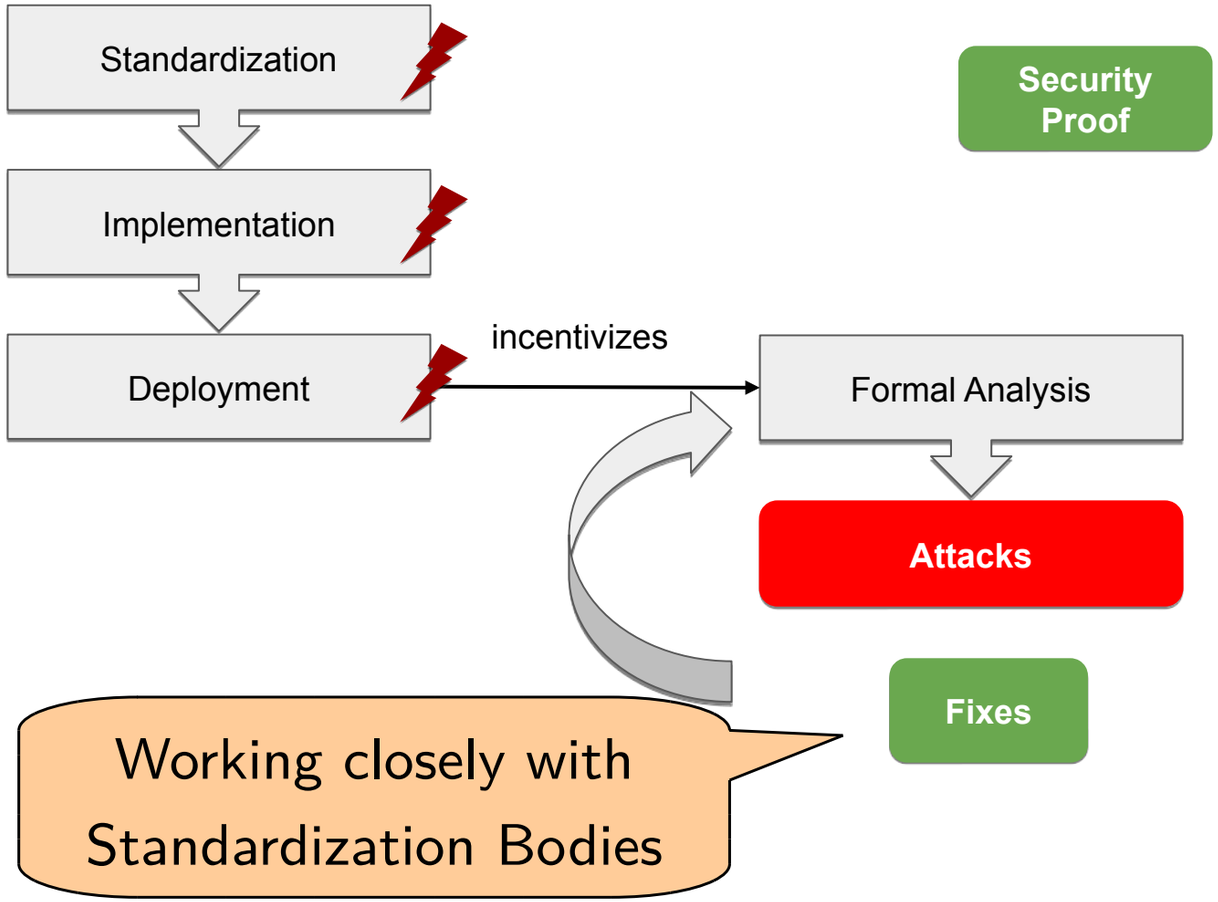

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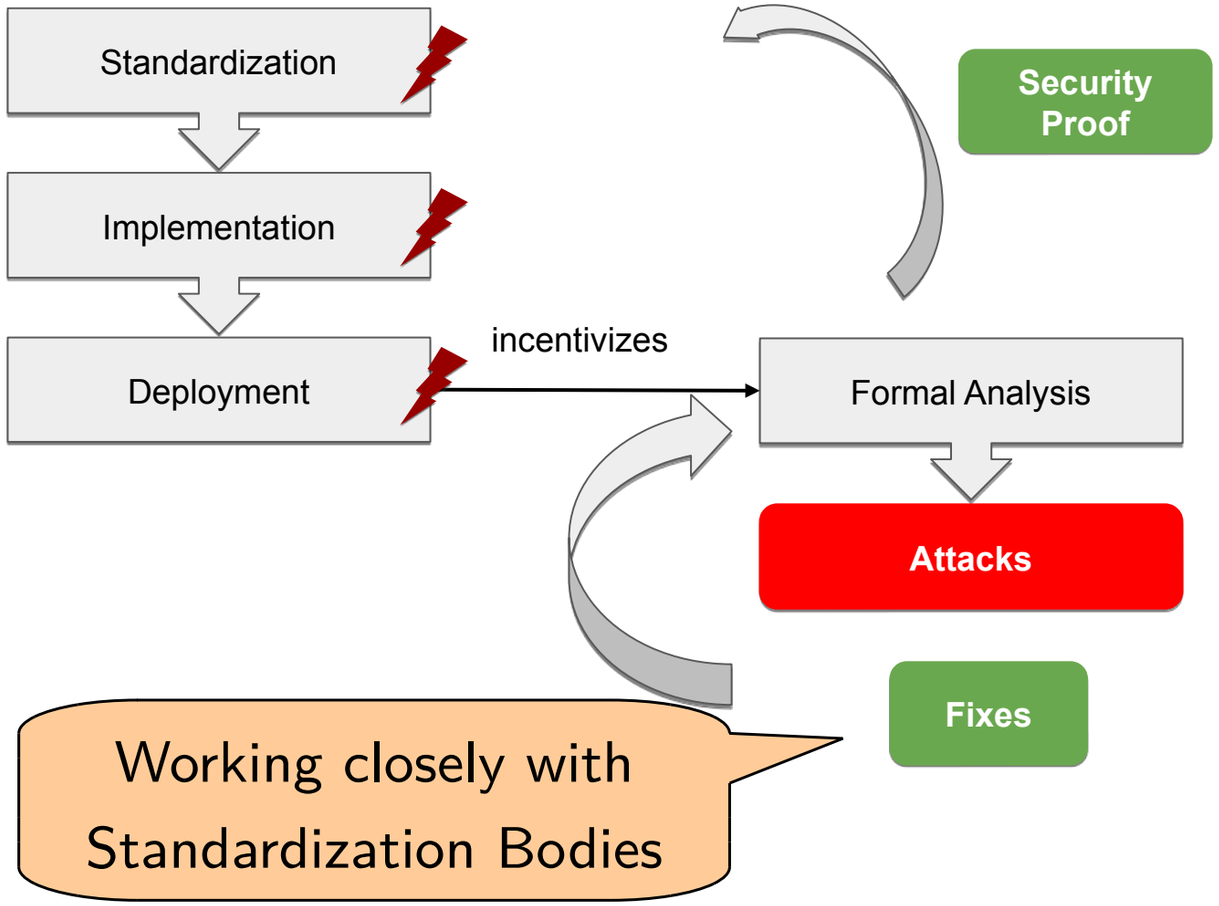

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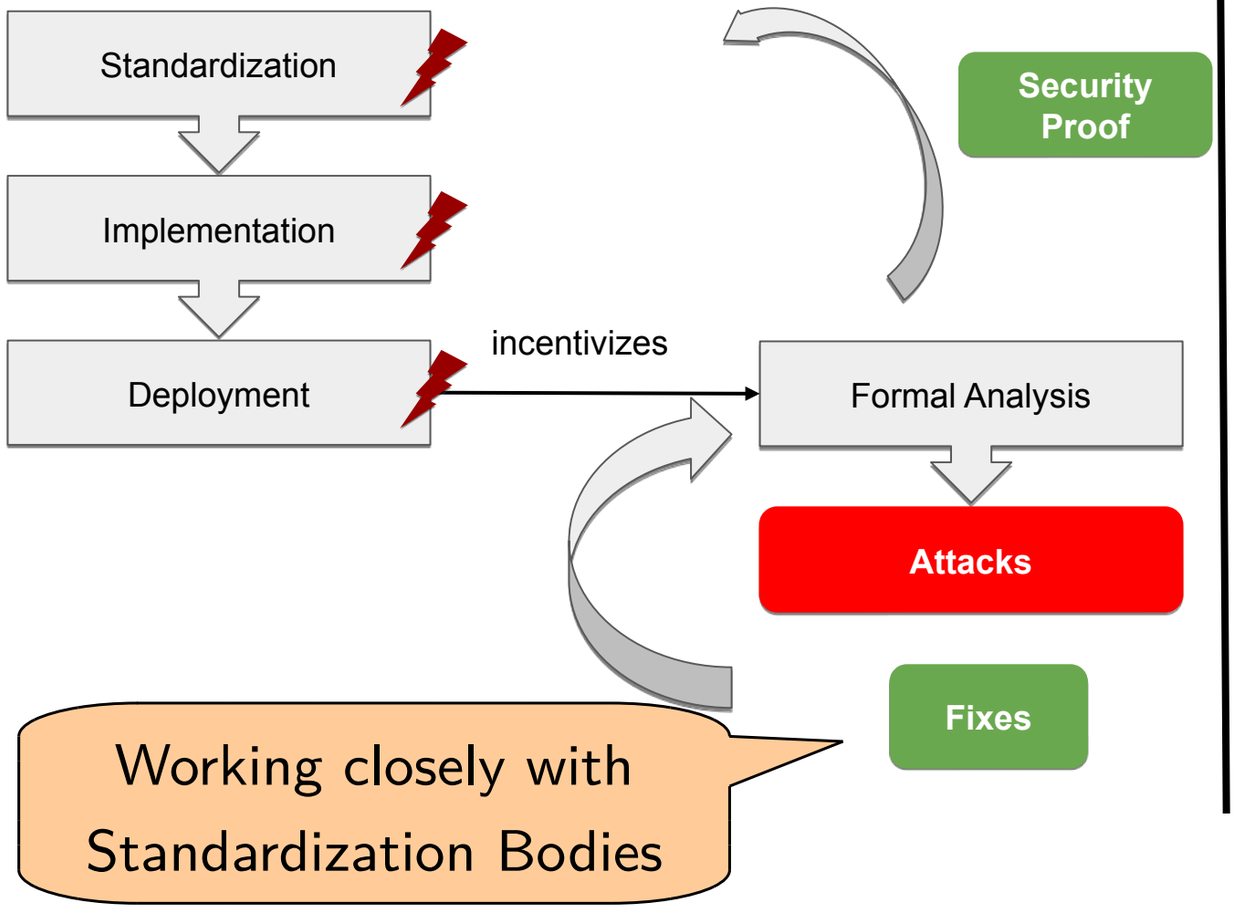
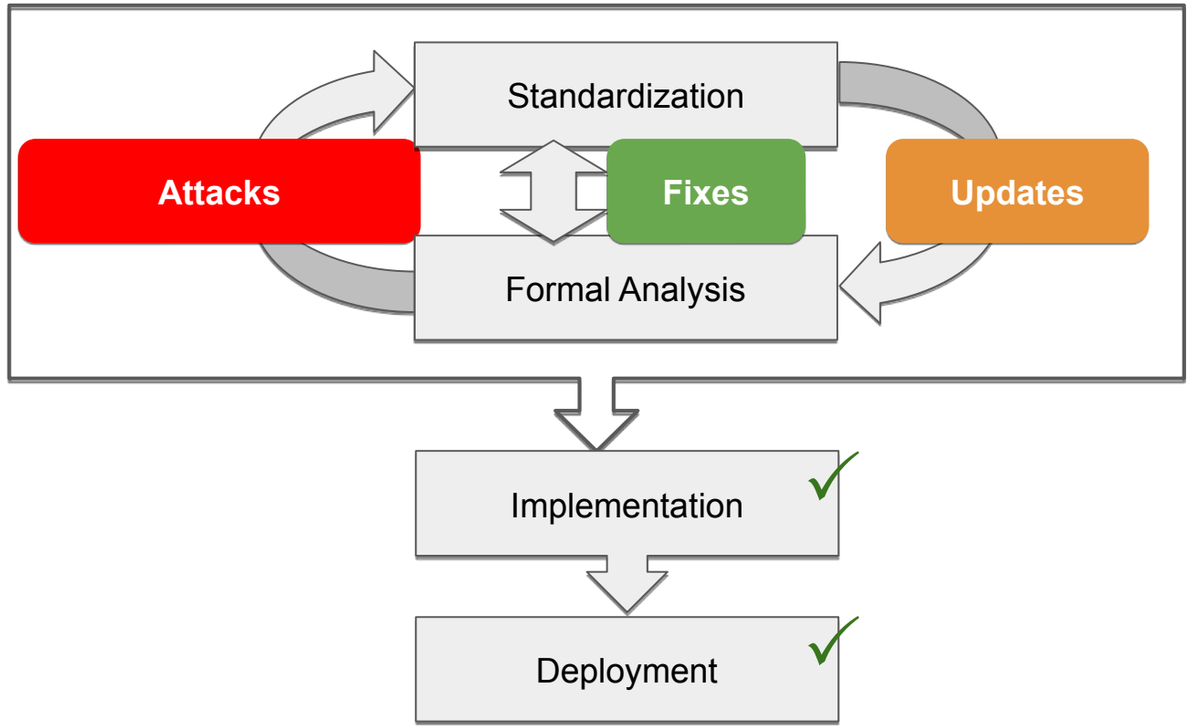
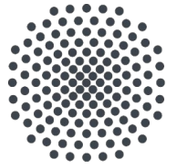


Now we are often part of the standardization process (OpenID Foundation, IETF):



Towards Mechanizing the WIM

DY*



University of Stuttgart
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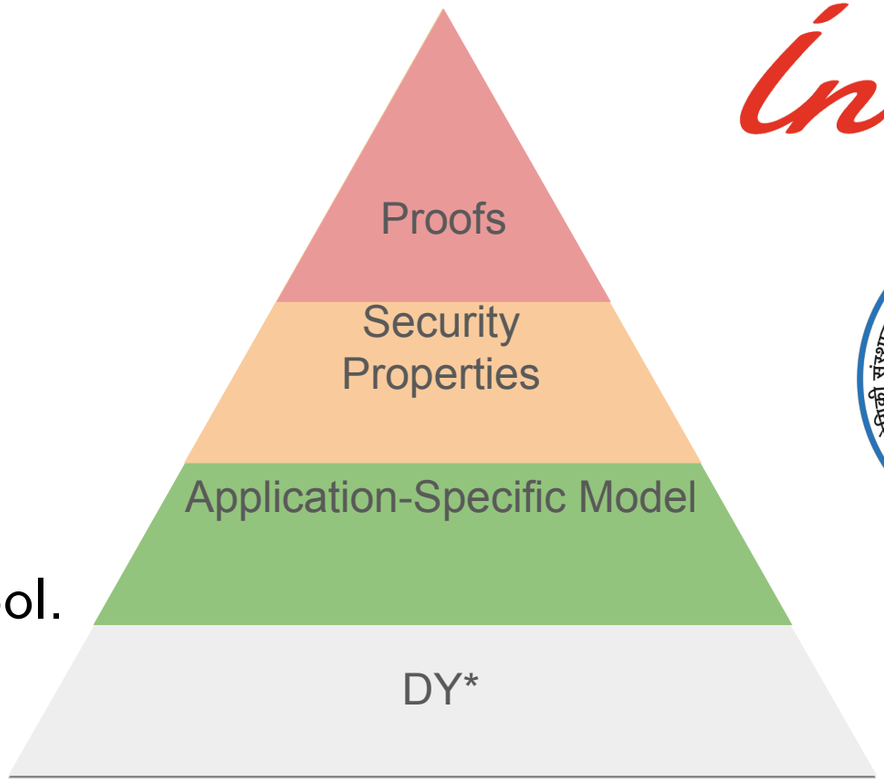
Lancaster University

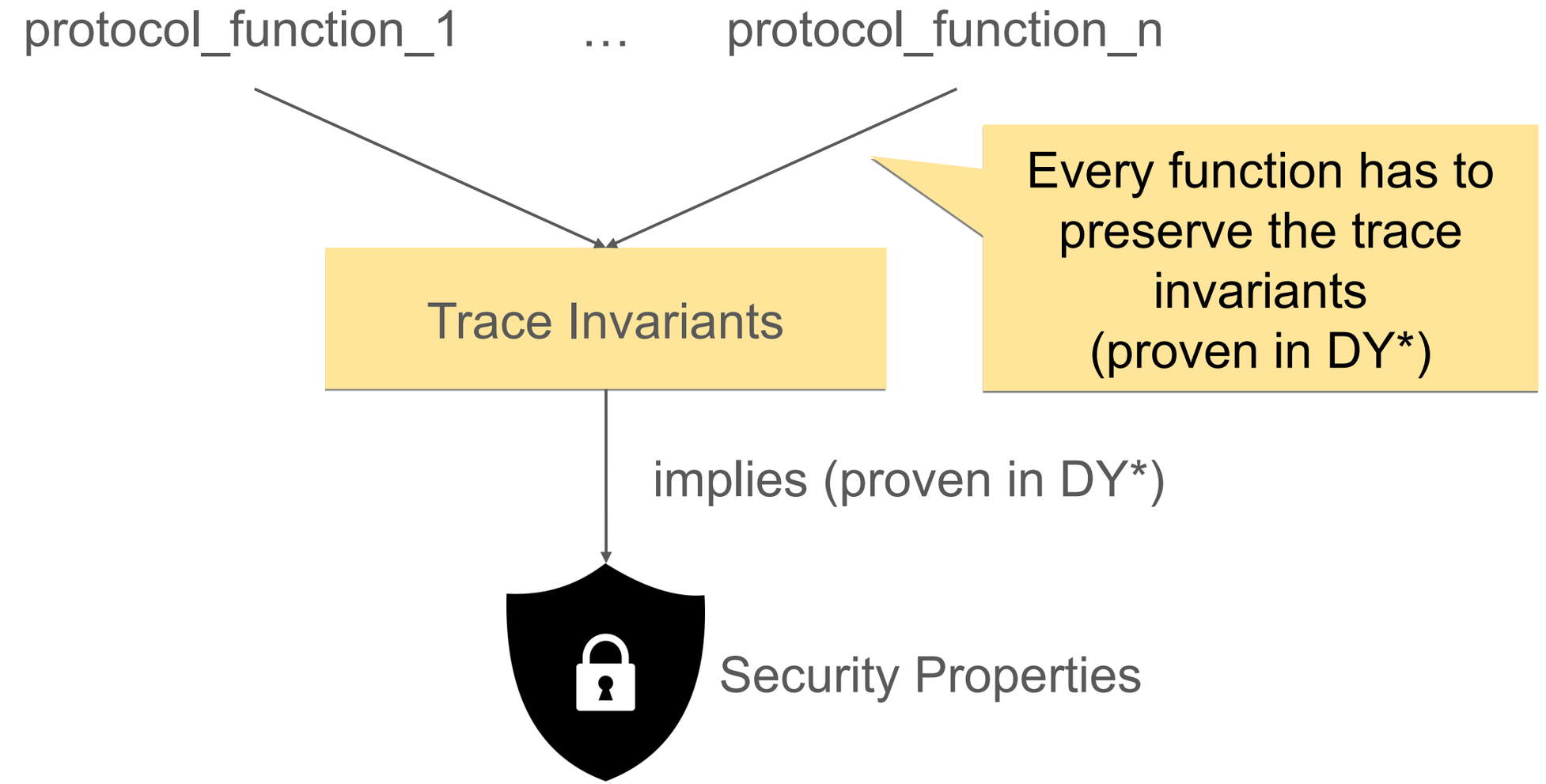
Inria

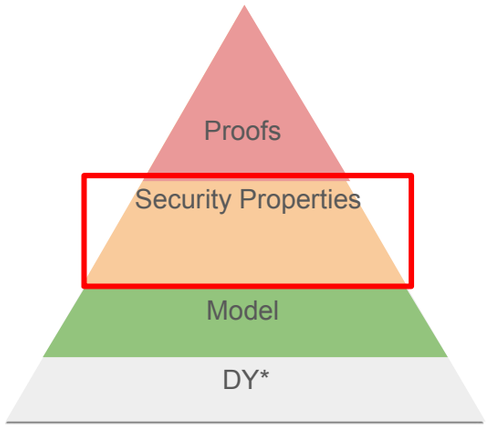
- Dolev-Yao model implemented in F*
- Enables fine-grained analysis up to implementation level
- Mechanized (tool checked) proofs
- Partially automated proofs
- Executable models
- Highly modular



At this point, general crypto protocol analysis tool.

Security Properties





Case Studies So Far




- Signal Messaging Protocol
  - Unbound number of rounds (ratcheting)
  - Forward Secrecy & Post Compromise Security
- Automatic Certificate Management Environment (ACME)
  - One of the largest & most in-depth formal security analyses in the literature (16.000 LoC)
  - ACME client model can interoperate with real-world server
- Needham-Schroeder(-Lowe), ISO-DH, and ISO-KEM

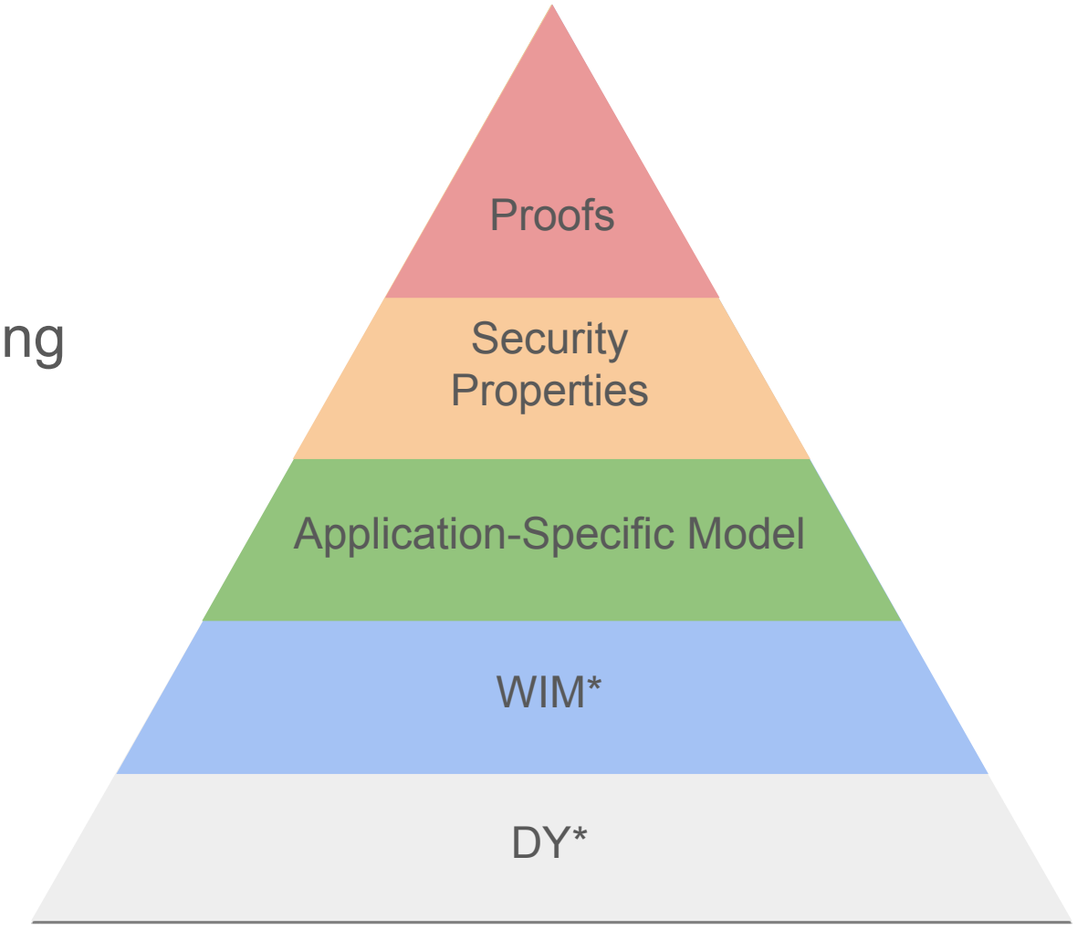
DY* Current/Future Work

Near-term

- Improving proof automation
- Database library
- HTTP library for sending requests and receiving responses

Long-term

- WIM*
  - Generic web server
  - Browser



Conclusion

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- SSO protocols and standards are fun!
- The WIM is the most comprehensive model of the web infrastructure to date
- And has proved to be instrumental for formal analysis
- Several standards analyzed based on the WIM
- (Almost always) found new attacks and/or attack classes
- Proposed fixes
- Proved fixed standards secure in the WIM
  (under precisely formulated assumptions)
- Direct impact on standards
- Close collaboration with standardization bodies
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Overview of formal methods for web security: Michele Bugliesi, Stefano Calzavara, Riccardo Focardi