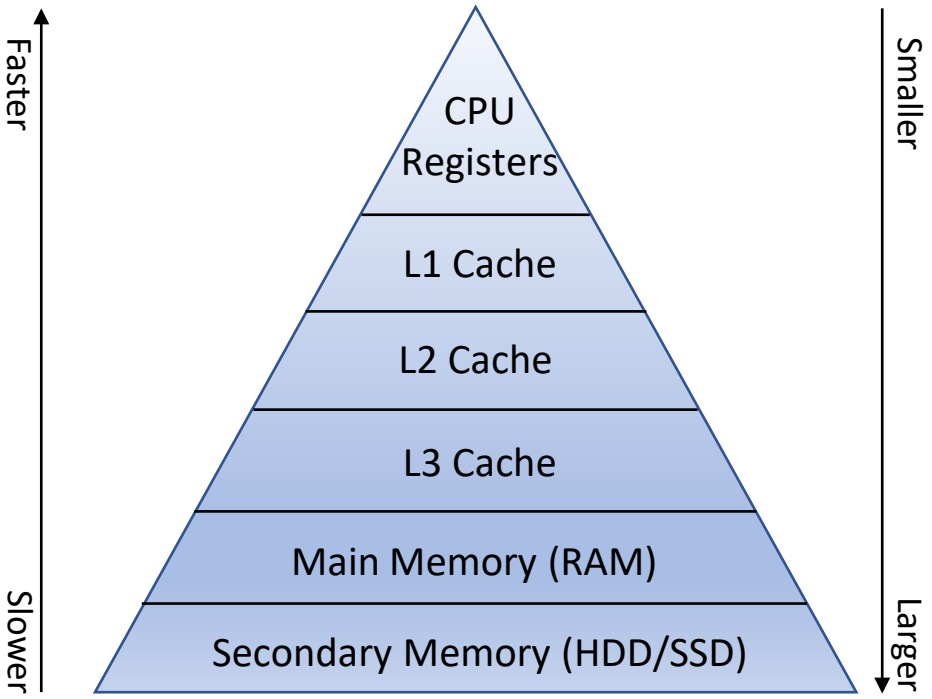


Diminisher: A Linux Kernel-based Countermeasure for TAA Vulnerability

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David Novo (LIRMM), Florent Bruguier (LIRMM), Pascal Benoit (LIRMM)

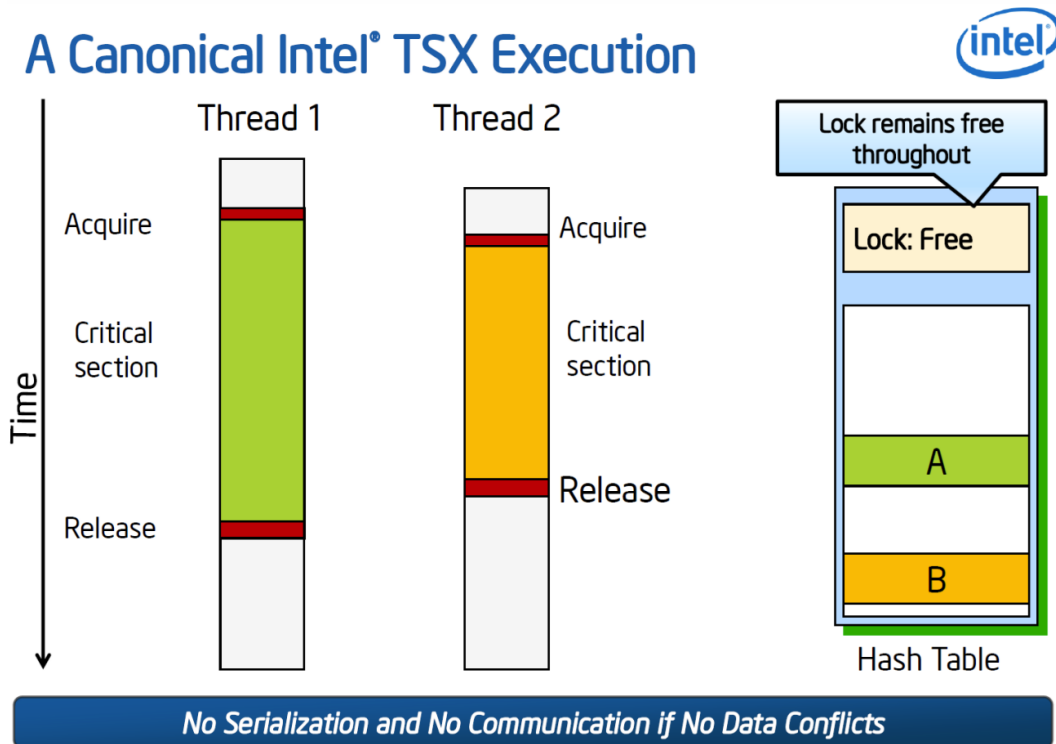
Background

Caches

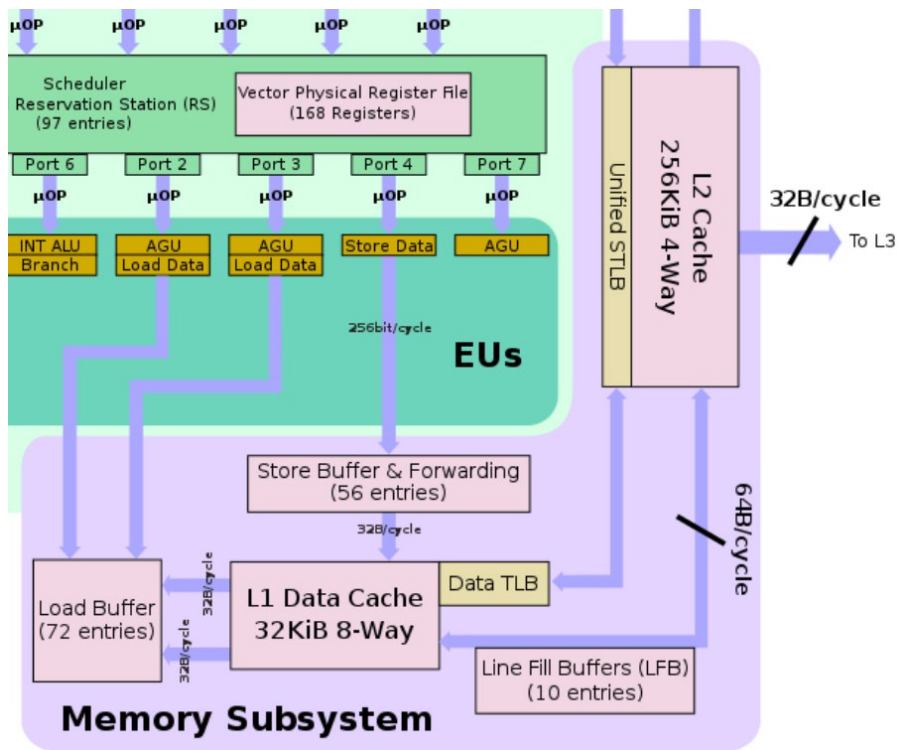


Intel TSX

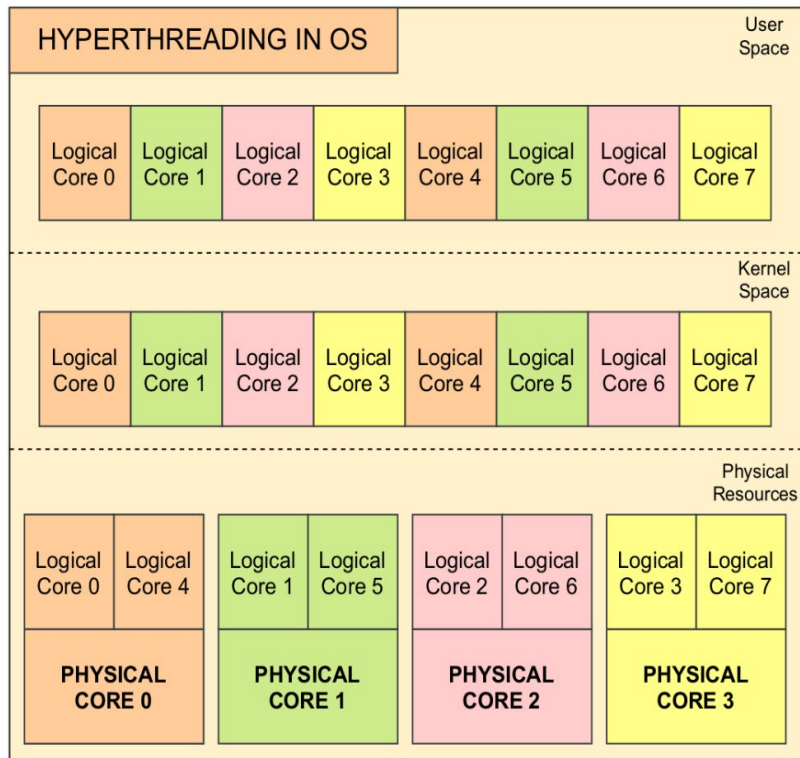
A Canonical Intel® TSX Execution



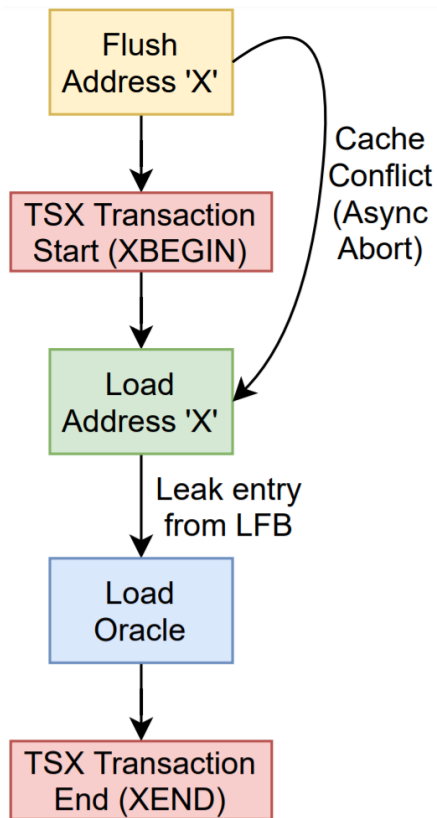
Line Fill Buffer



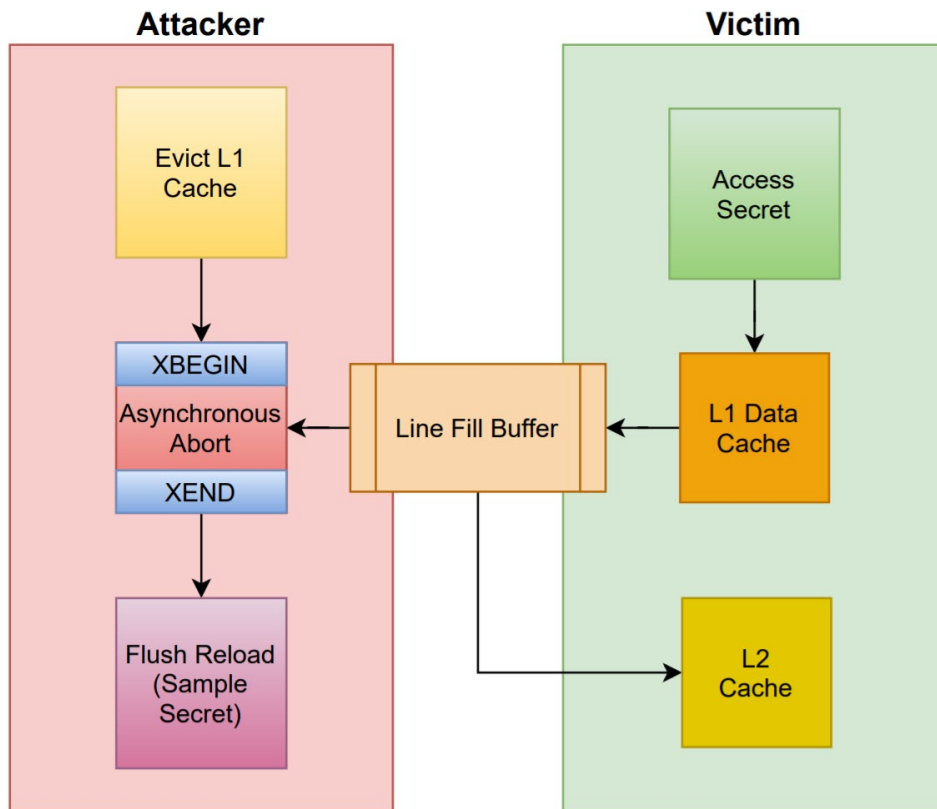
Intel Hyper-threading



Asynchronous Aborts



TAA Vulnerability



Related Papers & Existing mitigation

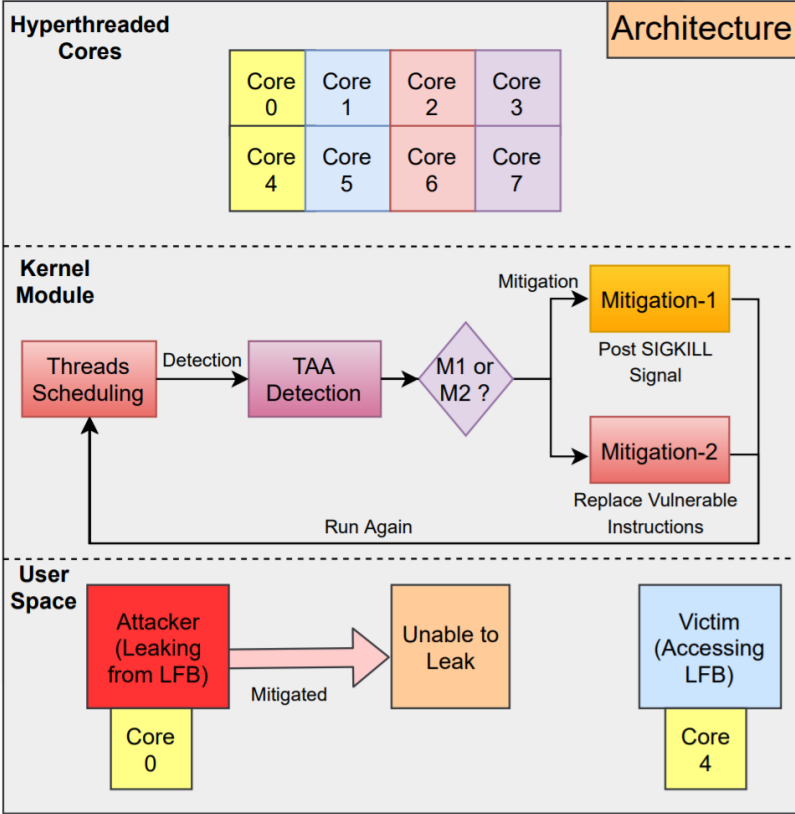
- CacheOut, RIDL and Zombieload
- VERW mitigation
- No solution for hyperthreaded TAA attacks

Contributions

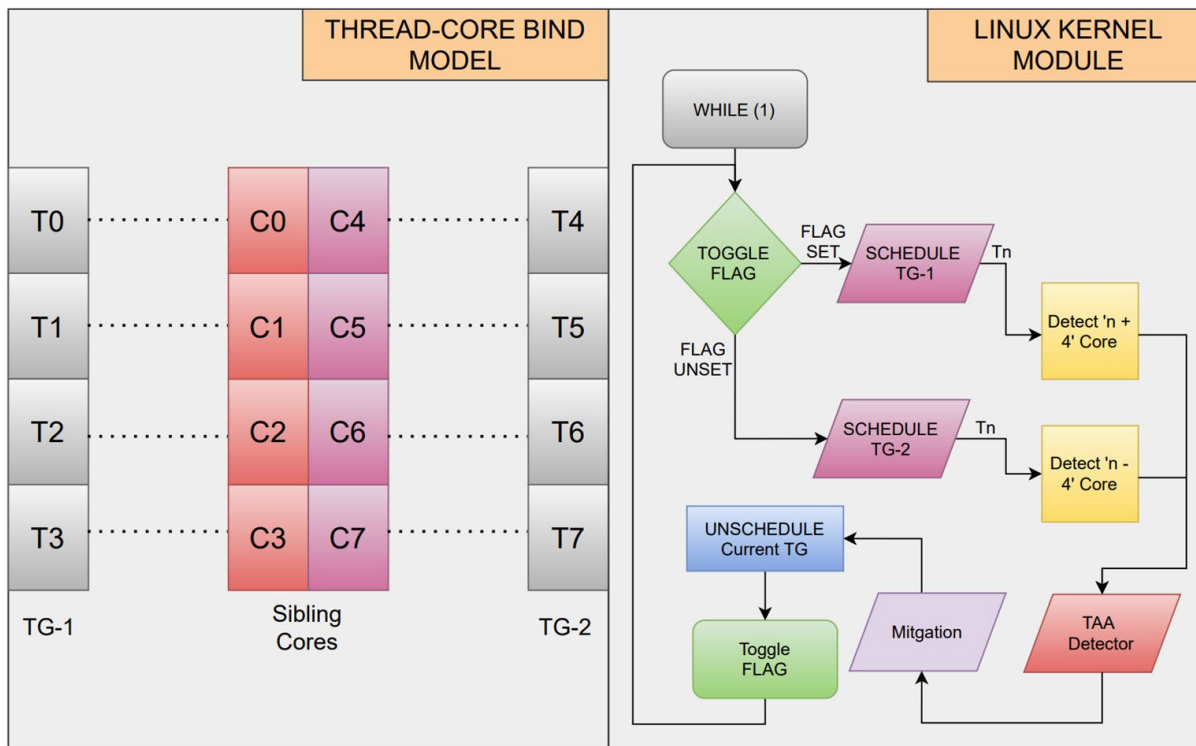
- Countermeasure for TAA vulnerability
- Hyperthreaded & Cross-cores TAA attacks
- Efficient solution
- Scalable solution
- Noise resilient

Methodology

Methodology



Scheduling



Detection

- Analyzing cache conflicts
- Feature Selection
- Threshold based Detection

Detection: Feature Selection

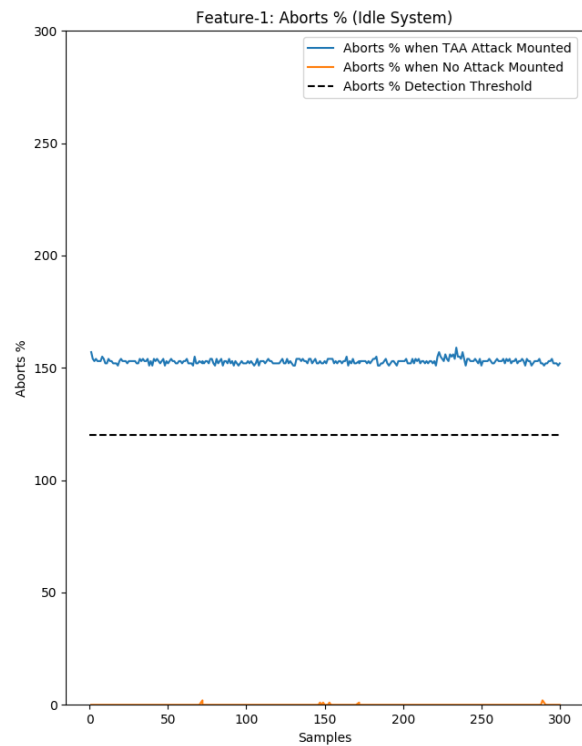
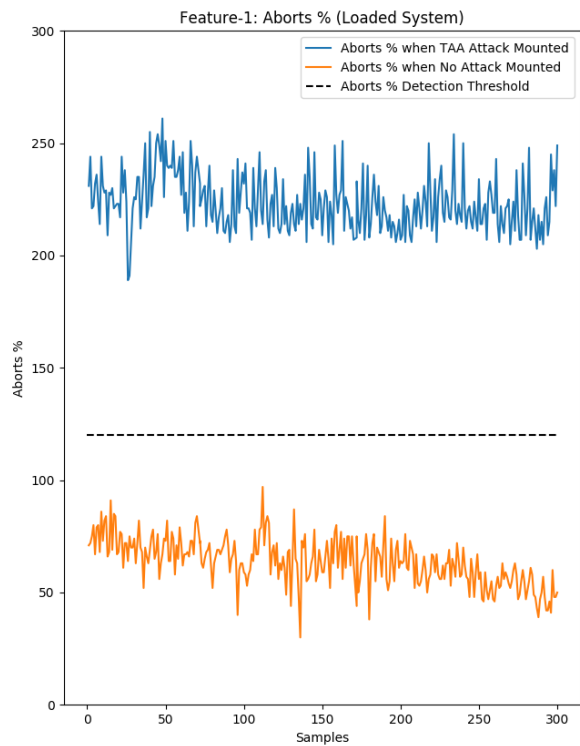
- Feature-1
 - Total number of cache conflicts
- Feature-2
 - Maximum number of cache conflicts occurred on a cache set
- Feature-3
 - Minimum number of cache conflicts occurred on a cache set

Detection: Algorithm

Algorithm 1: TAA Attack Detection

```
abrt ← 0, no_abrt ← 0, abrt_per_set[SETS] ← 0, samples ← 0
while samples ≤ 50000 do
  for set ← 0 to 64 by 1 do
    Begin_TSX_Transaction();
    Access_All_Ways(set);
    End_TSX_Transaction();
    if ABORT_REASON_CACHE_CONFLICT then
      | abrt ← abrt + 1;
      | abrt_per_set[set] ← abrt_per_set[set] + 1;
    else if NO_ABORT then
      | no_abrt ← no_abrt + 1;
  samples ← samples + 1;
  Sleep_Detection();
if ((abrt / no_abrt) * 50000 ≥ F1_THRESHOLD &&
  Get_Max(abrt_per_set) ≥ F2_THRESHOLD &&
  Get_Min(abrt_per_set) ≥ F3_THRESHOLD) then
  | /* Attack Detected */
  | return 1;
/* No-Attack Detected */
return 0;
```

Detection: Threshold Calculation



Mitigation

- Mitigation-1 (SIGKILL to Process):
 - Posts SIGKILL signal from kernel to Vulnerable process to kill it
- Mitigation-2 (Instruction Replacement):
 - Replaces vulnerable instruction that causes the attack within Linux Kernel

Mitigation: SIGKILL to Vulnerable Process

Algorithm 2: Mitigation-1 (SIGKILL to Vulnerable Process)

```
if Is_TAA_Detected() then  
    task_struct  $\leftarrow$  Get_Task_Struct();  
    Task_Lock(task_struct);  
    ret  $\leftarrow$  Send_SIG_KILL_Signal(task_struct);  
    Task_Unlock(task_struct);  
    if ret == SUCCESS then  
        /* Killed vulnerable process */  
        return 1;  
    /* Unable to kill vulnerable process */  
    return 0;
```

Mitigation: Vulnerable Instructions Replacement

Algorithm 3: Mitigation-2 (Vulnerable Instruction Replacement)

```
if Is_TAA_Detected() then
    task_struct ← Get_Task_Struct();
    Task_Lock(task_struct);
    user_pages ← Read_Text_Section(task_struct);
    text_ptr ← Map_Kernel_Space(user_pages);
    for i ← 0 to code_size by 1 do
        if text_ptr[i] == VULNERABLE_INSTRUCTION then
            Replace_To_NOP(text_ptr[i]);
            ret ← SUCCESS;
    Unmap_And_Release_Pages(user_pages);
    Task_Unlock(task_struct);
    if ret == SUCCESS then
        /* Replaced Vulnerable Instruction */
        return 1;
    /* Unable to find Vulnerable Instruction */
    return 0;
```

Results

Experimental Results

System State	Type	Accuracy (%)	FP (%)	FN (%)	Overhead (%)	Latency (us)
IDLE	Detection	97.31	2.64	0.03	2.5	5264232
	Mitigation-2	99.94	0.03	0		306
	Mitigation-1	99.85	0.03	0.18		86
LOADED	Detection	98.26	1.73	0.03	2.5	6916110
	Mitigation-2	99.91	0.03	0.06		452
	Mitigation-1	99.82	0.03	0.21		106

Conclusion

- Countermeasure for TAA Vulnerability
- Works for both cross & hyper-threaded cores
- High accuracy, low overhead and low latency
- Standalone Mitigation
- Scalable solution